Chapter I Introduction

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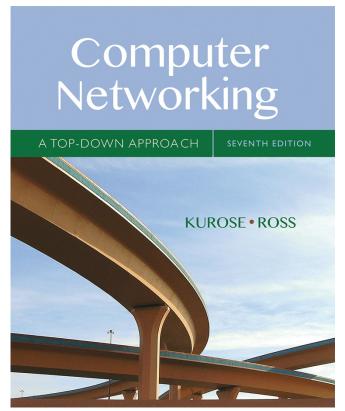
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Computer Networking: A Top Down Approach

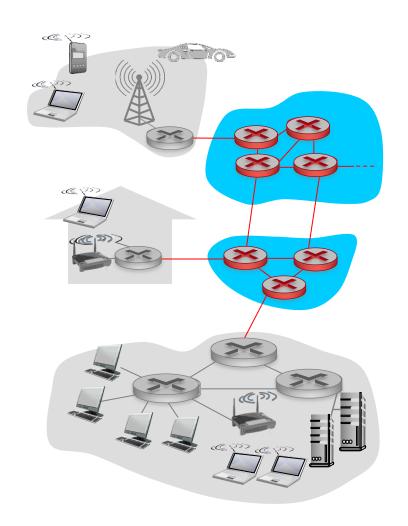
7th edition
Jim Kurose, Keith Ross
Pearson/Addison Wesley
April 2016

Chapter I: roadmap

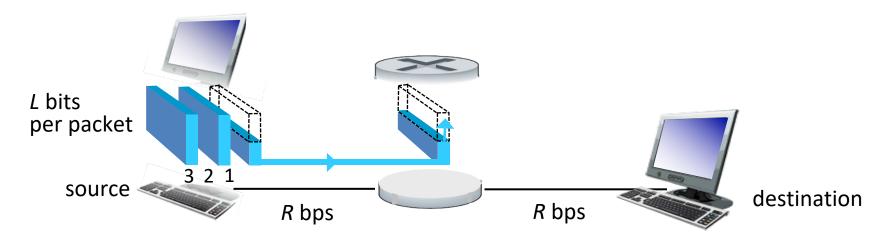
- I.I what is the Internet?
- 1.2 network edge
 - end systems, access networks, links
- 1.3 network core
 - packet switching, circuit switching, network structure
- 1.4 delay, loss, throughput in networks
- 1.5 protocol layers, service models
- 1.6 networks under attack: security
- 1.7 history

The network core

- mesh of interconnected routers
- packet-switching: hosts break application-layer messages into packets
 - forward packets from one router to the next, across links on path from source to destination
 - each packet transmitted at full link capacity



Packet-switching: store-and-forward



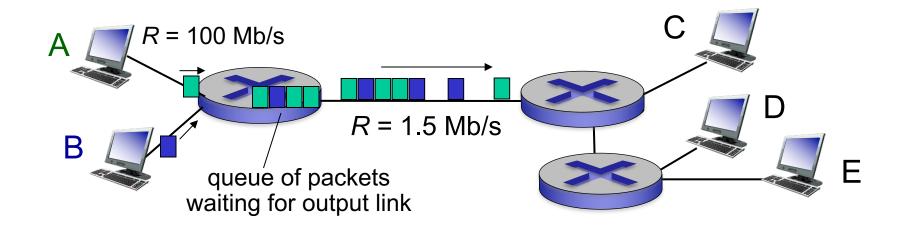
- takes L/R seconds to transmit (push out) L-bit packet into link at R bps
- store and forward: entire packet must arrive at router before it can be transmitted on next link
- end-end delay = 2L/R (assuming zero propagation delay)

one-hop numerical example:

- L = 7.5 Mbits
- R = 1.5 Mbps
- one-hop transmission delay = 5 sec

more on delay shortly ...

Packet Switching: queueing delay, loss



queuing and loss:

- if arrival rate (in bits) to link exceeds transmission rate of link for a period of time:
 - packets will queue, wait to be transmitted on link
 - packets can be dropped (lost) if memory (buffer) fills up

Two key network-core functions

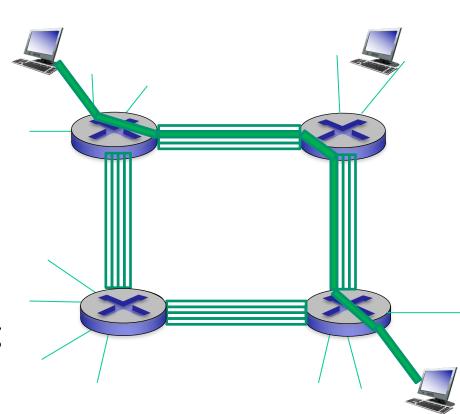
routing: determines sourcedestination route taken by forwarding: move packets from packets router's input to appropriate routing algorithms router output routing algorithm local forwarding table header value output link 0100 3 0101 0111 1001 destination address in arriving packet's header

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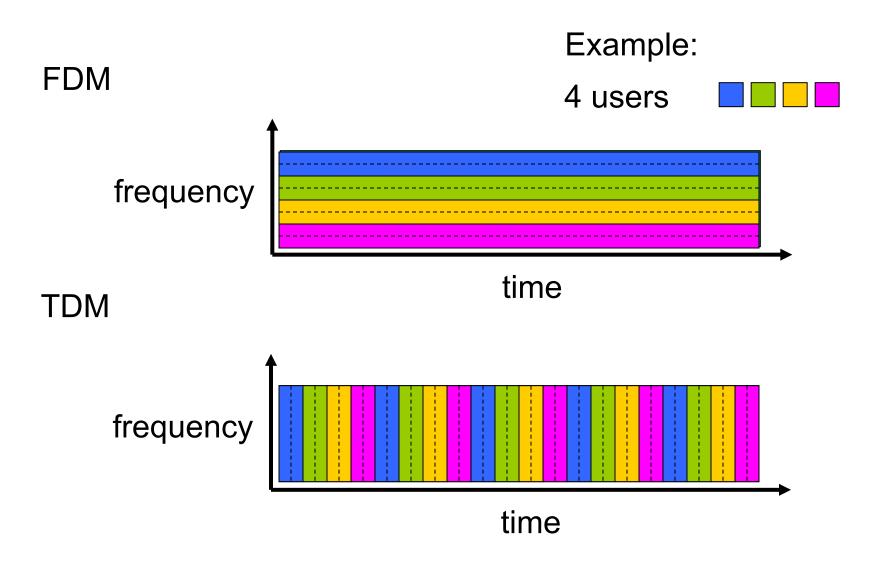
Alternative core: circuit switching

end-end resources allocated to, reserved for "call" between source & dest:

- in diagram, each link has four circuits.
 - call gets 2nd circuit in top link and Ist circuit in right link.
- dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- circuit segment idle if not used by call (no sharing)
- commonly used in traditional telephone networks



Circuit switching: FDM versus TDM



Packet switching versus circuit switching

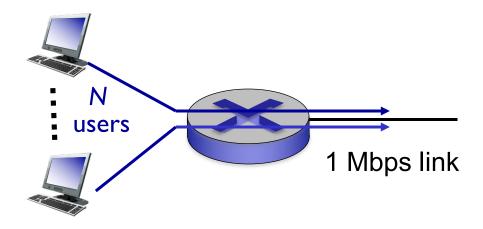
packet switching allows more users to use network!

example:

- I Mb/s link
- each user:
 - 100 kb/s when "active"
 - active 10% of time



- 10 users
- packet switching:
 - with 35 users, probability > 10 active at same time is less than .0004 *



Q: how did we get value 0.0004?

Q: what happens if > 35 users?

^{*} Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive/

Packet switching versus circuit switching

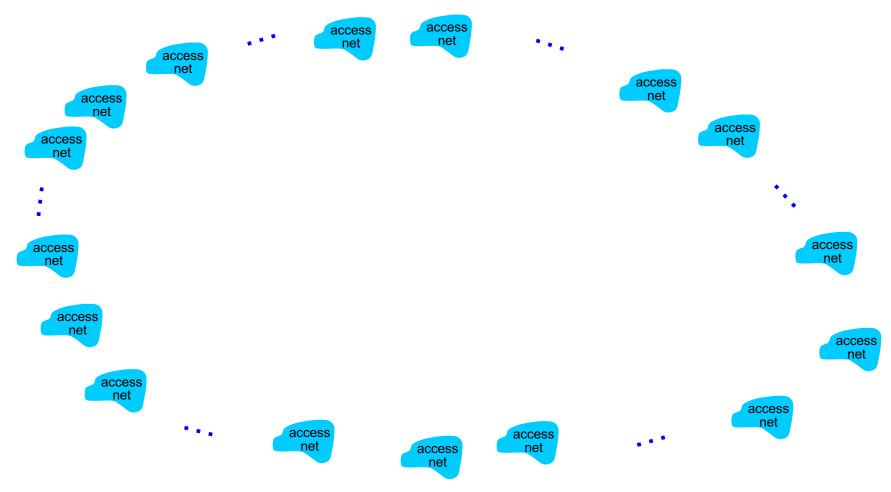
is packet switching a "slam dunk winner?"

- great for bursty data
 - resource sharing
 - simpler, no call setup
- excessive congestion possible: packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem (chapter 7)

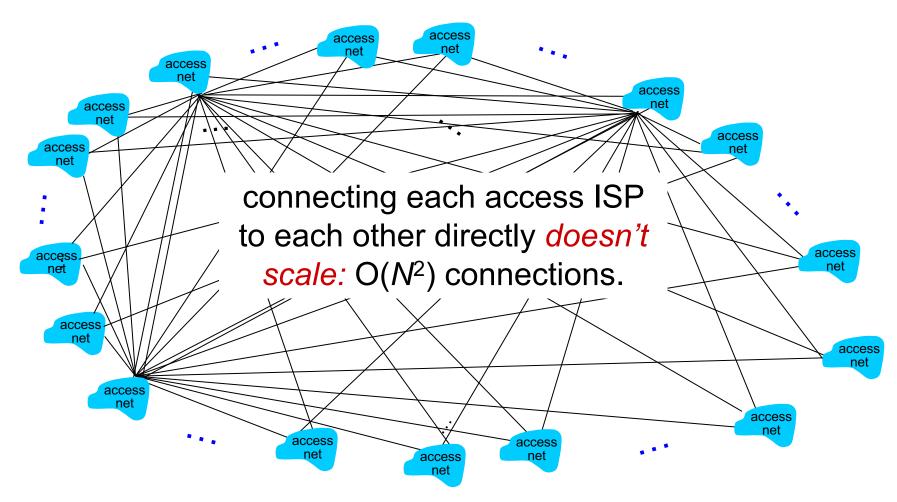
Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet-switching)?

- End systems connect to Internet via access ISPs (Internet Service Providers)
 - residential, company and university ISPs
- Access ISPs in turn must be interconnected.
 - so that any two hosts can send packets to each other
- Resulting network of networks is very complex
 - evolution was driven by economics and national policies
- Let's take a stepwise approach to describe current Internet structure

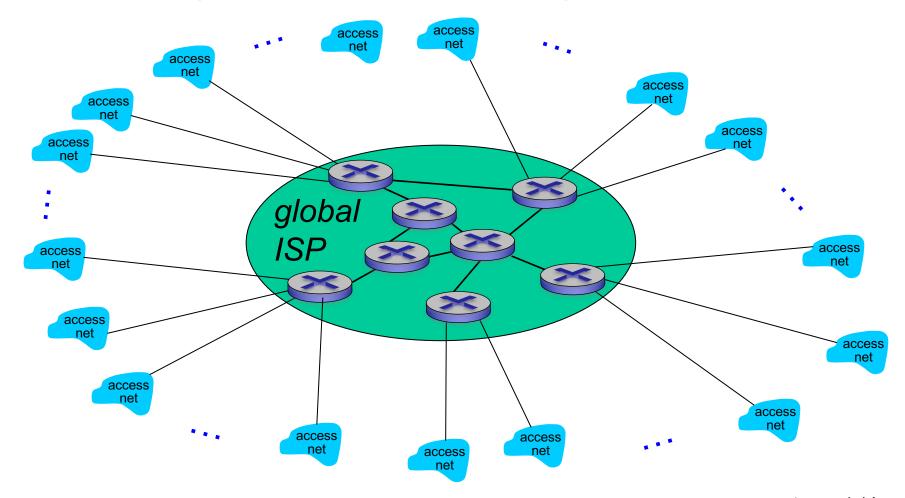
Question: given millions of access ISPs, how to connect them together?



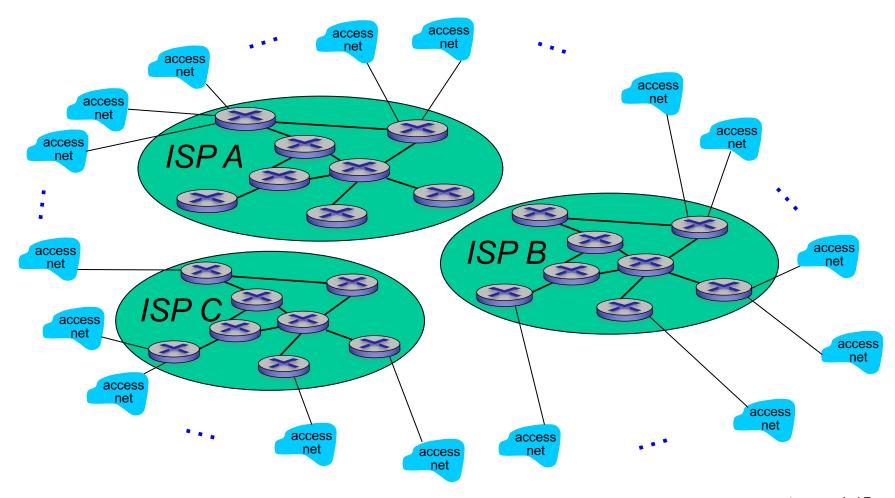
Option: connect each access ISP to every other access ISP?



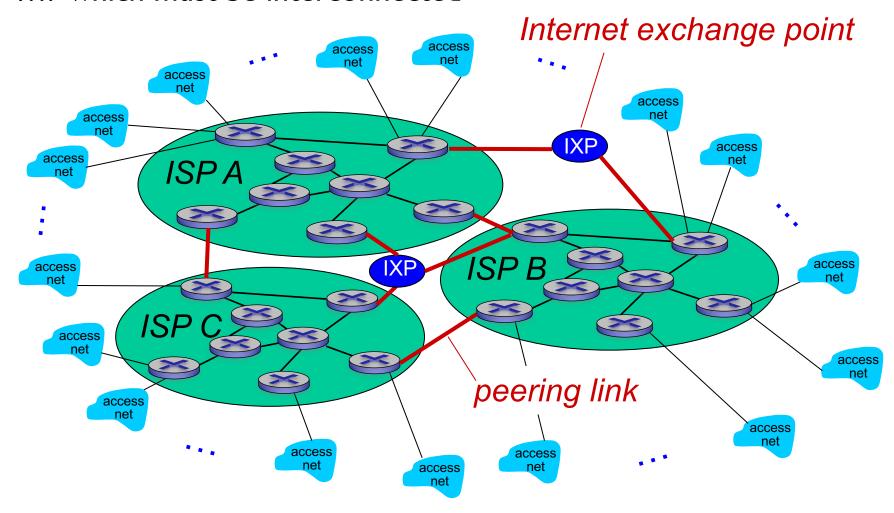
Option: connect each access ISP to one global transit ISP? Customer and provider ISPs have economic agreement.



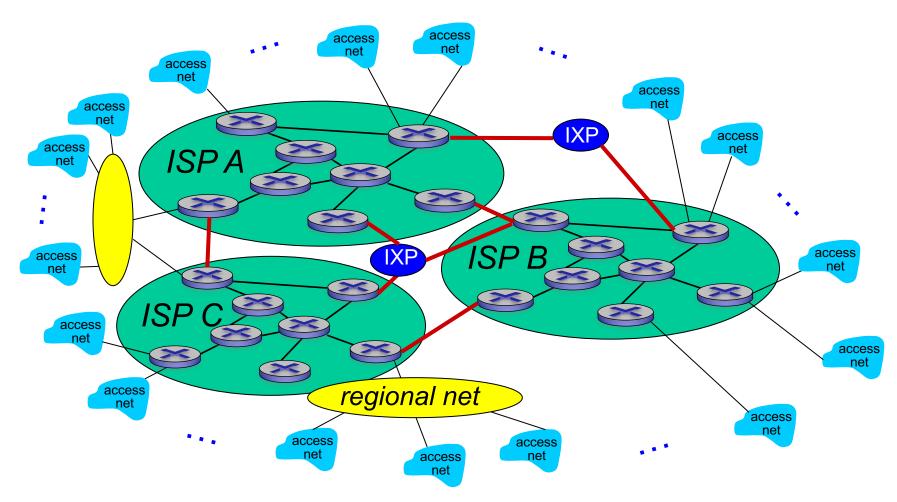
But if one global ISP is viable business, there will be competitors



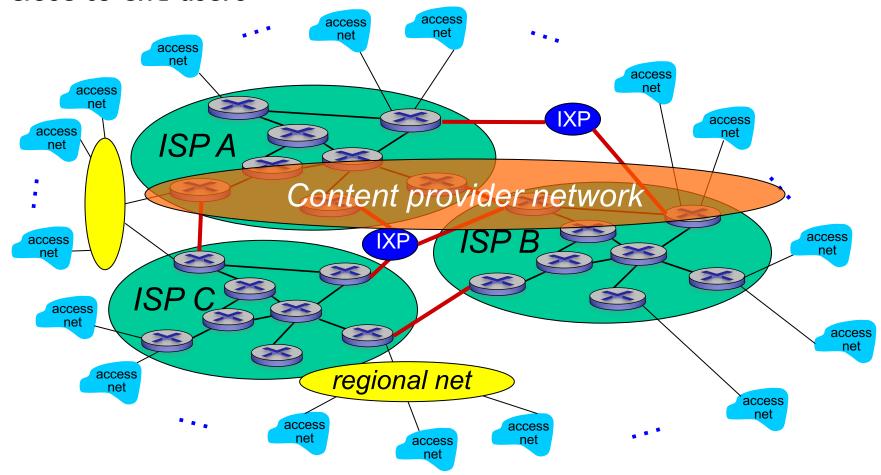
But if one global ISP is viable business, there will be competitors which must be interconnected

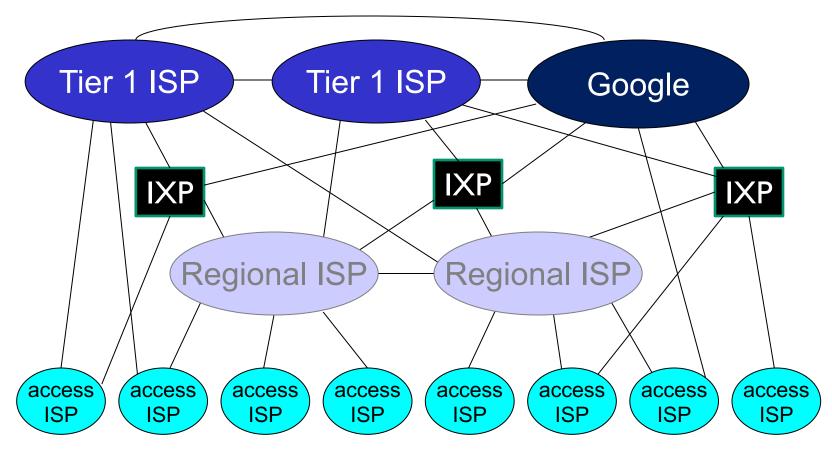


... and regional networks may arise to connect access nets to ISPs



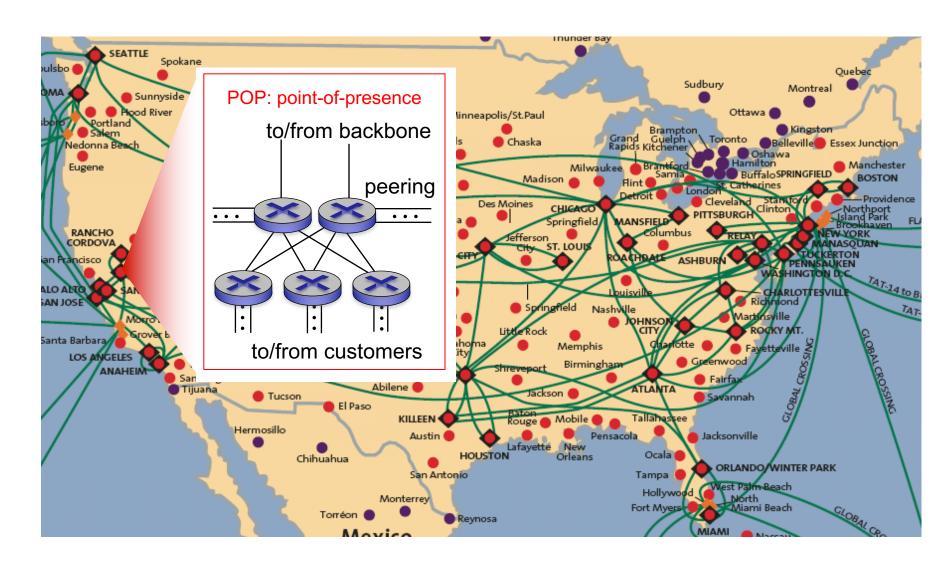
... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users





- at center: small # of well-connected large networks
 - "tier-I" commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
 - content provider network (e.g., Google): private network that connects it data centers to Internet, often bypassing tier-I, regional ISPs Introduction 1-19

Tier-I ISP: e.g., Sprint



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