

CACHE CONTENT PLACEMENT USING TRIANGULAR NETWORK CODING

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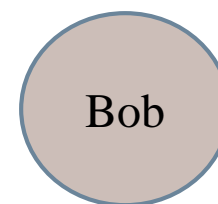
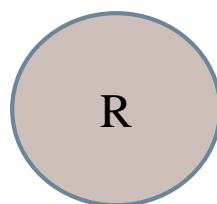
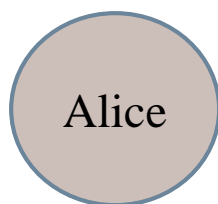
Agenda

2

- Introduction
- Motivation
- Content placement algorithm
- Simulation
- Conclusion

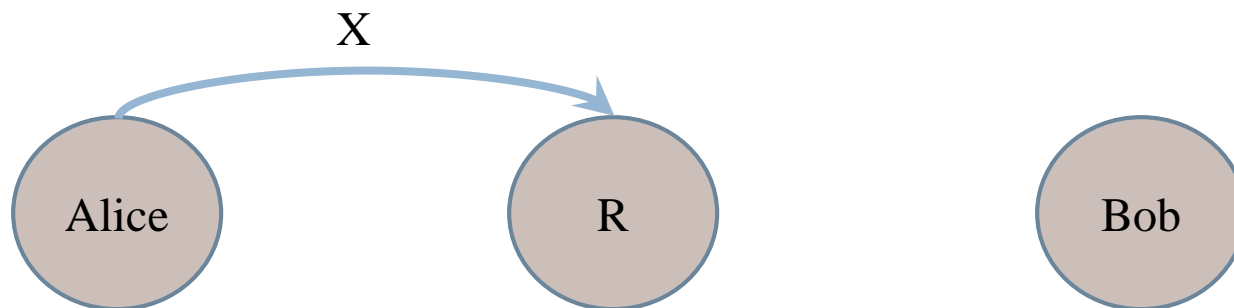
Alice and Bob (No coding)

3



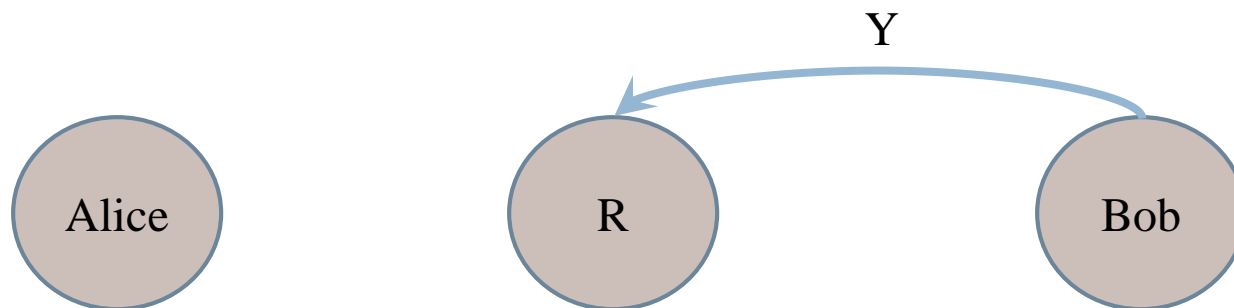
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3



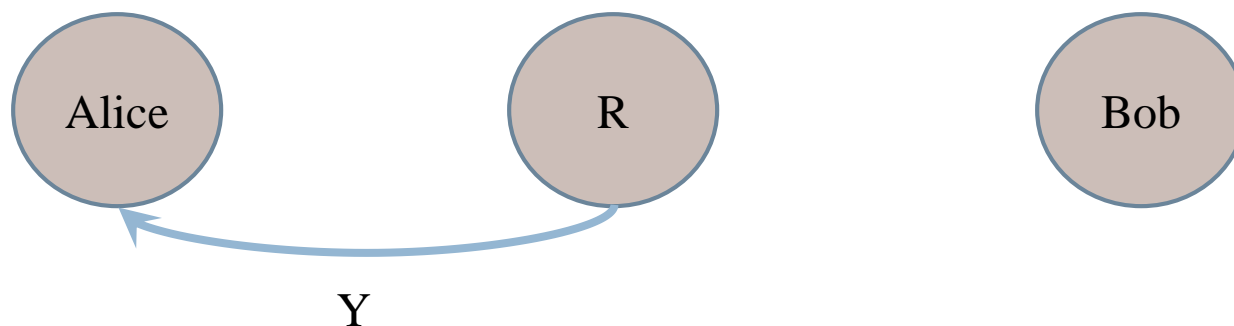
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3



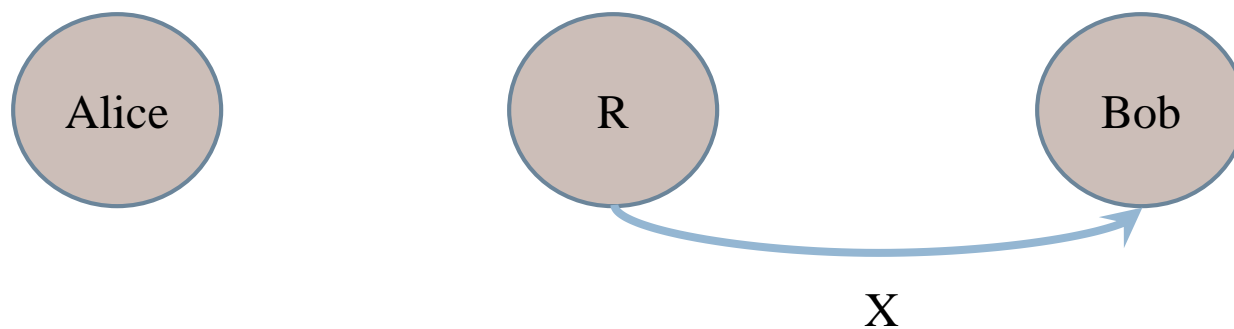
Alice and Bob (No coding)

3



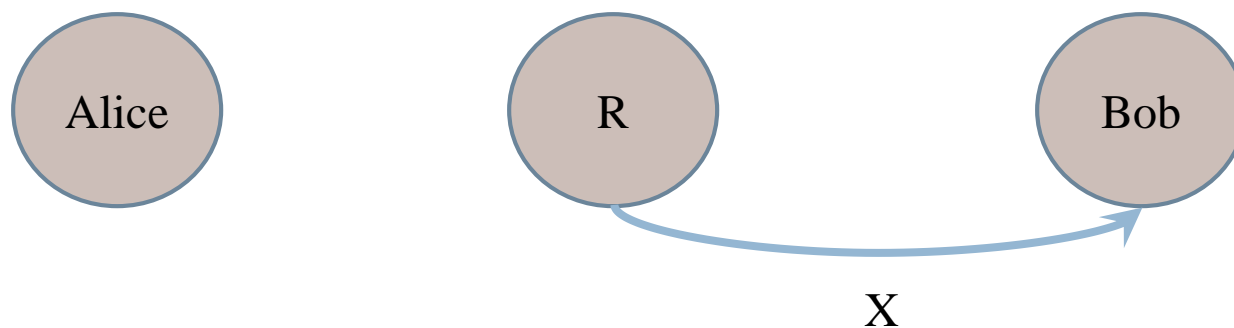
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Alice and Bob (No coding)

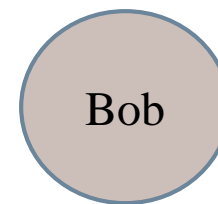
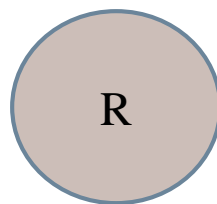
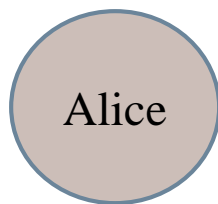
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4 transmissions

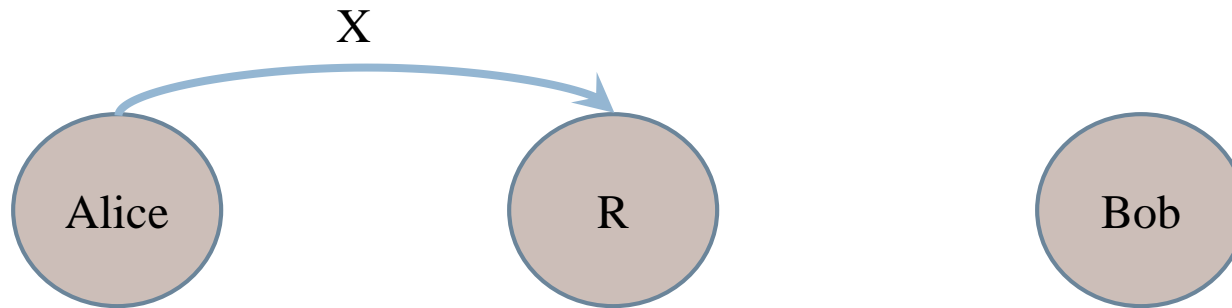
Alice and Bob (Coding)

4



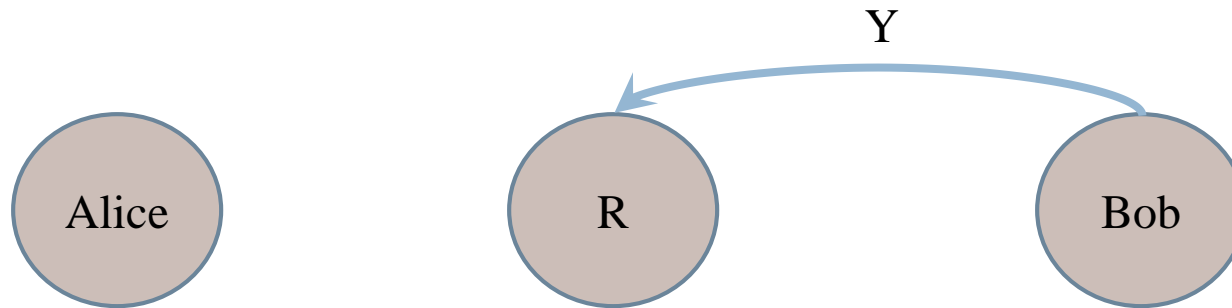
Alice and Bob (Coding)

4



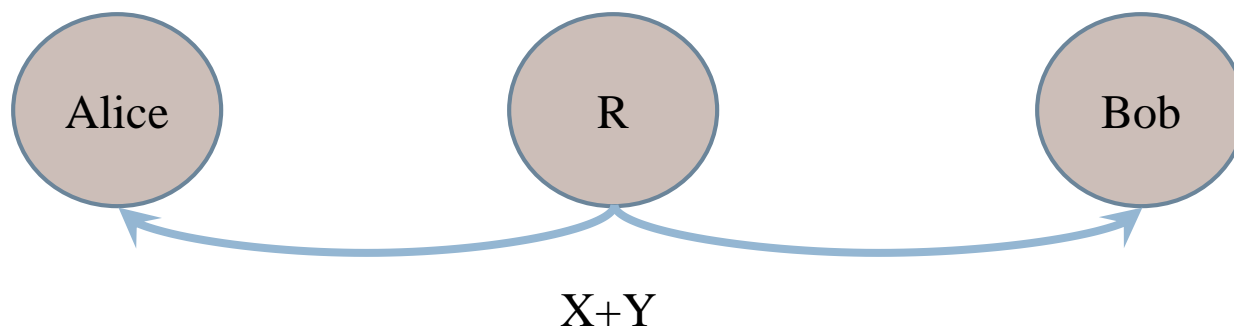
Alice and Bob (Coding)

4



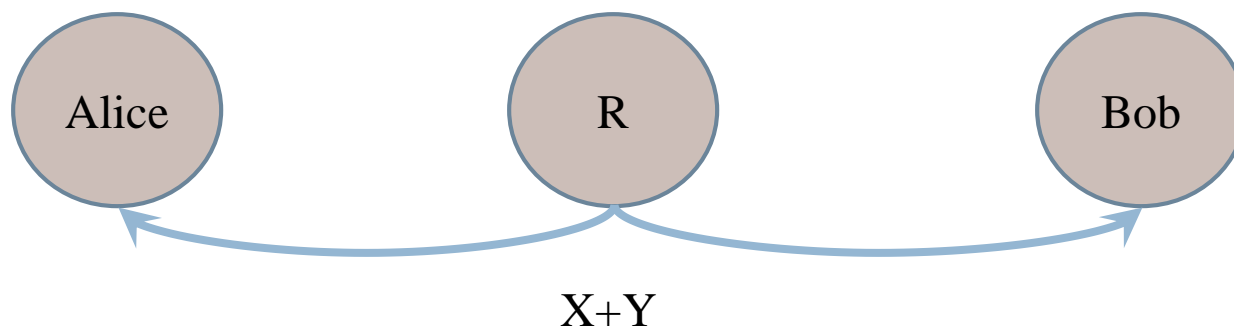
Alice and Bob (Coding)

4



Alice and Bob (Coding)

4

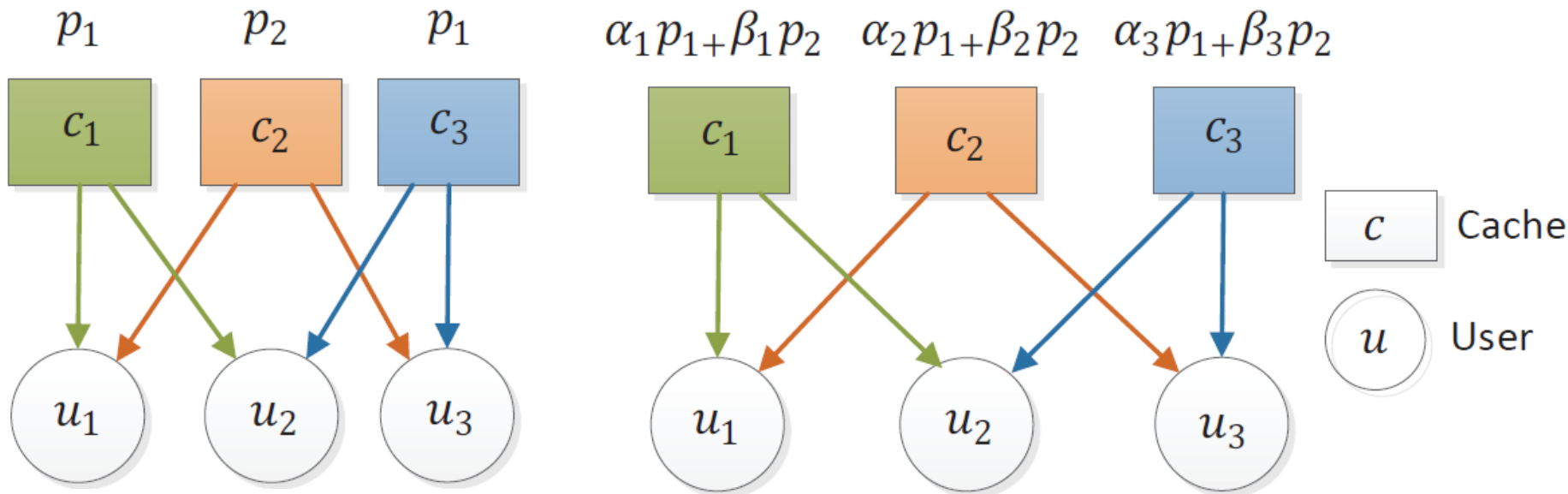


3 transmissions

Motivation

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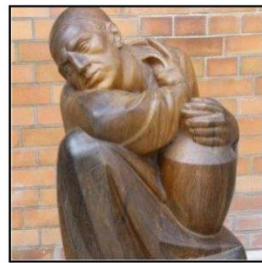
- Providing more amount of data to the users.



Setting

6

- h video layers on the server: p_1, \dots, p_h
- Layer p_i is not useful without the layers with a smaller index.



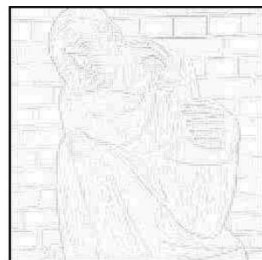
(a) Original



(b) Layer 1



(c) Layer 2



(d) Layer 3



(e) Layers 1 & 2



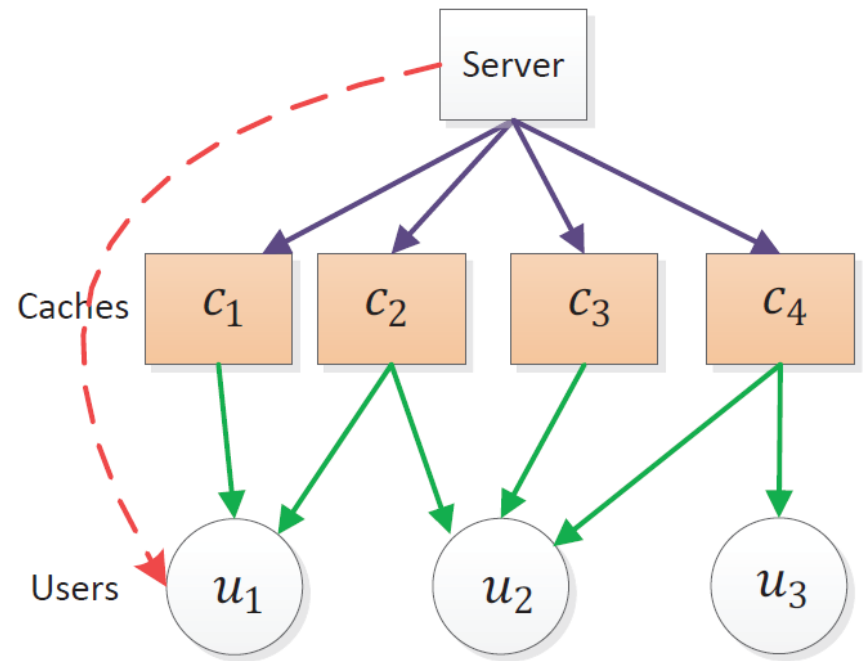
(f) Layers 2 & 3

Setting

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- Capacity=size of the video layers
- Objective: maximizing the total number of available layers.

$$\max \sum_{i=1}^h \sum_{j=1}^n z_{ij}$$



Triangular Coding

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□ Linear Coding

- ▣ $2^h - 1$ ways to code h layers.
- ▣ $(2^h - 1)^n$ different possible placements for n caches.

□ Triangular network coding

- ▣ The encoded video layers are in the form $\sum_{j=1}^k \alpha_j p_j$.

Original packets

$$\begin{cases} p_1 \\ p_2 \\ p_3 \end{cases}$$

Linear coding

$$\begin{cases} p_1, p_2, p_3 \\ p_1 + p_2, p_1 + p_3, p_2 + p_3 \\ p_1 + p_2 + p_3 \end{cases}$$

Triangular coding

$$\begin{cases} p_1 \\ p_1 + p_2 \\ p_1 + p_2 + p_3 \end{cases}$$

Content Placement Algorithm

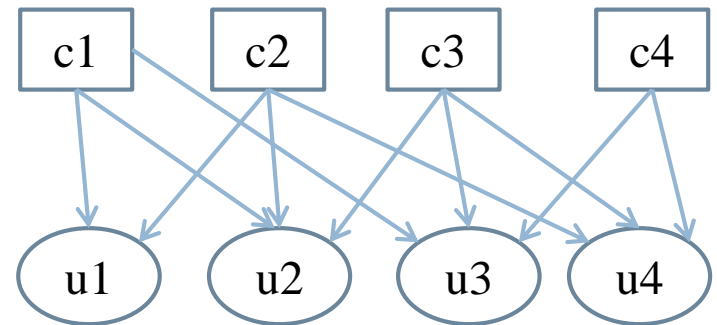
9

- The problem of efficient content placement on the caches is an NP-complete problem.
- The greedy algorithm fills-up the caches in rounds.
- In each round, we select a user and fill-up its adjacent caches.
- Selection rules
 - ▣ **Rule 1:** the user with the minimum degree.
 - ▣ **Rule 2:** the user with a larger number of filled-up caches.
 - ▣ **Rule 3:** the user whose adjacent caches have less cumulative ranks.
- The algorithm fills-up the empty adjacent caches to user u_i with a random linear combination of the first $d_i - v_i + r_i$ video layers.

Example

10

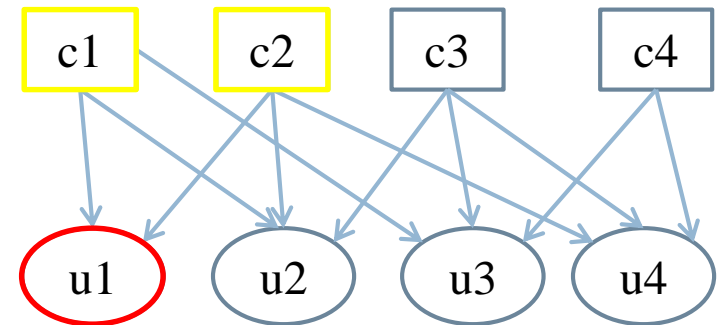
- Step 1: user u_1 has the minimum degree.
 - ▣ $2-0+0=2$
- Step 2: user u_2 has 2 filled adjacent caches.
 - ▣ $3-2+2=3$
- Step 3: select u_3 or u_4 randomly (assume u_3).
 - ▣ $3-2+2=3$



Example

10

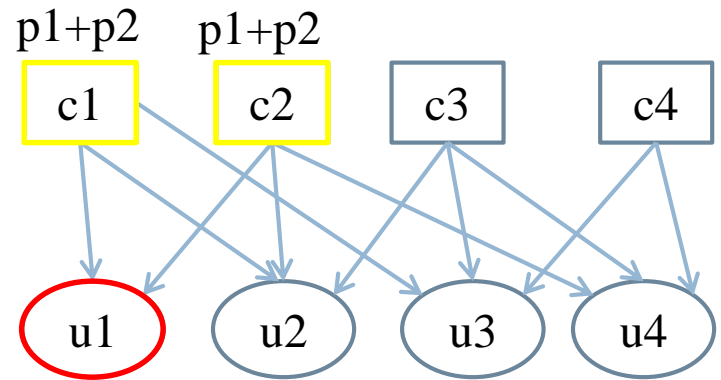
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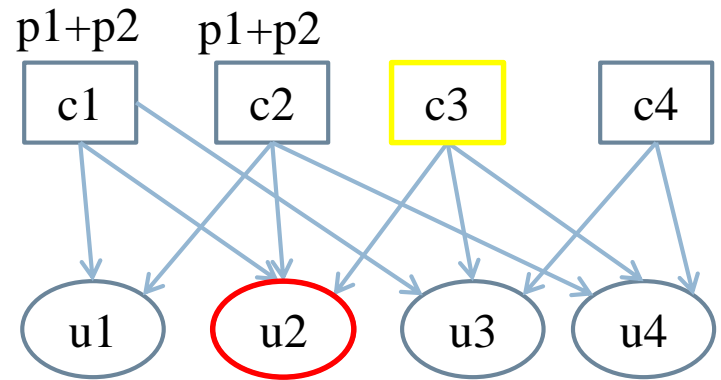
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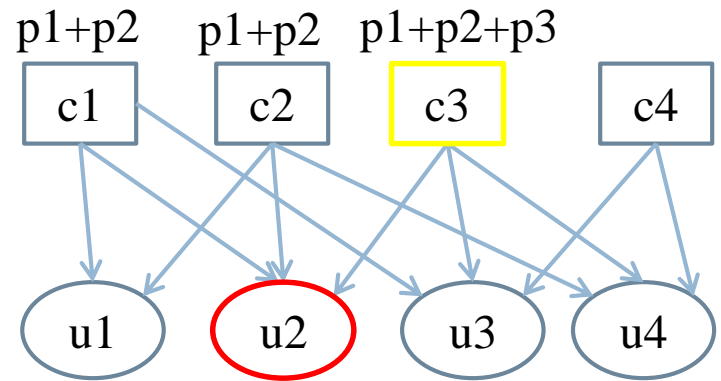
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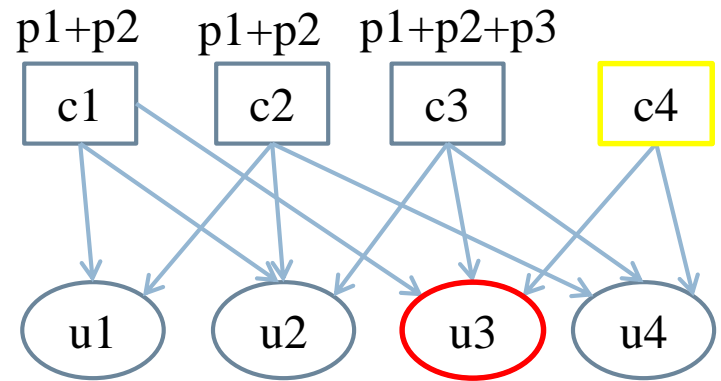
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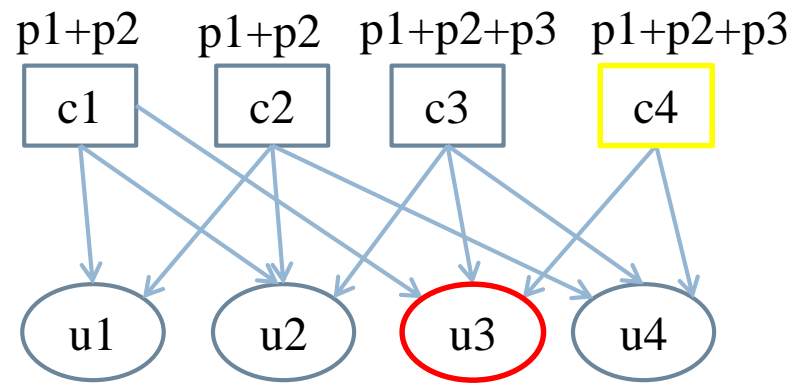
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Simulation Setting

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- Simulator in the MATLAB environment.
- Comparison
 - ▣ Number of available layers to the users.
 - ▣ Average utility: the number of available layers to a user divided by its degree.
 - ▣ Fairness: we define unfairness as the average difference between the number of available layers to each user and the average number of available layers to the users.

$$f' = \frac{\sum_{i=1}^m |q_i - e|}{m}$$

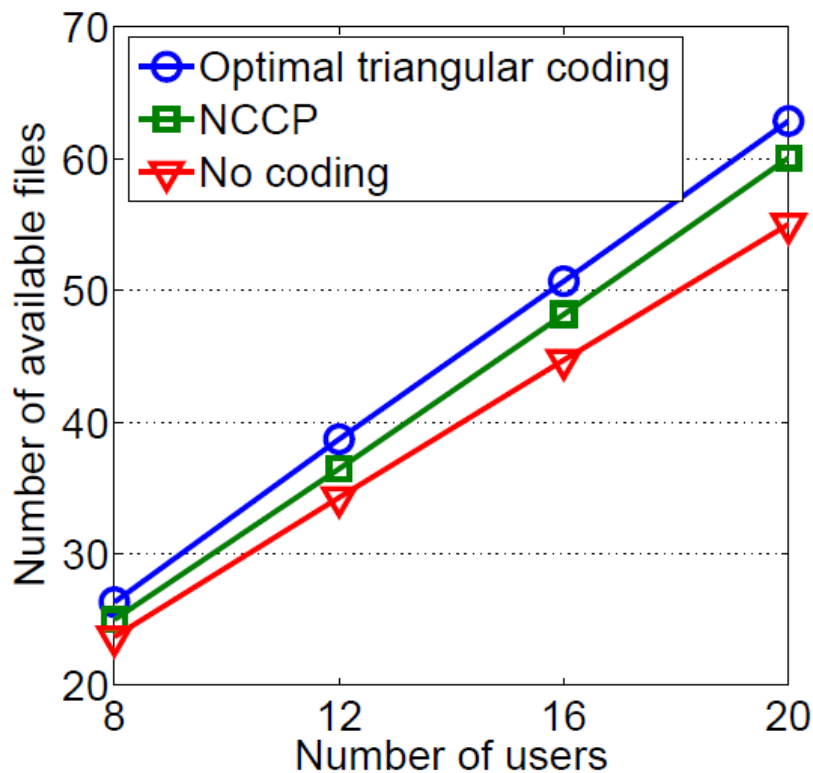
$$e = \frac{\sum_{i=1}^m q_i}{m}$$

$$f = \frac{1}{f'}$$

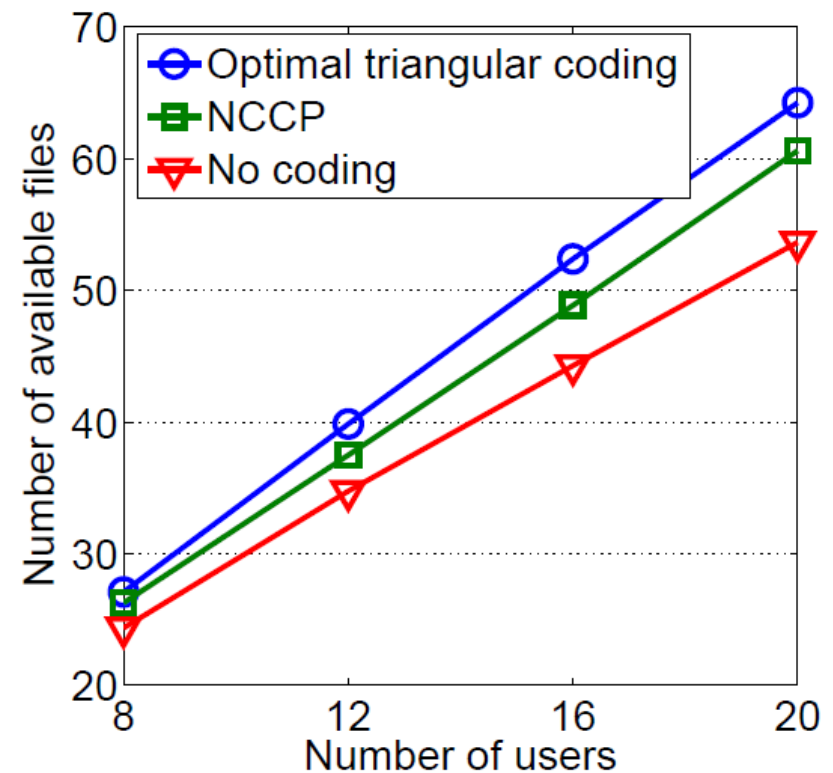
Simulations

12

- Number of caches: 5
- Number of layers: 4



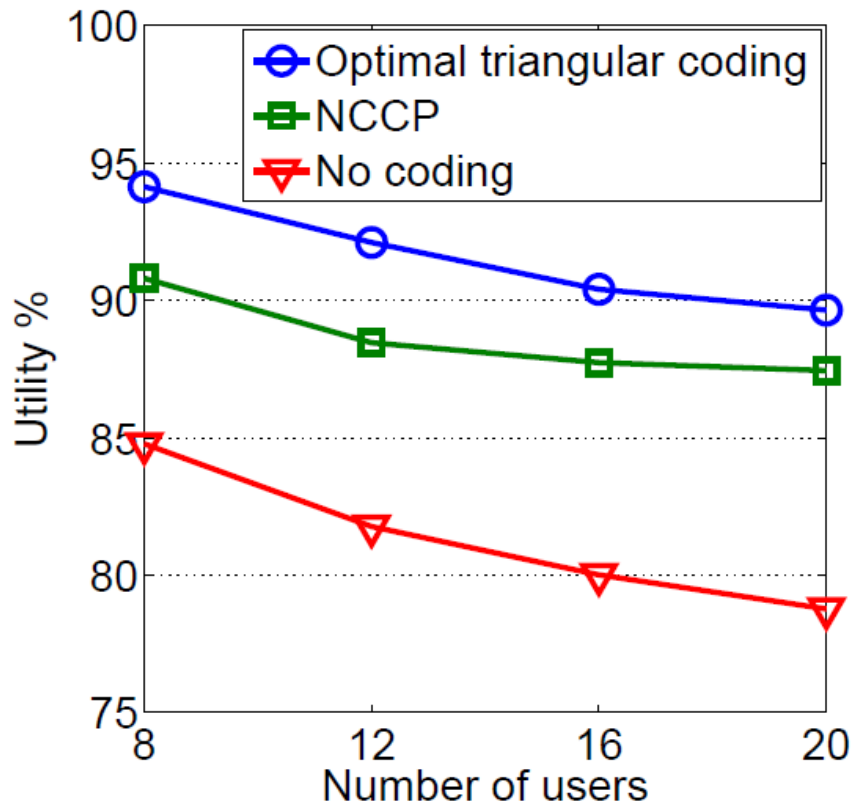
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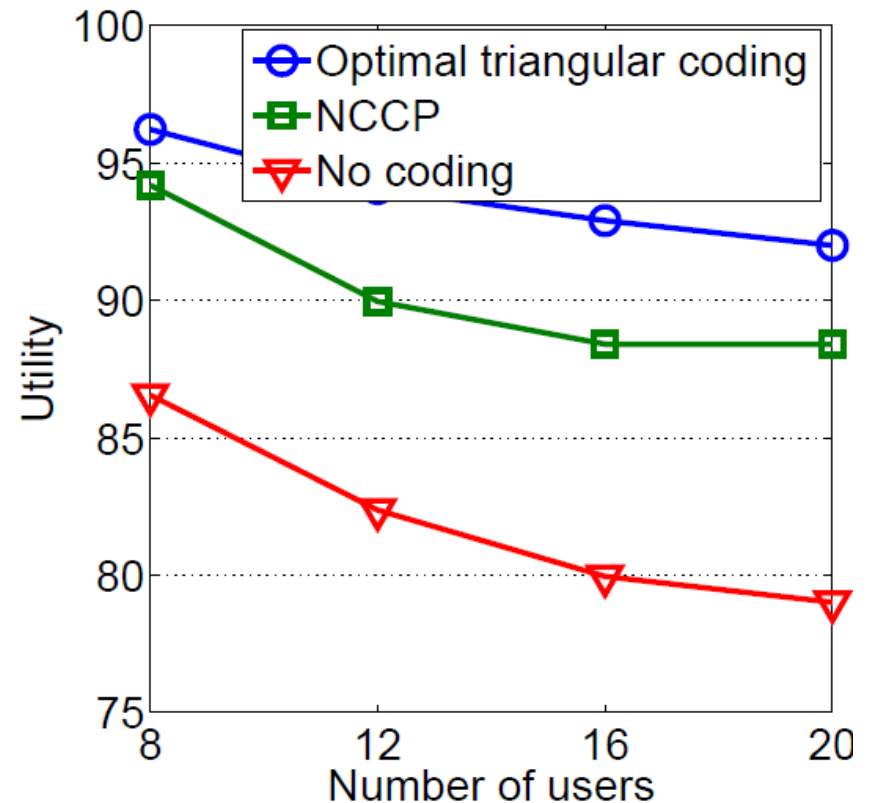
Simulations

13

- Number of caches: 5
- Number of layers: 4



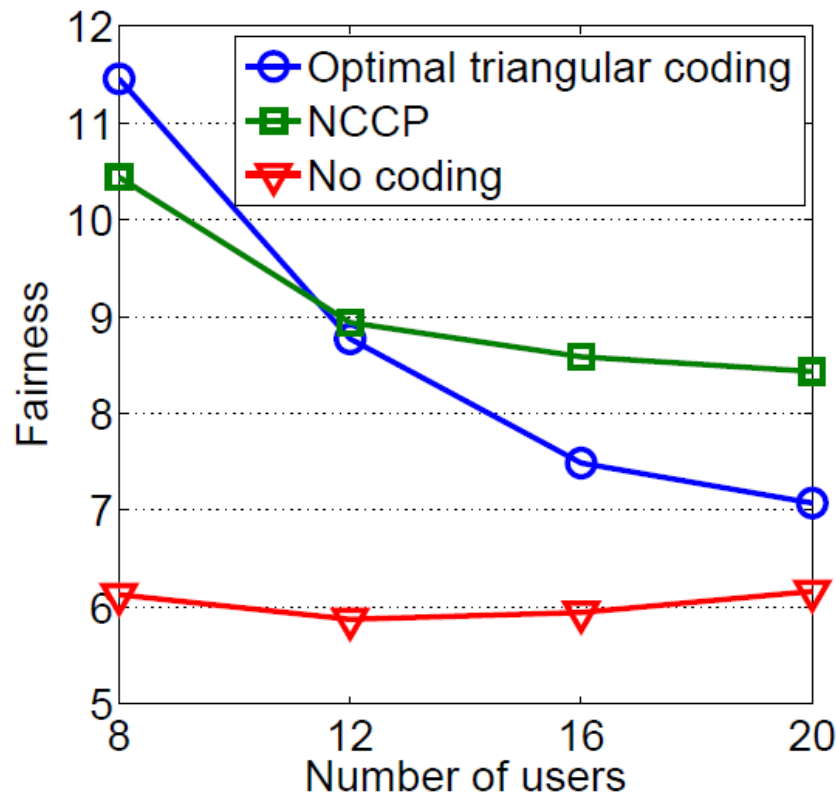
- Number of caches: 5
- Number of layers: 4



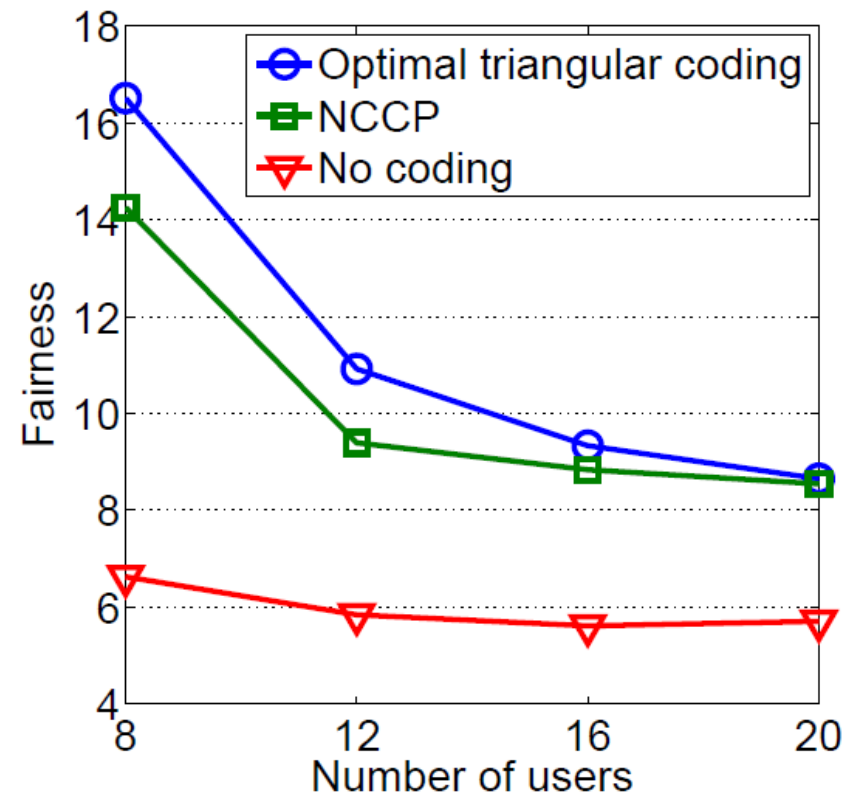
Simulations

14

- Number of caches: 5
- Number of layers: 4



- Number of caches: 5
- Number of layers: 4



Summary

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- The problem of efficient content placement on the caches is known as an NP-complete problem.
- Triangular network coding can reduce the complexity of content placement compared to the general form of coding.
- We propose a heuristic algorithm to solve the problem.

Questions