40th IEEE International Conference on Data Engineering

Joint Mobile Edge Caching and Pricing: A Mean-Field Game Approach







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Background & Motivation

>Related Work & Problem Formulation

Basic Idea & Solution

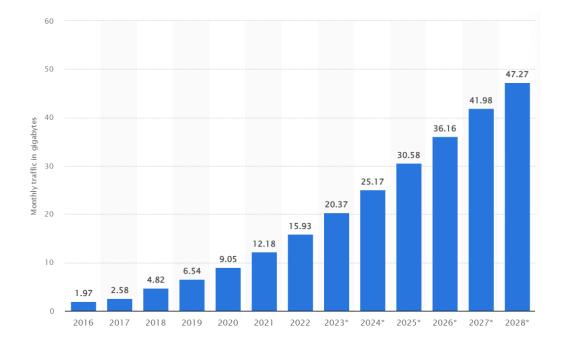
Evaluation & Conclusion





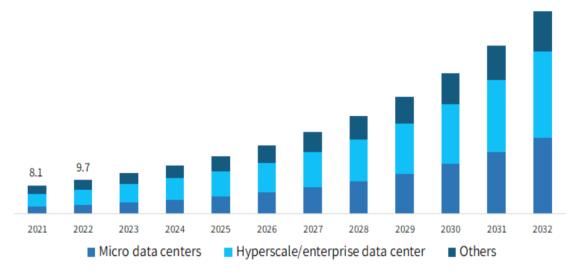
Explosive growth of mobile data traffic

- Fast deployment of edge devices
- 10x increase from 2016 to 2023
- Significant stress on back-haul links



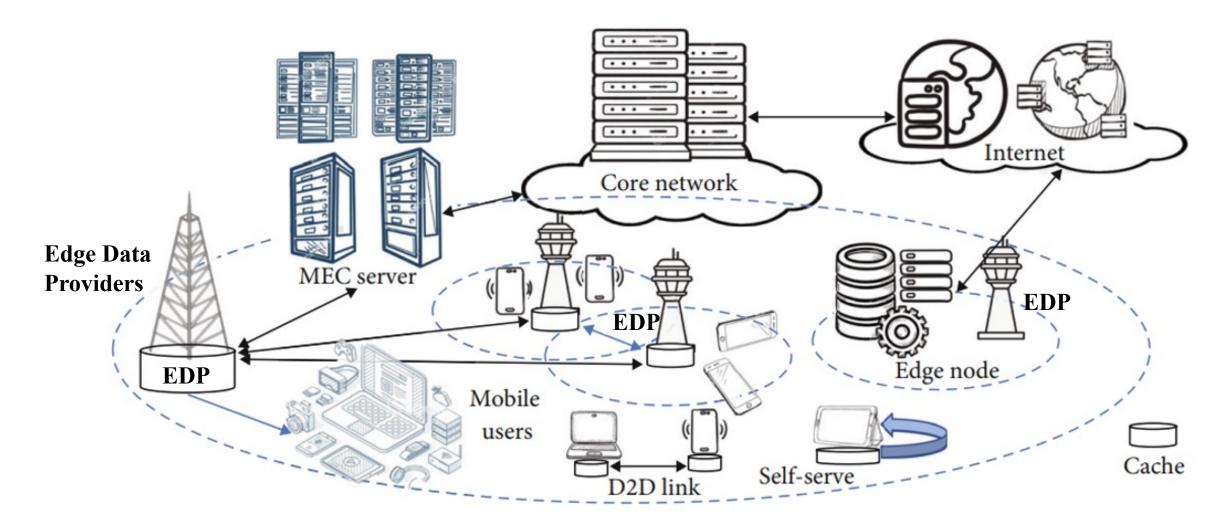
- Content caching and trading in the edge
 - By 2032, the market size of edge data centers is expected to exceed \$50 billion
 - Enhance Quality of Experience (QoE)

Global Edge Data Center Market Size, By Data Center, 2021-2032, (USD Billion)





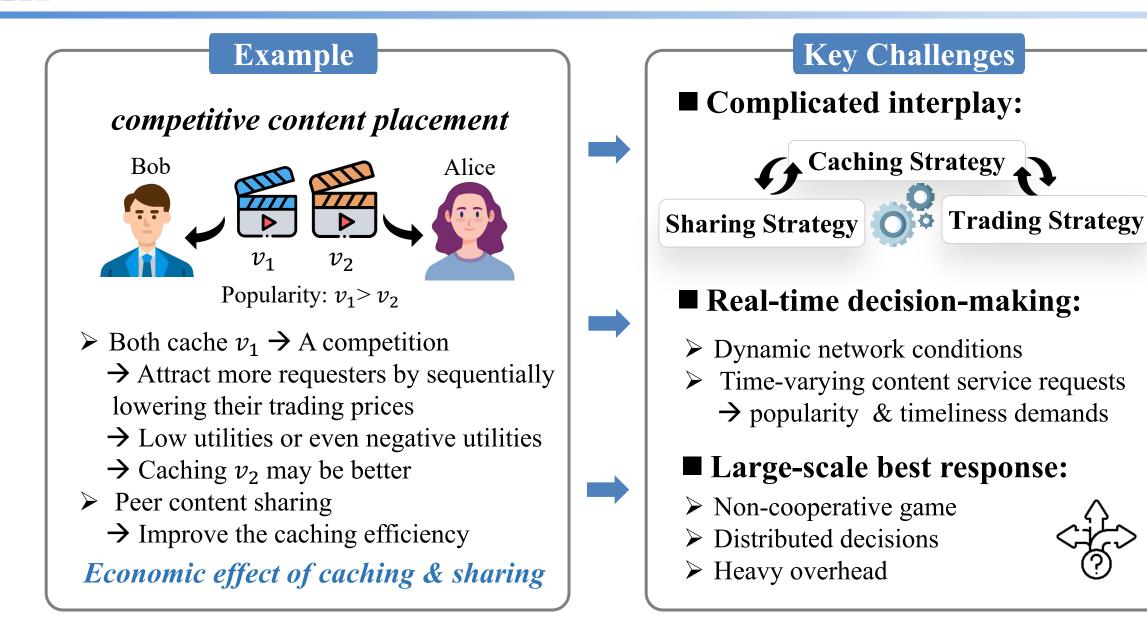




General architecture of Mobile Edge Caching systems

Motivation









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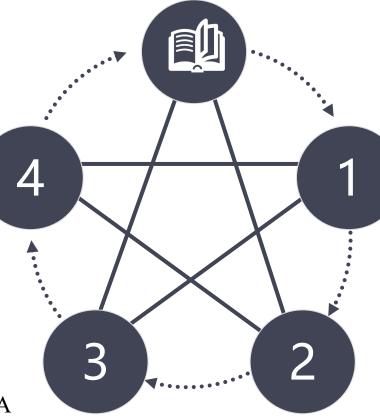
Goal: edge caching optimization while considering pricing and content sharing

Mean-Field Game

[5] D. Narasimha, et al., "Agedependent distributed MAC for ultra-dense wireless networks," IEEE INFOCOM2021
[6] H. Feng, et al, "Mean-field game theory based optimal caching control in mobile edge computing," IEEE TMC2022.

Edge caching and pricing

[4] J. Zou, et al, "Joint pricing and cache placement for video caching: A game theoretic approach," IEEE JSAC2019.



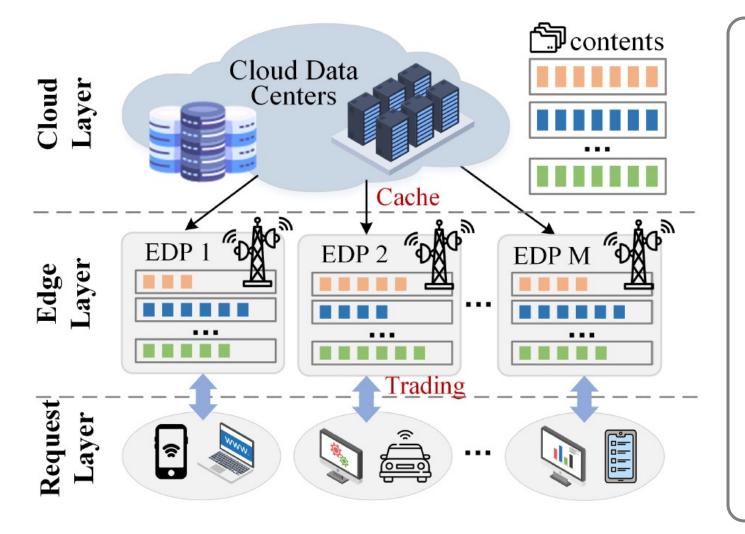
Mobile Edge Caching

[1] H. Sun, et al, "A proactive ondemand content placement strategy in edge intelligent gateways," TPDS2023.
[2] T. Zong, et al, "Cocktail edge caching: Ride dynamic trends of content popularity with ensemble learning ," IEEE/ACM ToN2023.

Content Trading

[3] Y. Huang, et al, "Profit sharing for data producer and intermediate parties in data trading over pervasive edge computing environments," IEEE TMC2023. **System Model**





System Model

- Edge Data Providers (EDPs) {1,...,i,...,M}: cache certain contents in advance and are allowed to share the cached contents with each other.
- Content requesters {1,...,j,...,J}: buy desired contents from EDPs and pay a suitable trading price.
- Contents $\{1, \dots, k, \dots, K\}$: the data size of content k is denoted by Q_k .
- Caching strategy vector of EDP *i*: $x_i(t) = [x_{i,1}(t), \dots, x_{i,K}(t)] \rightarrow$ the instantaneous caching rate of contents
- Remaining Storage capacity of *i*: $q_i(t) = [q_{i,1}(t), \dots, q_{i,K}(t)]$





> Network Model: the evolution of $h_{i,j}(t) \rightarrow$ the channel fading coefficient between EDP *i* and requester *j*

 $\mathbf{d}h_{i,j}(t) = \frac{1}{2}\varsigma_h(v_h - h_{i,j}(t))\mathbf{d}t + \varrho_h\mathbf{d}W_{i,j}(t).$ $\varsigma_h: \text{ changing rate } W_{i,j}(t): \text{ a standard Brownian motion}$ $v_h \& \varrho_h: \text{ the long-term mean and standard deviation}$

➢ Edge Caching Model: Content Popularity + Timeliness → The dynamics of the caching state

 $\mathbf{d}q_{i,k}(t) = Q_k[-w_1 x_{i,k}(t) - w_2 \prod_{i,k}(t) + w_3 \xi^{L_{i,k}(t)}]\mathbf{d}t + \varrho_q \mathbf{d}W_i(t)$

Content popularity: the frequency at which the content k is requested

Content timeliness: the level of urgency with which requesters acquire content *k*



The fewer the requests for content k, the faster it is removed from the caching storage. The more urgent the requests for content k, the faster it is added into the caching storage.

➤ Trading Model: $p_{i,k}(t)$ → the unit price customized by EDP *i* for selling content *k* Sharing price \bar{p}_k → a uniform unit price for obtaining content *k* from other peer EDPs







> Trading income: the dynamic price * the number of requesters

$$\Phi_{i,k}^{1}(t) = I_{i,k}(t)p_{i,k}\mathbb{P}^{1}\left(Q_{k} - q_{i,k}(t)\right) + I_{i,k}(t)p_{i,k}\mathbb{P}^{2}\left(Q_{k} - q_{-i,k}(t)\right) + I_{i,k}(t)p_{i,k}\mathbb{P}^{3}Q_{k}$$
The set of requesters who
ask for content k at time t
$$p_{i,k}(t) = \begin{cases} \hat{p}, & M = 1 \\ \hat{p} - \frac{\eta_{1}\sum_{i'=1,i'\neq i}^{M}Q_{k} \cdot x_{i',k}(t)}{M-1}, & M \ge 2 \end{cases}$$
The principle
of supply and demand

Sharing benefit: the monetary benefit formed when EDP *i* shares its contents with other EDPs $\mathcal{M}_{i,k}(t)$

$$\Phi_{i,k}^2(t) = \sum_{i' \in \mathcal{M}_{i,k}(t)} \overline{p}_k \left(q_{i',k}(t) - q_{i,k}(t) \right)$$



 $C_{i,k}^3(t$



> Content placement cost $C_{i,k}^{1}(t) = w_4 x_{i,k}(t) + w_5 x_{i,k}^{2}(t)$ [Placing and storing cached contents]

Staleness cost: a penalty function of the total request service delay

$$C_{i,k}^{2}(t) = \eta_{2} \left\{ \frac{Q_{k}x_{i,k}(t)}{H_{c}} + \sum_{j \in I_{i,k}(t)} \left[\mathbb{P}^{1} \frac{Q_{k} - q_{i,k}(t)}{H_{i,j}(t)} + \mathbb{P}^{2} \frac{Q_{k} - q_{-i,k}(t)}{H_{i,j}(t)} + \mathbb{P}^{3} \left(\frac{q_{i,k}(t)}{H_{c}} + \frac{Q_{k}}{H_{i,j}(t)} \right) \right] \right\}$$

Sharing cost

$$) = \mathbb{P}^{2} \overline{p}_{k} \left(q_{i,k}(t) - q_{-i,k}(t) \right)$$
 Buy content k from an adjacent EDP

Game

$$\begin{aligned}
\max_{x_{i,k}(t),t:0\to T} \mathcal{U}_{i,k}(x_{i,k}, S_{i,k}, S_{-i,k}) &= \Phi_{i,k}^{1}(t) + \Phi_{i,k}^{2}(t) - C_{i,k}^{1}(t) - C_{i,k}^{2}(t) - C_{i,k}^{3}(t) \\
s. t., \quad \mathbf{d}_{i,j}(t) &= \frac{1}{2}\varsigma_{h}(v_{h} - h_{i,j}(t))\mathbf{d}t + \varrho_{h}\mathbf{d}W_{i,j}(t) \\
\mathbf{d}_{i,k}(t) &= Q_{k}[-w_{1}x_{i,k}(t) - w_{2}\Pi_{i,k}(t) + w_{3}\xi^{L_{i,k}(t)}]\mathbf{d}t + \varrho_{q}\mathbf{d}W_{i}(t)
\end{aligned}$$





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≻Basic Idea & Solution

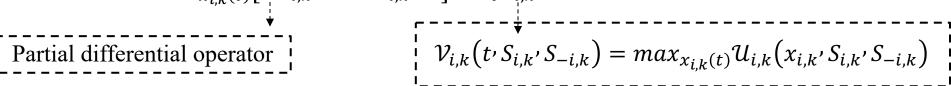
Evaluation & Conclusion





Solving a series of Hamilton-Jacobi-Bellman (HJB) equations \succ

 $max_{x_{i,k}(t)} \left[\mathcal{OV}_{i,k}(t) + \mathbf{U}_{i,k}(t) \right] + \partial_t \mathcal{V}_{i,k}(t) = 0$



Solve



There exists a complicated coupling of the M equations; Unknown mutual influence: a dynamic pricing policy and a peer-to-peer sharing policy.



Approximate the **collective impact** of large-scale EDPs on caching, trading, and sharing, and then iteratively solve multiple coupled equations to determine the optimal caching strategy for each EDP.

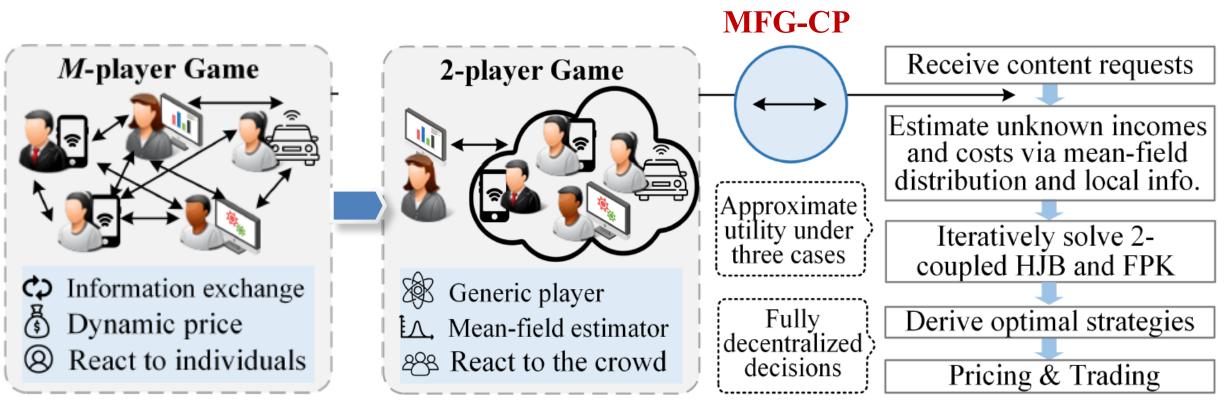
The original stochastic differential game \rightarrow A Mean-Field Game with a generic-player representation

A mean-field estimator \rightarrow Estimate the unknown **Estimate** incomes & costs of EDPs A generic player \rightarrow Adjust its caching strategy

An iterative best response learning algorithm to solve the coupled equations in a distributed way







- ✓ The number of partial differential equations in the MFG-CP is reduced from M×K to 2×K
 → The computing time required to derive the optimal strategies will not incur excessive complexity.
- ✓ The optimal caching strategy can be determined during the time horizon using only local information and a mean-field distribution.

>>>> Determining the Optimal Caching Strategy

λ

X



Mean-field Estimator

- Construct the Fokker-Planck-Kolmogorov (FPK) equation to estimate the statistical distribution of states $\lambda(S_k(t))$: $\partial_t \lambda(S_k(t)) + \frac{1}{2}\varsigma_h(\upsilon_h - h(t))\partial_h \lambda(S_k(t))$ $+Q_k[-w_1x_k(t) - w_2\Pi_k(t) + w_3\xi^{L_k(t)}]\partial_q \lambda(S_k(t))$ $-\frac{1}{2}\varrho_h^2 \partial_{hh}^2 \lambda(S_k(t)) - \frac{1}{2}\varrho_q^2 \partial_{qq}^2 \lambda(S_k(t)) = 0.$
- Estimate the dynamic price:

 $p_k(t) \approx \hat{p} - \eta_1 Q_k \int_h \int_q \lambda (S_k(t)) x_k^* (S_k(t)) \mathbf{d}h \mathbf{d}q$

- Estimate the average caching state: $\bar{q}_{-,k}(t) \approx \int_{h} \int_{q} q_k(t) \lambda(S_k(t)) \mathbf{d}h \mathbf{d}q_k$
- Estimate the average sharing benefit:

$$\overline{\Phi}_k^2(t) = \overline{p}_k \overline{\Delta q}(t) \left(\frac{M - M'_k(t)}{M_k(t)} - 1 \right)$$

Generic player

- Solve the optimization problem: $\max_{x_k(0\to T)} \quad \mathcal{U}_k(x_k, S_k, \lambda)$
- Employ the mean-field estimator to quickly obtain the trading income, sharing benefit, staleness cost, and sharing cost.
- Reconstruct HJB equation to determine the value function $\mathcal{V}_{i,k}(t)$:

 $\max_{x_k(t)} \left[\frac{1}{2} \varsigma_h(\upsilon_h - h(t)) \partial_h \mathcal{V}_k(t) + \frac{1}{2} \varrho_h^2 \partial_{hh}^2 \mathcal{V}_k(t) + Q_k [-w_1 x_k(t) - w_2 \Pi_k(t) + w_3 \xi^{L_k(t)}] \partial_q \mathcal{V}_k(t) + \frac{1}{2} \varrho_h^2 \partial_q^2 \mathcal{V}_k(t) + \mathbf{U} \left[(t - w_1 - \xi_h) \right] + \partial_h \mathcal{V}_h(t) = 0$

- $+\frac{1}{2}\varrho_q^2 \partial_{qq}^2 \mathcal{V}_k(t) + \mathbf{U}_k(t, x_k, S_k, \lambda) \bigg] + \partial_t \mathcal{V}_k(t) = 0.$
- The optimal caching strategy:

$$x_{k}^{*}(t) = \left[-\left(\frac{w_{4}}{2w_{5}} + \frac{\eta Q_{k}}{2H_{c}w_{5}} + \frac{Q_{k}w_{1}\partial_{q}\mathcal{V}_{k}(t)}{2w_{5}} \right) \right]^{+}$$



Algorithm 2: Iterative Best Response Learning Scheme **input** : The iterative number ψ ; iterative threshold ψ_{th} ; **output:** The optimal caching strategy; 1 Initialize: $\psi = 0$; $x_k^{\psi}(0)$; $\mathcal{V}_k^{\psi}((\sigma+1)T)$; λ ; 2 while $\psi < \psi_{th}$ do $\psi = \psi + 1;$ 3 Solve the HJB equation using $x_k^{\psi-1}(t)$ and λ ; 4 Update the strategy vector $x_k^{\psi}(t)$ based on Eq. (21); 5 if $|x_k^{\psi}(t) - x_k^{\psi-1}(t)| < a$ preset threshold then 6 Break the iterative learning process; 7 Solve the FPK equation using $x_k^{\psi}(t)$ to update λ ; 8 Update the mean-field estimator: $p_k(t), \bar{q}_{-,k}(t), \overline{\Delta q}(t);$ 9 Update the utility function according to Eq. (10); 1011 return $x_k^{\psi}(t)$ for $t \in [\sigma T, (\sigma+1)T)$.



Core: iteratively solve the coupled HJB equation and FPK equation: the backward HJB models the induction process of the optimization of each individual, while the forward FPK models the evolution of the mean-field as a whole.

The solution of the HJB equation has a great impact on the FPK equation to update the mean field.

⁺ The solution of the FPK equation is also required by the HJB equation to estimate the generic player's utility and update strategies.

Framework Description



A	Igorithm 1: The Proposed Framework (MFG-CP)	*		
	nput : The parameters $w_1 \sim w_5$, η_1 , η_2 , ς_h , v_h , ϱ_h , ϱ_q , and ξ ; the maximum epoch number σ_{max} ; nitialize: the state $S(0)$; the data size of each content Q_k ;	 Core: Each EDP makes decisions in parallel Mobile Edge Caching + Pricing and Trading 		
2 f	for each EDP $i = 1, 2, \dots, M$ in parallel do			
3	while each optimization epoch $\sigma \leq \sigma_{max}$ do			
4	Record requests $\{I_{i,k}(t) k \in \mathcal{K}, t \in [\sigma T, (\sigma+1)T)\};$			
5	Determine the content set $\mathcal{K}' = \{k q_{i,k}(t) < Q_k$	The change in requesters' demands occurs at a		
	and $\sum_{\sigma T}^{(\sigma+1)T} I_{i,k}(t) > 0$ } that needs to be cached;	relatively slow rate		
6	for each content $k \in \mathcal{K}'$ do			
7	// Mobile Edge Caching;			
8	Compute content popularity $\Pi_{i,k}(t)$ and con-			
	tent timeliness $L_{i,k}(t)$ for all contents;			
9	Call for Alg. 2 to obtain the optimal caching	Invokes Alg. 2 to give the best response:		
	strategy $x_{i,k}^*(t)$ as its best response;	customizing its optimal caching strategy		
10	Update the current state based on $x_{i,k}^*(t)$;			
11	// Pricing and Trading;	<u>.</u>		
12	Case 1: sell the content at a unit price $p_{i,k}(t)$;	The trading process between the EDP and the		
13	Case 2: buy the uncached content from an	corresponding requesters will be carried out,		
	adjacent EDP at \bar{p}_k and sell it to requesters;			
14	Case 3: download the uncached content from	where the EDP needs to take different actions		
	the center and sell it to requesters;	under various cases		
		·'		



Lemma

✓ In the MFG-CP framework, there exists the unique value function V_k(t) of the HJB equation.
 Proof: (i) The caching strategy space is a compact subset of R;
 (ii) The drift term of the state dynamics & the utility function: bounded+Lipschitz continuous.

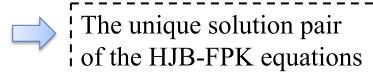
✓ There exists the **unique mean-field distribution** $\lambda(S_k(t))$ of the FPK equation. Proof: A parabolic partial differential equation $\rightarrow \partial_t \lambda(S_k(t)) + \Theta = d$

 $\Theta = -\sum_{i,j=1}^{k} a_{i,j} \lambda \left(S_i S_j \right) + \sum_{i,j=1}^{k} b_i \lambda \left(S_i \right) + c\lambda \rightarrow a_{i,j} = a_{j,i}$

Theorem

✓ In the MFG-CP framework, there exists a unique Nash equilibrium.
 Proof:

Given this contraction mapping, there exists a unique fixed-point based on the fixed-point theorem









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Real Dataset

♦ YouTube → The number of requests for each category
→ content id, tags, views, comment count, description.

Parameter settings

- ♦ M=300 EDPs, K=20 contents
- The conversion parameter η change from [0.1, 0.4]
- $\blacklozenge \ \lambda(0) \sim \mathcal{N}(0.7, 0.1^2)$



Experimental Settings

- Random Replacement (RR)
 Most Popular Caching (MPC)
 MFG: do not consider sharing
- WIPO: do not consider sharing
 Ultra-Dense Caching Strategy

(UDCS): consider content overlap

Compared Algorithms

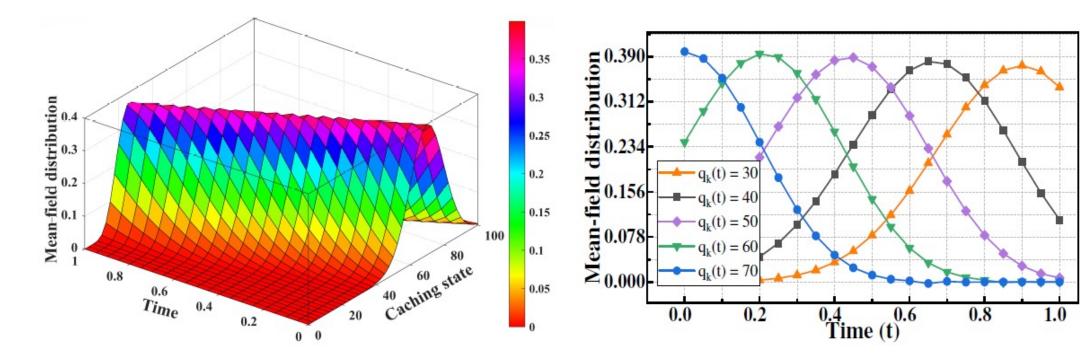
- ◆ Mean-Field Equilibrium
- ◆ Caching state & Utility
- Computation time







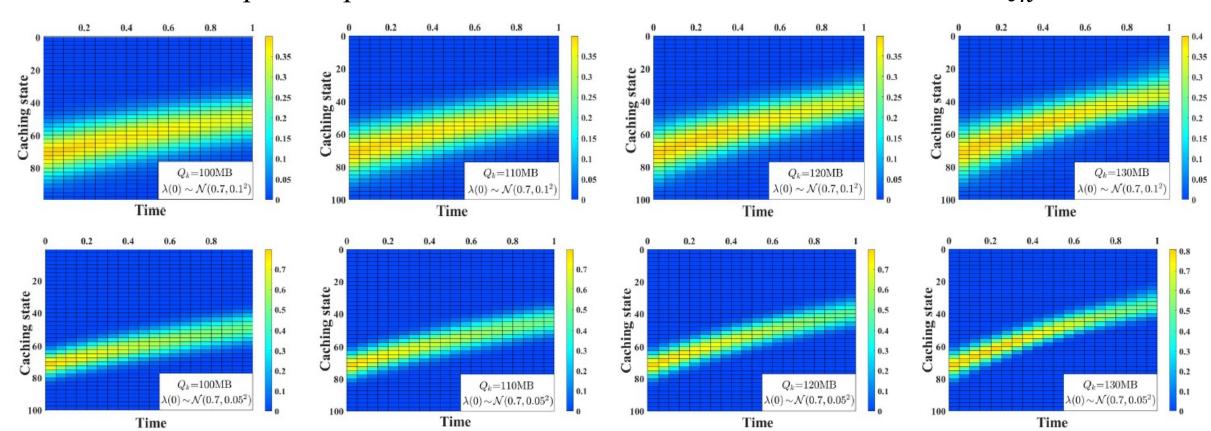
• Evolution of the mean-field distribution at the equilibrium:



When we fix the time slot, the size of the remaining caching space will increase first and then decrease. As the time evolves, the remaining caching space with {60MB, 70MB} will vanish due to the improvement of space utilization.

Evaluation





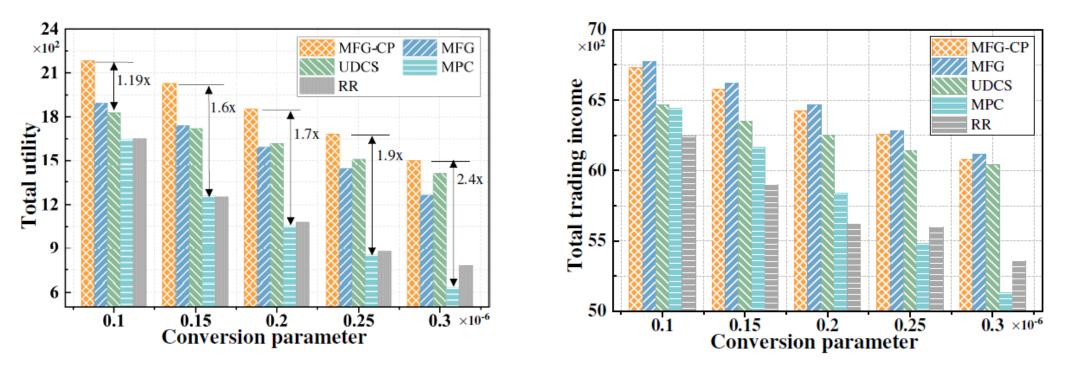
A heat map description of the mean-field distribution under different Q_k :

The caching space will gradually reach saturation with the increase of Q_k . When decreasing the variance values, the heat map displays more concentrated results.





□ Comparisons under different conversion parameters:

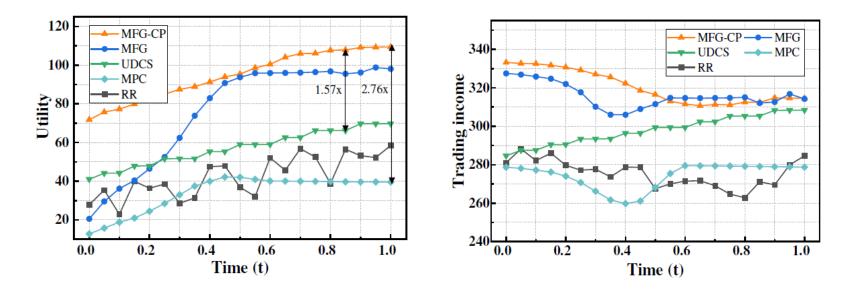


- Improving the value of the conversion parameter leads to a reduction in total utility;
- The staleness cost of MFG obviously exceeds that of MFG-CP;
- The behavior of content sharing is beneficial for improving the utility of each EDP;
- MFG-CP enables each EDP to possess a higher utility than these compared algorithms.





□ Comparisons on utility, trading income, and time:



The proposed framework MFG-CP offers significant advantages in maximizing the utilities of EDPs

COMPARISONS ON COMPUTATION TIME (SECOND)

Number Methods	50	100	200	300
MFG-CP	0.4319	0.4442	0.4336	0.5121
RR	0.1697	0.5527	0.9766	1.7832
MPC	0.1657	0.3157	0.8694	1.7094

The computational complexity of MFG-CP does not increase with the number of EDPs





- ✓ We study the competitive content placement issue in large-scale dynamic MEC systems.
- ✓ To facilitate decentralized decision-making, we propose the <
 MFG-CP framework for joint content caching and pricing.
- ✓ Extensive simulations on realworld traces validate its great performance.



- ✓ We model the problem as a noncooperative stochastic differential game, which considers the heterogeneous content demands, real-time trading prices, and paid content sharing.
- ✓ We develop an iterative best response learning scheme to determine the optimal caching strategy for each EDP.



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Thank you for your attention!

Question?

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