### Maximizing the User's Benefit in Mobile Cloud Computing

MULTIN MITTIT

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# Mobile Cloud Computing

- Concept of Mobile Cloud Computing (MCC):
  - Offload the computation complexity task from mobile devices to the cloud
  - Combination of cloud computing, mobile computing, and wireless networks to bring rich computational resources to mobile users





# Mobile Cloud Computing

- Different computation abilities
  - Cloud (e.g., Amazon EC2, Azure):
    Rich computation resource --- short processing delay
    Omitted in most scenarios
  - Mobile devices:
    - Limited computation resource --- long processing delay
      E.g., natural language processing, face recognition
- Cost of cloud computing
  - extra offloading/transmission delay
    - Cellular networks (low bandwidth, high energy consumption)
    - Wi-Fi (high bandwidth, low energy consumption)



# Mobile Cloud Computing

- Task finishing time
  - Task utility decay (reward for finishing the task)
    - E.g. hot news spreading, Siri
      - $\Box$  Linear decay model in this paper,  $U_t = U_0 wt$
- Trade-off in the MCC:
  - Smaller task-finishing time
  - More energy consumption is possible

#### • Challenge

- Maximize the gained utility with the given energy budget



# Problem Formulation

#### Network Model

- Limited battery of the mobile device
- Non-preemptive task
  - Cannot be interrupted once it is scheduled

### Scheduling problem

- Given the arrival time, local processing time, transmission time of each task, and the battery constraint of the smartphone, find a scheduling so that the utility gain is the maximum
- NP complete



### Challenges

#### A motivational example

- Every second, a new task arrives,  $U_t = 10 2t$
- Each task consumes 1 energy unit in local or 2 units to transmit
- The energy budget is 4 units







- How can we optimize the mobile device and the cloud scheduling jointly?
  - Traditional approaches will lead to a bad performance in certain scenarios







- Simple greedy scheduling does not have a performance bound
  - Cost effective

Unbalanced task assignment

- Balance the finishing time
  - □ Low energy usage efficiency

 We propose an approximation algorithm with a good performance bound!





### Solution

- LP Rounding algorithm
  - Assume tasks can be interrupted
    - $\hfill \mathbf{y}_{ijt}$  , task i is scheduled at j at time t
    - $\hfill \square$  Problem reduces into the linear programming





### Solution

- LP Rounding algorithm
  - Calculate the LP rounding result
  - Assign the task i to j at time t with a probability based on y<sub>ijt</sub>.
  - Each job is assigned at the earliest feasible time

Theorem: the LP rounding algorithm has a 2approximation ratio on expectation



# Experimental Setting

- Task arrival time
  - Average data size is 2 MB, 1 to 3 MB.
  - Follows the updates in Sina Weibo trace
    - Power-law task arrival time
- Smartphone

- Processing time/energy is 1 MBps, 500 mW.
- Transmission bandwidth/energy 2 MBps, 2000 mW.
- 100 mWh energy budget



# Algorithm Comparison

### • Algorithms:

- Cloud-only (CO) algorithm, which only utilizes the cloud for computation, referred to as the AllServer algorithm in [5]
- Smartphone-only (SO) algorithm, which only utilizes the smartphone for computation, and is referred to as the AllMobile algorithm in [5]
- Random (RD) algorithm, which randomly assigns the new task to a device
- Proposed LP rounding algorithm



### Simulation Result



SO and CO algorithms only achieve a good performance in certain scenarios. The proposed LP algorithm achieves a good performance or the best performance in a wide scenario. The rounding algorithm leads to a bad performance.



- We investigate mobile cloud offloading in a general scenario, where both mobile and cloud are used to maximized the gained utility
- An LP rounding algorithm is proposed with the performance bound

- Future Work
  - Online scheduling, real testbeds.



# Thank you!

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