

Incentive-Driven and Freshness-Aware Content Dissemination in Selfish Opportunistic Mobile Networks

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Presenter: Huan Zhou



- Introduction
- Motivations
- Network Model
- Proposed Scheme ConDis
- Performance Evaluation
- Conclusion

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Introduction

- Delay Tolerant Networks, Intermittently Connected Networks
- Consists of portable devices with bluetooth or wifi, such as smart phone, PDAs, laptops
- Intermittent connectivity
 - High dynamic time-varying topology
 - ➤ Contact: two nodes within the transmission range of each other
 - > Store-carry-forward

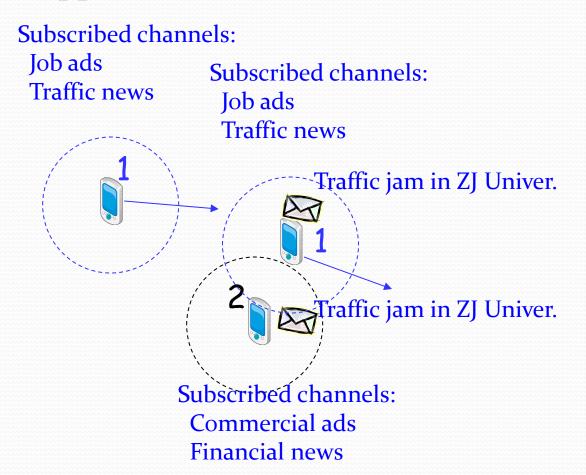
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- Content Dissemination: publish/subscribe scheme.
- Pub/sub scheme: high flexibility and adaptability when dealing with highly dynamic network topologies.
- Subscriber: express their interest without knowledge of the content generators' specific ID.
- Publisher: generate contents to the network without specifying the destination ID.
- Goal: deliver contents from publishers to subscribers.
- Contents to be disseminated fall into several channels:
 - > Sport news, Traffic news, Pop music and so on

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 An Example of Pub/sub Content Dissemination in OppNets.





Subscribed channels:
Movie news
Financial news

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Motivations

- Behaviors of nodes:
 - Cooperative: voluntarily store, carry, and forward others contents
 - ➤ Selfish: only carry its own interested contents; not be willing to share its resources
- Existing works addressing selfishness:
 - ➤ Reputation-based; Credit-based; and Tit-For-Tat (TFT)-based schemes.
 - > TFT only requires the principle of equal amounts of service.
- How to stimulate selfish nodes to participate into content dissemination and improve network performance under the TFT scheme?

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Network Model

- Node Contact Model: pair-wise inter-contact time realistic traces follows an exponential distribution.
- Channel and Content Model:
 - Nodes need to express their interests towards different kinds of contents in a certain way, and accordingly subscribe to those channels;
 - \triangleright Each published content includes (d, c, T_d , T);
- Assumptions:
 - ➤ The buffer size of nodes is the same;
 - Contents have the same volume capacity;

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-Architecture of ConDis:



- The Architecture of ConDis:
 - ➤ Subscribed Channel Manager: keeps the channel information subscribed by its one-hop neighbors.
 - ➤ Buffer State Collector: collects the content information stored at its one-hop neighbors.
 - ➤ Content Utility Estimator: gives the definition of content utility.
 - ➤ Buffer Manager: manages contents in the buffer, and the cache management of a node is mainly based on the content utility.

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-Content Utility Estimator

- Each content has an initial freshness value *V* when it is published by a certain node.
- The freshness value of a certain content *d* for a subscribed node *i* can be expressed as:

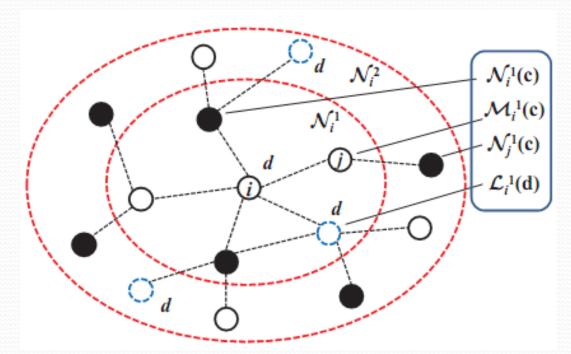
$$v_i(d) = \frac{R_d V}{T},$$

• R_d is the remaining valid time of content d; T is the TTL of content d.

-Content Utility Estimator

• Illustrating the utility definition:





- We divide node *i*'s one-hop neighbors into three kinds:
 - \triangleright Black-solid nodes $\mathcal{N}_i^1(c)$: interested in content d;
 - \triangleright Black-hollow nodes $\mathcal{M}_i^1(c)$: not interested in content d;
 - \triangleright Blue-dashed nodes $\mathcal{L}_i^1(d)$: already have obtained content d_i :

-Content Utility Estimator



- The first kind will absolutely choose to request content d from node i;
- The second kind may also request content *d* from node *i*;
- $\mathcal{N}_i^1(c)$ \longrightarrow direct subscribed value $U_{di}(d)$.
- $\mathcal{M}_{i}^{1}(c)$ indirect subscribed value $U_{indi}(d)$.
- The utility of a content *d* in channel *c* for node *i* is defined as:

$$U_i(d) = wU_{di}(d) + (1-w)U_{indi}(d),$$

• *w* in the range of [0, 1].

-Content Utility Estimator

- For node $j \in \mathcal{N}_i^1(c)$, j will request content d from node i only when the following three conditions are met:
 - ➤ Node *j* will not receive content *d* from its one-hop neighbors before this content is time out.
 - ➤ Node *i* can deliver content *d* to *j* before this content is time out.
 - ➤ When content *d* is delivered to node *j* , node *i* should guarantee that this content is fresh.
- Then, we express the direct subscribed value $U_{di}(d)$ as:

$$U_{di}(d) = \sum_{j \in \mathcal{N}_i^1(c)} \frac{(R_d - ED_{ij})V}{T} \prod_{k \in \mathcal{L}_j^1(d)} [1 - Pr_{jk}(R_d)],$$
(3)

-Content Utility Estimator

Similarly, the indirect subscribed value U_{indi}(d) can be expressed as:

$$U_{indi}(d) = \sum_{j \in \mathcal{M}_i^1(c)} \sum_{k \in \mathcal{N}_j^1(c)} \frac{(R_d - ED_{ik}^j)V}{T}, \quad (4)$$

- Here, node *j* is *i*'s one-hop neighbors which do not subscribe to channel *c* and have not obtained content *d*. Node *k* is j's one-hop neighbors which subscribe to channel *c* and have not obtained content *d*.
- We ignore node k's other one-hop neighbors.

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- Buffer Manager
- Objective of nodes in the network:



$$Max \ U_i = \sum_{c=1}^{C} (\sum_{d \in \theta(c)} U_i(d) - \sum_{d \in \phi(c)} U_i(d)),$$

• U_i is the utility function of node i; C is the total number of channels; and $\theta(c)$ and $\phi(c)$ are the set of contents associated to channel c in their buffer after and before exchange, respectively.

-Computing Content Utility



- Contact Probability Prediction: the pair-wise inter-contact time follows an exponential distribution.
 - ➤ PDF of the inter-contact time X_{ij} between nodes i and j can be expressed as:

$$f_{X_{ij}}(x) = \lambda_{ij} e^{-\lambda_{ij}x}.$$
 (8)

Then, the contact probability between nodes i and j within the remaining valid time Rd of content d can be expressed as:

$$Pr_{ij}(R_d) = Pr(X_{ij} \le R_d) = \int_0^{R_d} f_{X_{ij}}(x) dx$$

= 1 - $e^{-\lambda_{ij}R_d}$.

 $> \lambda_{ij}$ is the contact rate between nodes *i* and *j*.

-Computing Content Utility

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- Expected Delay Prediction:
 - ➤ Based on Eq. (8), the expected delay Ed_{ij} for transmitting a certain content from node i to j can be calculated as:

$$ED_{ij} = E[X_{ij}] = \int_0^\infty x f_{X_{ij}}(x) dx = \frac{1}{\lambda_{ij}}.$$

➤ the total time to transfer a content from *i* to *k* through *j* is $X^{j}_{ik} = X_{ij} + X_{jk}$, the PDF $f_{X^{j}_{ik}}(t)$ can be calculated as:

$$f_{X_{ik}^{j}}(x) = f_{X_{ij}}(x) \otimes f_{X_{jk}}(x)$$

$$= \lambda_{ij}\lambda_{jk} \int_{0}^{x} e^{-(\lambda_{ij} - \lambda_{jk})t} e^{-\lambda_{jk}t} dt$$

$$= \frac{\lambda_{ij}\lambda_{jk}(e^{-\lambda_{ij}x} - e^{-\lambda_{jk}x})}{\lambda_{jk} - \lambda_{ij}},$$
(11)

 $\triangleright \otimes$ is the convolution operator

-Computing Content Utility

• The expected delay Ed^{i}_{ik} for transmitting a certain content from i to k through j can be expressed as:

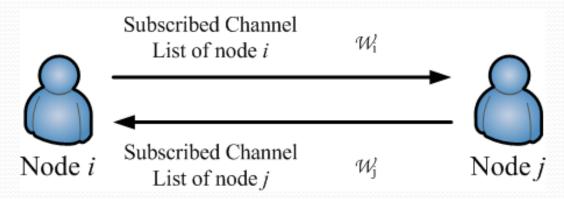
$$\begin{split} ED_{ik}^j &= E[X_{ik}^j] = \int_0^\infty x f_{X_{ik}^j}(x) dx \\ &= \int_0^\infty \frac{x \lambda_{ij} \lambda_{jk} (e^{-\lambda_{ij}x} - e^{-\lambda_{jk}x})}{\lambda_{jk} - \lambda_{ij}} dx \\ &= \frac{\lambda_{ij} + \lambda_{jk}}{\lambda_{ij} \lambda_{jk}}. \end{split}$$

• Then, the content utility $U_i(d)$ of content d in channel c for node i can be expressed as follows:

$$U_{i}(d) = w \sum_{j \in \mathcal{N}_{i}^{1}(c)} \frac{(R_{d} - \frac{1}{\lambda_{ij}})V}{T} \prod_{k \in \mathcal{L}_{j}^{1}(d)} e^{-\lambda_{jk}R_{d}}$$
$$+ (1 - w) \sum_{j \in \mathcal{M}_{i}^{1}(c)} \sum_{k \in \mathcal{N}_{j}^{1}(c)} \frac{(R_{d} - \frac{\lambda_{ij} + \lambda_{jk}}{\lambda_{ij}\lambda_{jk}})V}{T}.$$

-The Content Exchange Protocol

 Exchange control message, including subscribed channel list and collected buffer state of each other.



• Generate the candidate request list.

Node
$$i$$

$$L_{i} = S_{i} - (S_{i} \cap S_{j})$$

$$L_{i} = S_{j} - (S_{i} \cap S_{j})$$

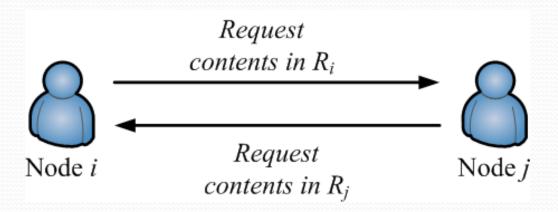
$$L_{j} = S_{j} - (S_{i} \cap S_{j})$$

$$R_{j} = L_{j} + L_{j}$$

$$R_{j} = L_{j} + L_{j}$$

-The Content Exchange Protocol

 Exchange content in the decreasing order of priority and content utility one by one.

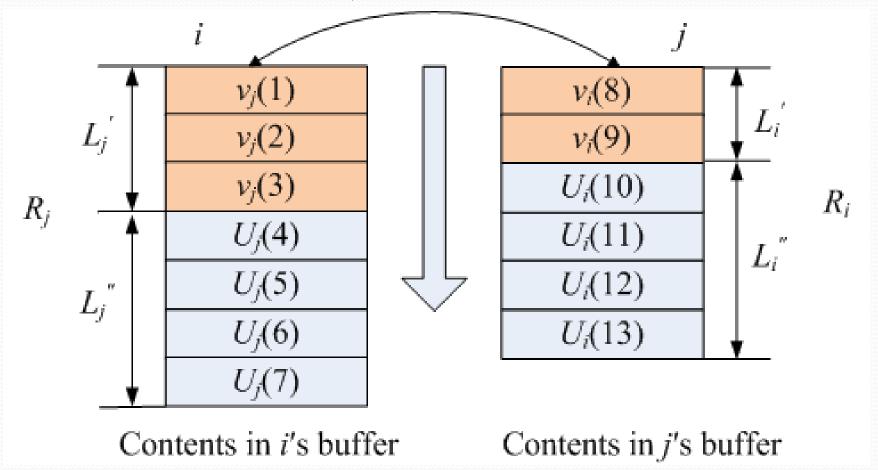


 Finish the trading until one side does not have contents, which can increase the total content utility in the buffer for the other side.

-The Content Exchange Protocol



• An example about the content exchange process between nodes *i* and *j* in ConDis.



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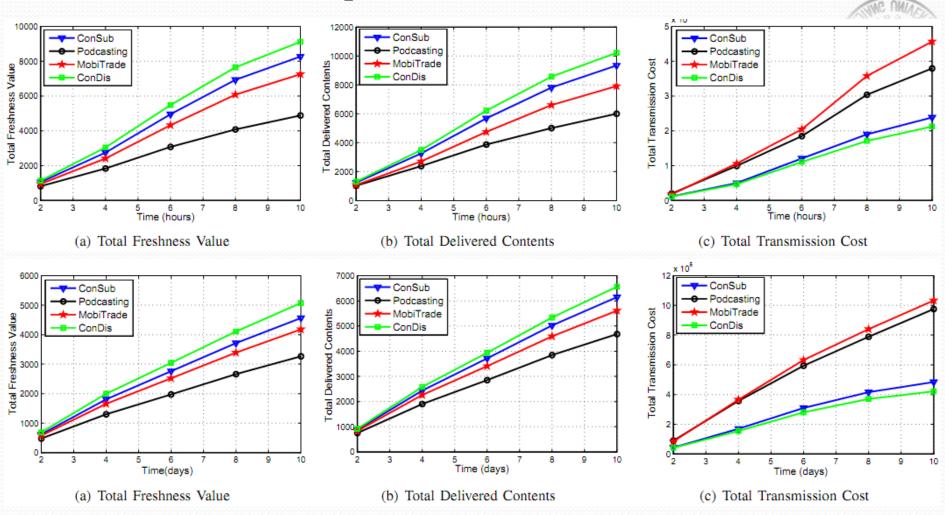
Performance Evaluation

- We compare our work with Podcasting [17], MobiTrade
 [9], ConSub [8].
- Three Performance Metrics: Total freshness value, Total delivered contents and Total transmission Cost.
- Two experimental traces: Infocom o6 and MIT Reality.
- There are 5 channels in the network, and each node only expresses interest, randomly, in one channel.

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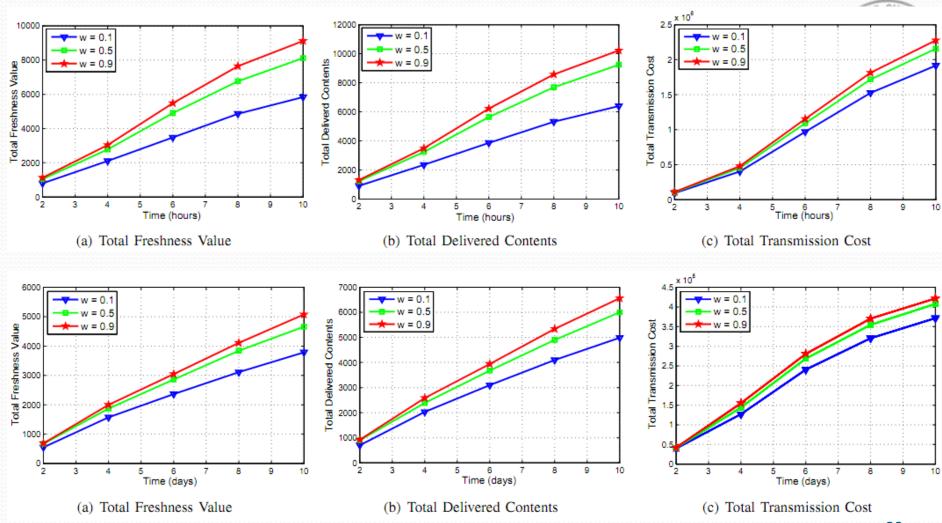
• Performance comparison



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Conclusion

- We propose an incentive-driven and freshness-aware pub/sub content dissemination scheme, called ConDis, for selfish OppNets.
- we also propose a novel content exchange protocol when nodes are in contact.
- Extensive realistic trace-driven simulations are conducted to evaluate the performance of our proposed scheme



Thanks! Questions?