

lecture 23:

background on symbolic execution

5590: software defined networking

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TTLMAN 401B, R 17:30-20:00

symbolic execution and program testing

<http://dl.acm.org/citation.cfm?id=360252>

the symbolic execution approach

the symbolic execution approach

towards large production of reliable programs

- two extreme alternatives

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program testing

- best effort: run program on sample inputs
- sample inputs?

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program verification

- logical proof based on precise specification
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symbolic execution: a compromise

- assures “program meets its requirement” even when formal specifications are not given

symbolic execution



symbolic execution

normal
inputs



program execution



values

symbols
(class of inputs)




symbolic execution



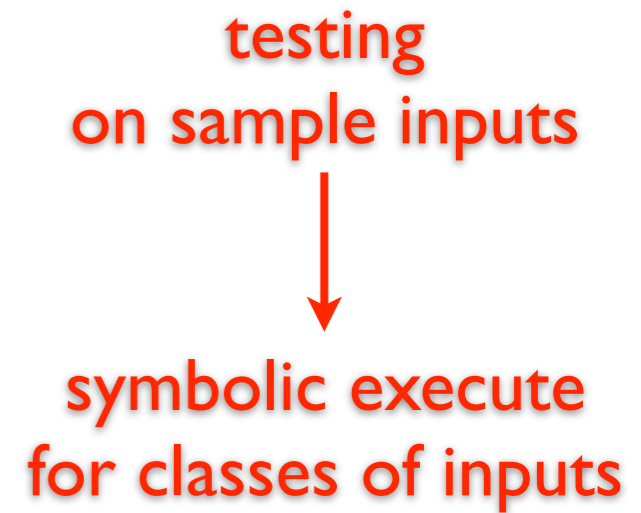
symbolic formulas
over input symbols

testing
on sample inputs



symbolic execute
for classes of inputs

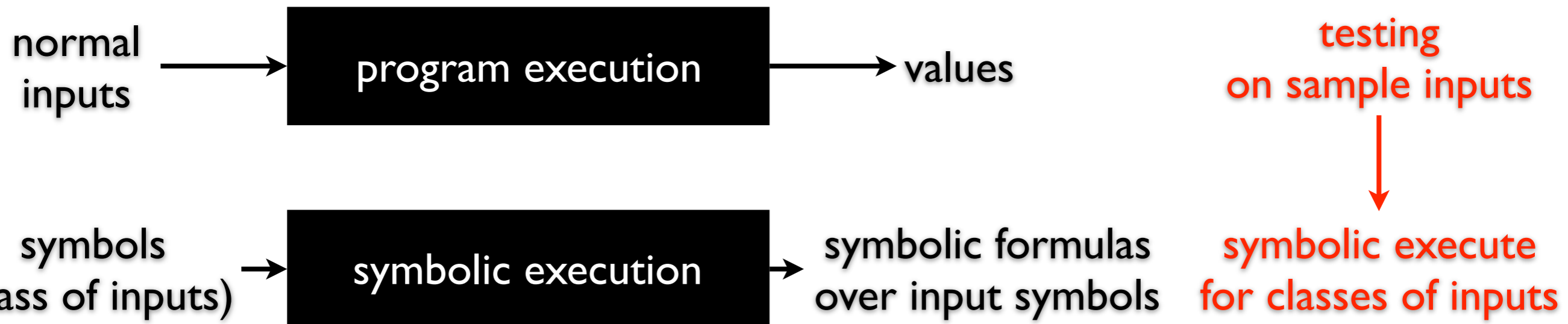
symbolic execution



class of inputs

- characterized by each symbol execution

symbolic execution



class of inputs

- characterized by each symbol execution

coverage

- determined by the dependence of program's control
 - e.g., control flow independent of inputs — a single symbolic execution

execution semantics

for normal program execution

- data objects
- how statements manipulate data objects
- how control flows through the statements

state of a program execution

- values of program variables
- statement counter

execution semantics

for symbolic execution

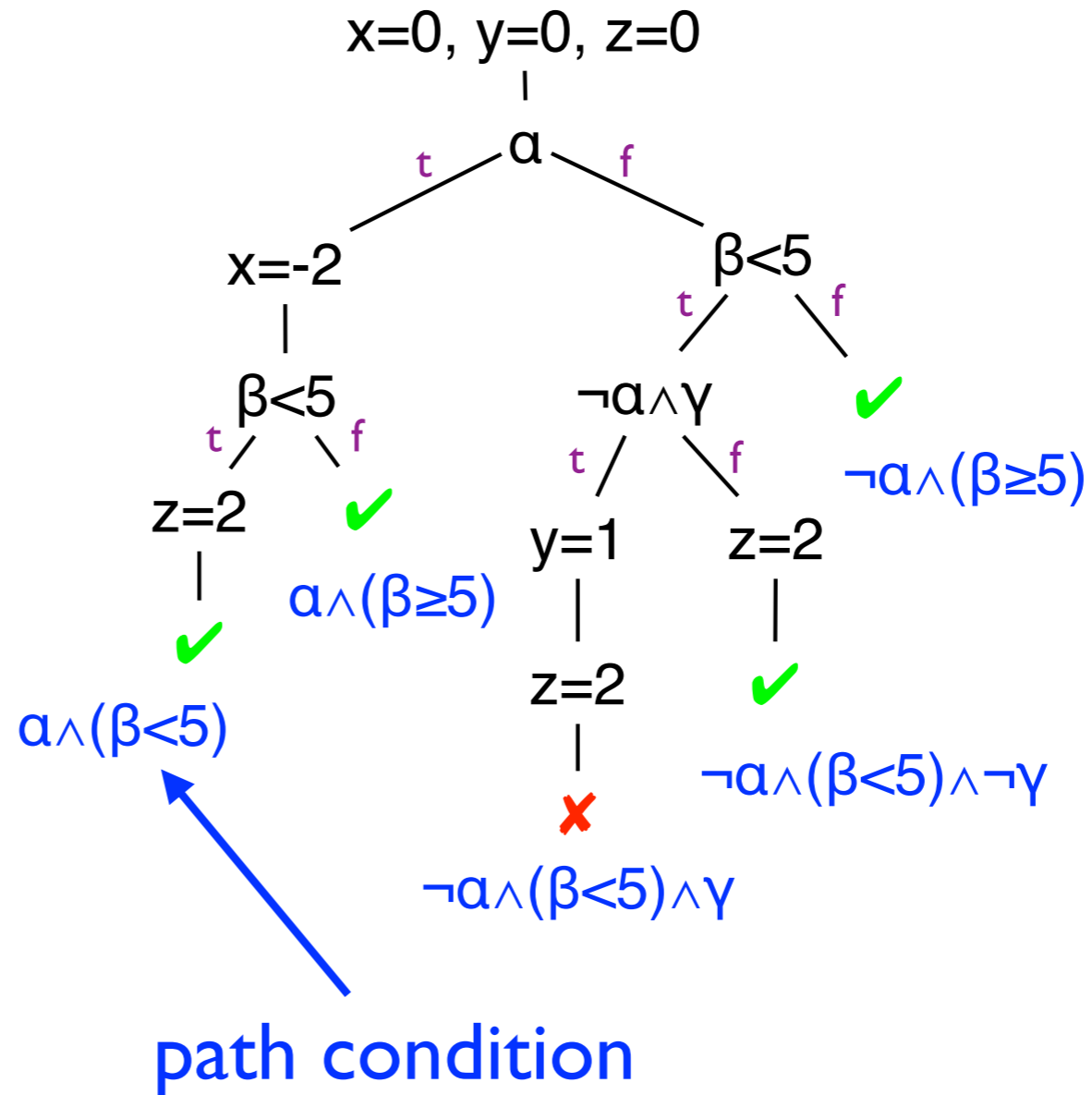
- symbolic data objects
- evaluation rules

state

- path condition (pc)
 - accumulator of properties which the inputs must satisfy for an execution to follow the path

symbolic execution tree

```
1. int a = a, b = β, c = γ;
2.           // symbolic
3. int x = 0, y = 0, z = 0;
4. if (a) {
5.   x = -2;
6. }
7. if (b < 5) {
8.   if (!a && c) { y = 1; }
9.   z = 2;
10.}
11. assert(x+y+z!=3)
```



IF-statements

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IF $q(\text{inputs} \dots)$

- in choosing between alternative paths, assumptions about the inputs are made and added (conjoined) to pc
- pc initialized to **TRUE**

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non-forking

- if $pc \supset q$ or $pc \supset \neg q$

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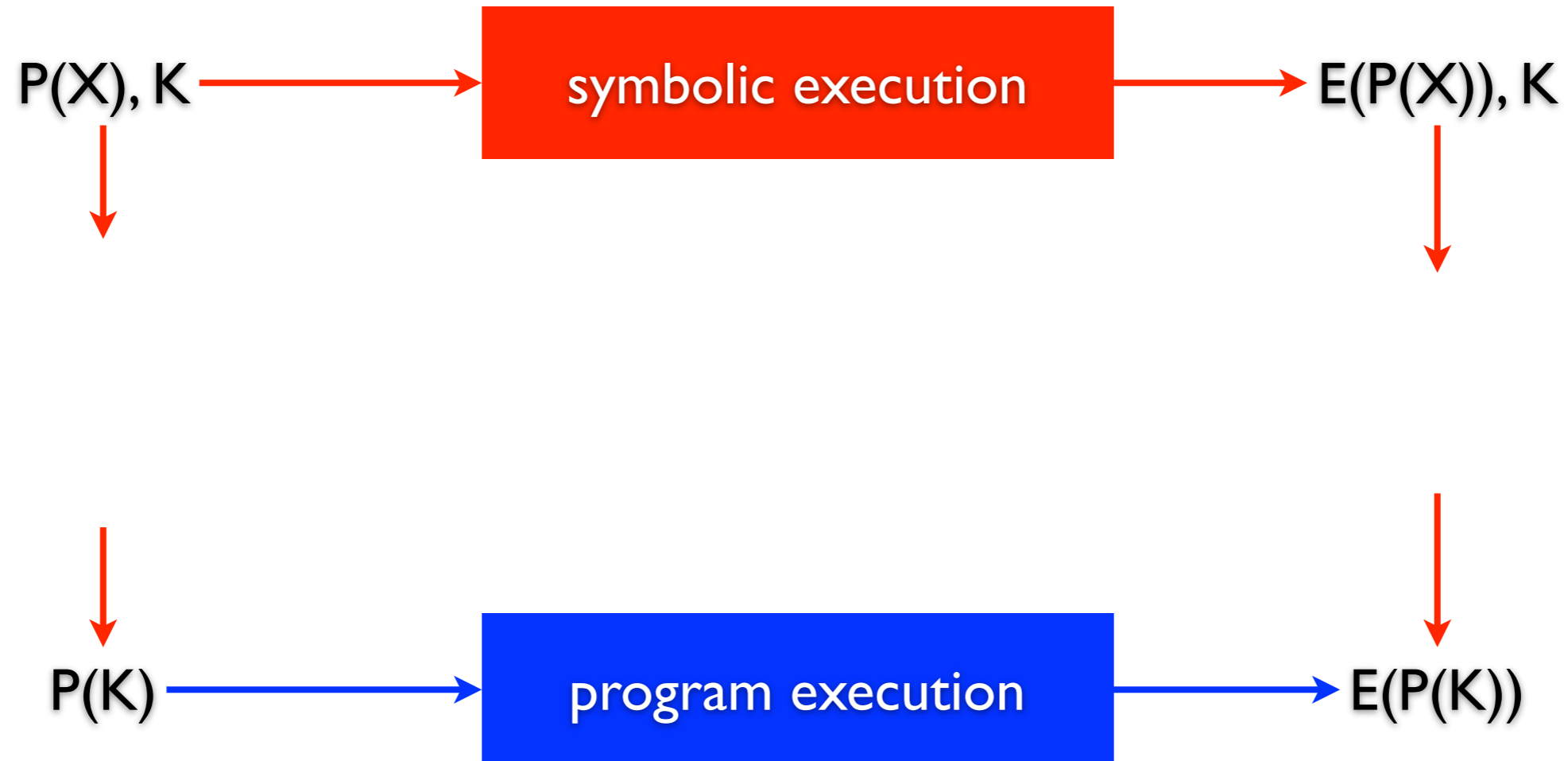
- if $pc \supset q$ or $pc \supset \neg q$

forking

- if neither $pc \supset q$ or $pc \supset \neg q$ is true
- THEN branch: $pc \leftarrow pc \wedge q$
- ELSE branch: $pc \leftarrow pc \wedge \neg q$

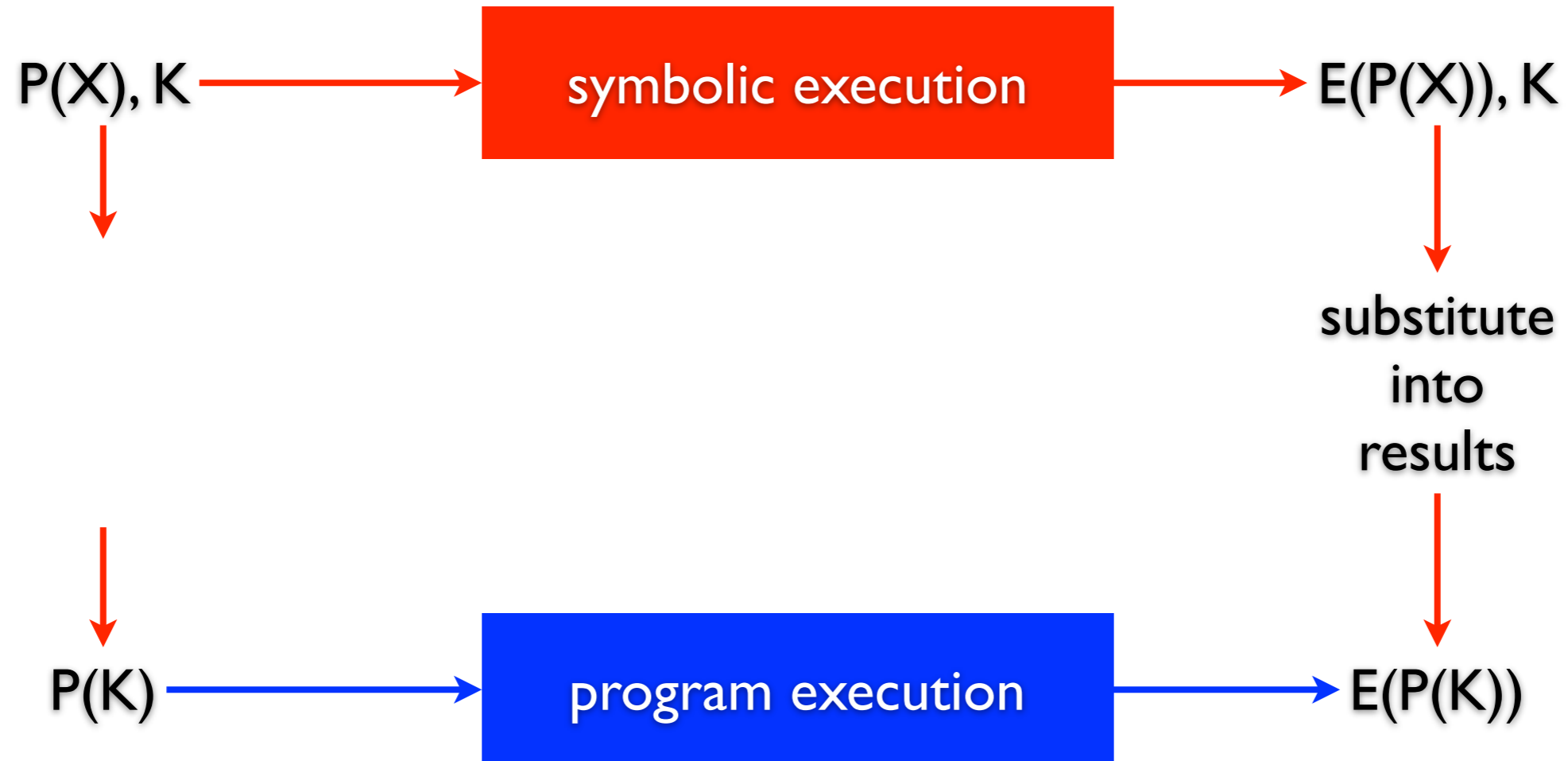
commutativity

P — program, $E(P(X))$ — execute P on input symbol X , K — concrete inputs



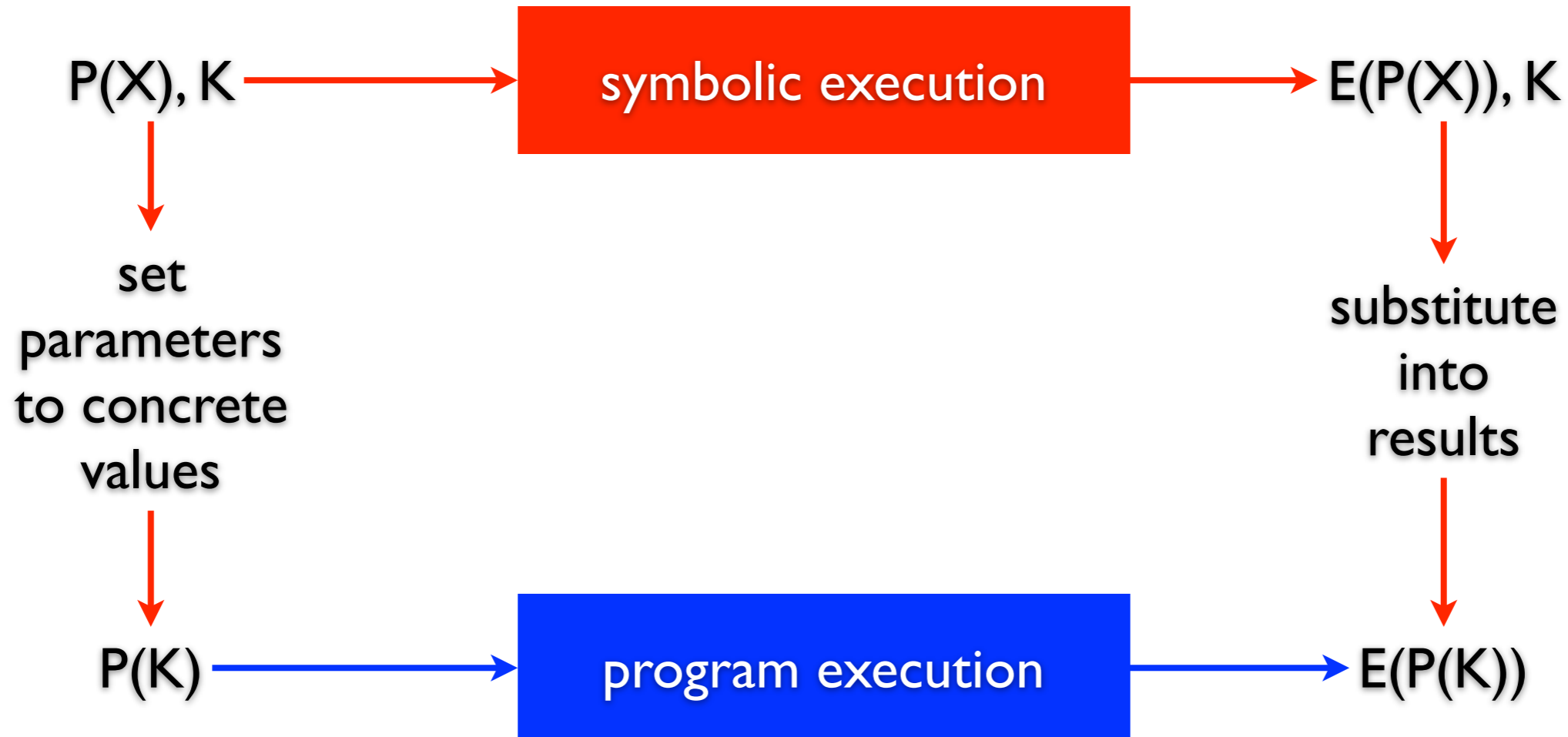
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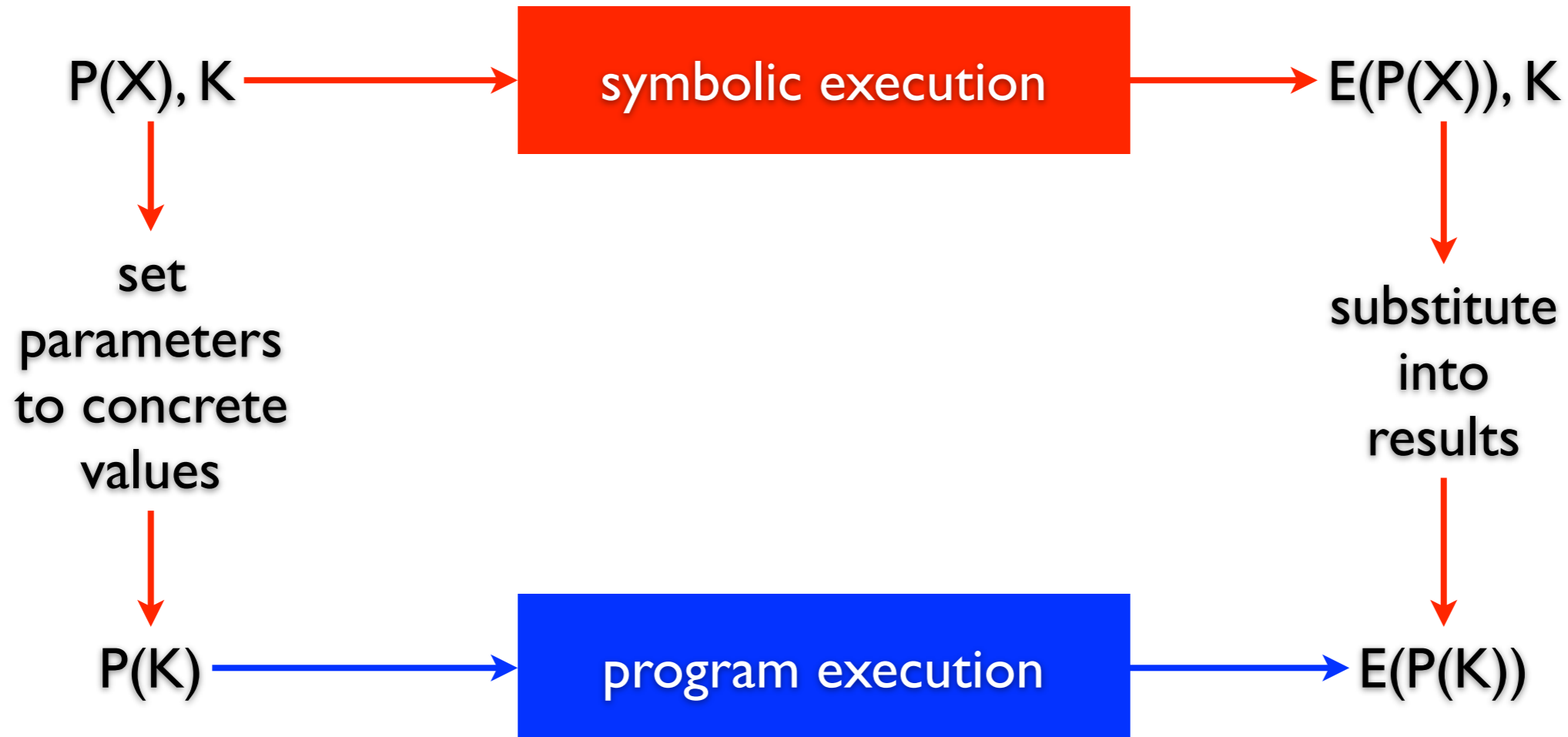
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- symbolic execution captures the same effect as conventional execution
- the specific computation of the program are generalized and delayed

the challenges

the interesting case — brach type statements

- executing “IF statement” requires theorem proving

still an enhanced testing methodology

- input classes needed (to exhaust all possible cases) is practically infinite
 - lots of program paths
 - program state has many bits

industrial strength tool

- Klee

final report

formatting requirements

no more than two (2) single-spaced page

SIGCOMM style

- 10pt font on 12pt leading formatted for printing on Letter-sized (8.5" by 11")

latex template

- <http://conferences.sigcomm.org/sigcomm/2017/misc/sig-alternate-10pt.cls>

what to present

motivation

- introduce the problem
- if a concrete problem, place it in a larger problem space

introduction, approach

- present your solution — the method, design, architecture
 - describe your solution and the rational behind it

status

- preliminary results? what are the next steps?
- challenges and discussions?

examples

- <https://www.usenix.org/conference/ons2014/technical-sessions/>

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