Shape Clustering: Common Structure Discovery

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Abstract
This paper aims to address the problem of shape clustering by discovering the common structure which captures the intrinsic structural information of shapes belonging to the same cluster. It is based on a skeleton graph, named common structure skeleton graph (CSSG), which expresses possible correspondences between nodes of the individual skeletons of the cluster. To construct the CSSG, we derive the correspondences by the optimal subsequence bijection (OSB). To cluster the shape data, we apply an agglomerative clustering scheme, in each iteration, the CSSGs are formed from each cluster and the two closest clusters are merged into one. The proposed agglomerative clustering algorithm have been evaluated on several shape data sets, including three articulated shape data sets, Torsello’s data set, and a gesture data set. In all experiments, our method demonstrates effective performance compared to other algorithms.

Keywords: Shape, Shape clustering, Skeleton, Common structure, Hierarchical clustering.

1. Introduction
Shape clustering, the task of unsupervised grouping of shapes, is a fundamental problem in computer vision and cognitive perception. It is useful in many applications including speeding up the database retrieval and automatically labeling of objects presented in image collections.
Compared to general clustering, there are two main problems in shape clustering: 1) Shape data has a wide range of intra-class variations like deformation and articulation, which lead to the difficulty of designing robust descriptor to represent the shape data. 2) Most of the shape descriptors are not vectors and it is difficult to convert them to vectors, since they are usually represented by graphs, strings, or trees. Consequently, the similarity between shape descriptors cannot be measured by standard metrics, such as Euclidean distance, directly, and the popular clustering algorithm, such as kmeans, cannot be directly used for shape data.

To capture the multiple variations of the shape, many effective descriptors have been introduced for shape representation [5, 15, 10, 33, 6]. All these shape descriptors are extracted from one single shape, so even the descriptors extracted from the shapes of the same class may be quite different due to the intra-class deformation. In addition, no common intra-class information is considered when computing the pairwise similarities. The performance of shape clustering suffers from these limitations. To solve these problems, we propose a hierarchical shape clustering approach. In statistics, hierarchical clustering is a method of cluster analysis which seeks to build a hierarchy of clusters. Algorithms for hierarchical clustering are generally either agglomerative, in which one starts at the leaves and successively merges clusters together; or divisive, in which one starts at the root and recursively splits the clusters [12]. The reason for applying the hierarchical clustering framework for the shape data is that it can find successive clusters using previously established clusters and the common intra-class information can be extracted from the current clusters to deliver more robust representation for the next iteration.

In this paper, we choose agglomerative strategy, which initializes each shape as a cluster at first, then iteratively merges two most similar clusters into a single cluster and produces one less cluster at the next iteration until the number of clusters is reduced to a desired number. Each cluster is represented by a common structure abstracted from the skeletons of the shapes in the cluster. Our motivation is that shapes in the same class have common structure encoded into their skeletons as shown in Fig. 1. The common structure is defined as a skeleton graph, called common structure skeleton graph (CSSG), in which each node consists of a set of matching nodes (instances) of the individual skeleton graphs of the same class and each edge consists of a set of edges (instances) in the individual skeleton graphs between the matching nodes. Nodes and edges in the CSSG have their weights
related to the number of the instances, which is explicitly defined in Sec. 4.2. The common structure corresponding to multiple shape instances represents the intrinsic information of the class. We utilize this information to improve the similarity measure between shapes. The outline of the proposed skeleton based hierarchical clustering is illustrated in Fig. 2. In the clustering process, once the shapes are merged into one cluster, they are taken as a whole and their CSSG is updated for the next iteration. The distance between clusters is calculated by a proposed measure between their CSSGs.

Our contribution can be divided into three aspects: 1) We represent a cluster of shapes by the CSSG and link the CSSG in our agglomerative shape clustering scheme. 2) We propose a distance measure between the skeletons of individual shapes which also can be used to match the nodes. 3) We extend the above-mentioned measure to computing the distance between clusters and locating the correspondence between the nodes of CSSGs.

The rest of this paper is organized as follows. In Sec. 2, we review the literature on shape clustering. Then we describe how to construct the common structure skeleton graph and propose our agglomerative shape clustering algorithm in Sec. 4. The experimental results on five shape data sets are presented in Sec. 5. Finally, we draw the conclusion in Sec. 6.

2. Related Work

In this section, we review the related works, including shape descriptor extraction, skeleton graph matching and shape clustering.

Generally, the existing shape descriptors can be classified into two types: contour based and skeleton based. Shape context [5] is the most popular
contour based shape descriptor, which describes the relative spatial distribution (Euclidean distance and orientation) of landmark points around feature points on the contour. Ling et al. [15] use inner distance, which is the length of the shortest path within the shape boundary, to replace the Euclidean distance in shape context. As an extension of shape context, inner distance is articulation insensitive. Felzenszwalb et al. [10] utilize a hierarchical model which represents the contour by the segments composed at multiple levels of resolution. This hierarchical representation can capture the important shape variations.

An alternative to contour based representation is to use a structural abstraction, in form of a shape skeleton. The skeleton, also known as medial axis, is defined as the set of centers of all maximally inscribed disks (disks that are contained inside the object but not contained in any other such disk) [6]. The skeleton is a very useful shape descriptor [22], and it can better capture articulation of shapes than contour [21]. Since skeletons contain the structural information of shapes, it is natural to organize them into attribute-relation graphs (ARG), which can be used to measure the similarities between skeletons. Shock graph is a kind of ARG proposed by Siddiqi et al. [26, 27, 25], which is obtained through specialized Shock Grammar [24]. In shock graph, shocks are the branch points, end-points, and those
skeleton segments, which contains both topology and geometry information of the shape [20, 19]. Bone graph is an extension of shock graph, which only retains the non-ligature structures of the shock graph and offers improved stability [16]. Matching ARG is an NPC problem, so several algorithms were proposed to obtain approximate solutions. In [20], Sebastian et al. define the edit distance between shock graphs, however the computational cost is expensive due to the complex edit operations. Methods such as [8, 18] could obtain the correspondence between nodes by converting the skeleton graph to a skeleton tree. However, such a conversion requires selecting one node as the root, which needs heuristic rules as in [8, 18]. As pointed out in [20], a relatively small change in the shape causes the root to change, leading to a significant change in the topology of the tree representation. In addition, the conversion to a tree may result in loss of important structural information, and consequently, negatively influence the matching result [3]. Baseski et al. [4] use rooted-depth-1 tree expressed from Aslan skeleton [1] to represent shape data structure which benefits from the robustness of the disconnection locations of Aslan skeletons with respect to articulations. The dissimilarity between two skeletons is measured by the tree-edit distance involved in the influence of the category contexts. Han et al. [11] apply an EM algorithm to learn both the structure of the supergraph and the correspondences between the nodes of the sample graphs. They mainly focus on the topological structure and the general graph matching problem. Bai et al. [3] represent each end point by the skeletal shortest paths emanating from it and address skeleton graph matching by matching the sequence of end points.

Our method employs the framework of agglomerative clustering. Agglomerative clustering iteratively merges two closest clusters into a single one, therefore, a distance measure between two clusters (groups of instances), also called linkage criteria, must be defined [12]. The linkage criterion is always a function of the pairwise distances between instances. Several linkage criteria have been designed for determining the distance between groups of instances, such as single linkage, complete linkage [29], and average linkage [28]. However, the shape skeleton data has its own intrinsic properties, Therefore, we will design a special distance measure for the shape data.

In recent years, several methods have been proposed for shape clustering. For example, Lakäemper et al. [13] design a distance measure between a single shape and a group of shapes. Then a soft k-means like framework is applied for shape clustering as the basis of the new distance measure. Srivastava et al. [30] use the geodesic paths constructed between shape boundaries
to measure the distance between shapes and perform clustering by using a minimum variance type criterion and a Markov process. In [35], shape is converted to a 1D time series represented by the distance from the centroid of the shape to the contour points. Then a nonlinear projection algorithm is used to group together similar shapes. In [17], the elastic properties of the shape boundaries are encoded in Riemannian metrics and the clustering is applied based on the elastic geodesic distance with DP alignment between shapes.

All the above clustering methods are contour based. Unlike these methods, we represent shapes in terms of their skeletons, which can better capture natural deformations of shapes. Another advantage of our method is that the common structure can be abstracted from skeletons which is useful for clustering. Learning a skeletal shape abstraction from a set of exemplars has been used for shape categorization. Demirci et al. [7] construct the class skeletal prototype based on a many-to-many correspondence between the nodes of skeletons obtained by the method in [8]. However, their prototype is still an exemplar, which is obtained by averaging all the exemplars in the class. Consequently, their prototype may not capture the intra-class structural variation well. Besides, the construction process for their prototype requires converting the skeleton graph to a skeleton tree, which needs heuristic rules as described above. Torsello et al. [32] attempted to find a mixture of tree unions that best accounts for the observed samples using a minimum encoding criterion. Our method is closely related to Torsello’s, however, there are two differences. First, Torsello’s method concentrated on trees, while we construct graphs for clusters. Second, in Torsello’s tree union, only node weight are considered, while in our CSSG, both node weight and edge weight are defined. Erdem and Torsello [9] also proposed a skeleton-based shape clustering method aiming at simultaneously extracting shape classes and learning class-specific shape similarities. While they bias the similarity measure between a shape and a shape cluster by making by using of some statistical values of the skeletal attributes in the cluster, we improve the similarity measure between the shape clusters based on the common structures abstracted from them.

3. Basic Skeleton Concepts

To better describe our approach, we give some basic skeleton concepts first. We follow the definitions of end point, junction point, connection
point, skeleton graph, skeleton path and path distance in [3]: The
skeleton point with only one adjacent point is an end point; the skeleton
point with more than two adjacent points is a junction point. If a skeleton
point is not an end point or a junction point, it is called a connection point.
(For skeletons in digital images, we assume the curves of the skeleton are one-
pixel wide.) The end/junction points and the sequences of connecting points
between two end/junction points form the nodes and the edges of the skeleton
graph respectively. An end point and a junction point in a skeleton graph are
called end node and junction node respectively. A skeleton path $P(u,v)$ is
the shortest path between a pair of nodes $u, v$ in a skeleton graph. Given are
two skeleton paths $P(u_x, v_x), P(u_y, v_y)$ represented by two vectors of the radii
of their maximal disks centered at the $M$ sample points: $(r_{xi}; i = 1, \ldots, M)$
and $(r_{yi}; i = 1, \ldots, M)$, respectively. The path distance between $P(u_x, v_x)$
and $P(u_y, v_y)$ is

$$pd(P(u_x, v_x), P(u_y, v_y)) = \sum_{i=1}^{M} \frac{(r_{xi} - r_{yi})^2}{r_{xi} + r_{yi}} + \alpha \frac{(l_x - l_y)^2}{l_x + l_y},$$

(1)

where $\alpha$ is the weight factor and $l_x$ and $l_y$ are the lengths of $P(u_x, v_x)$
and $P(u_y, v_y)$, respectively. Both the radii and the length are normalized to
ensure the path distance is scale invariant.

Besides the above definitions, we also give some others here: The skeleton
path between a pair of end nodes is called an end path. The skeleton path
between a junction node and an end node is called a junction path. For two
junction nodes $j_1, j_2$, if there are no other junction nodes in the skeleton path
between them, $j_1$ and $j_2$ are connected directly, denoted by $j_1 \sim j_2$.

4. Agglomerative Shape Clustering

In this section, we introduce our agglomerative shape clustering method.
First we propose our distance measure between shapes, then extend it to
the distance measure between clusters. The main idea of our approach is to
define distance measures between skeletons by extending the notion of the
path distance introduced in [3] which allows for the treatment of endpoints
first, followed by junction points, building up to the concept of a CSSG for
agglomerative clustering.

We utilize junction paths to represent shapes, since junction points con-
tain the structural information of shapes [34], as shown in Fig. 3. However,
Figure 3: A junction point and its junction paths.

Figure 4: The junction points on each shape are merged into three clusters (marked by the blue circles) to match well those on the other shape.

Junction points are not stable [3], they may suffer from ligature-induced instability [16]. For example, in Fig. 4, there are actually five junction points in the first horse’s skeleton. We observe that if junction points \( j_2 \) and \( j_3 \) are merged together as well as \( j_4 \) and \( j_5 \), then the skeleton graph structure will resemble the structure of the second horse. The instability of shape skeletons has been studied in the literature [1, 36, 16, 3, 34]. In [16], the instable ligatures are removed which leads to a more stable skeletal representation. Here, we modify the merging strategy in [34] as a preprocessing step for our approach. The detail of the preprocessing step are described in the Appendix.

4.1. Pairwise Distance between Shapes

Our basic idea for a distance measure between shapes is to find the optimal correspondence between end nodes first, then using the optimal correspondence between end nodes to determine the one between junction nodes and the distance between shapes. The reason why we do not match the junction nodes antecedent to matching end nodes is that finding the optimal correspondence between end nodes benefits from the prior that they have the order along the shape contour. Since junction nodes do not have a natural order, they suffer from instability [3], therefore it is much harder to match
junction nodes directly.

For two skeletons $S_x$ with $n_x$ end nodes and $m_x$ junction nodes and $S_y$ with $n_y$ end nodes and $m_y$ junction nodes, their end nodes are $E_x = \{e_{x1}, \ldots, e_{nx}\}$ and $E_y = \{e_{y1}, \ldots, e_{ny}\}$, respectively. The goal is to obtain an optimal correspondence $\varphi : \{e_{x1}, \ldots, e_{nx}\} \to \{e_{y1}, \ldots, e_{ny}, \phi\}$, where $e_x \in E_x$ is mapped to $\varphi(e_x) \in E_y$, and we allow a many-to-one mapping to $\phi$. This is a sequence matching problem, so we compute all distances between the end nodes in $E_x$ and those in $E_y$ and obtain a $n_x \times n_y$ matrix $\Phi(E_x, E_y)$:

$$\Phi(E_x, E_y) = \begin{pmatrix}
ed(e_{x1}, e_{y1}) & ed(e_{x1}, e_{y2}) & \cdots & ed(e_{x1}, e_{y}n) \\
ed(e_{x2}, e_{y1}) & ed(e_{x2}, e_{y2}) & \cdots & ed(e_{x2}, e_{y}n) \\
\vdots & \vdots & \ddots & \vdots \\
ed(e_{xn}, e_{y1}) & ed(e_{xn}, e_{y2}) & \cdots & ed(e_{xn}, e_{y}n) \\
\end{pmatrix}, \quad (2)$$

where $ed(\cdot, \cdot)$ is the distance between two end nodes computed by applying Optimal Subsequence Bijection (OSB) [14] to the matrix whose entry is the path distance between any two end paths emanating from the two end nodes, see [3] for more details.

In [3], the Hungarian algorithm is applied to $\Phi(E_x, E_y)$ to obtained the optimal correspondence between the end nodes in $E_x$ and those in $E_y$. Here, we make a small change, since the order of the sequence is not considered in the Hungarian algorithm, we apply OSB to the matrix $\Phi(E_x, E_y)$ to obtain the optimal correspondence, which ensures the consistency of the orders of the pair of end nodes along the contours.

Now we introduce an optimal correspondence between junction nodes and a pairwise distance measure between shapes. For two skeletons $S_x$ and $S_y$, their junction nodes are $J_x = \{j_{x1}, \ldots, j_{xm}\}$ and $J_y = \{j_{y1}, \ldots, j_{ym}\}$, respectively. The optimal correspondence $\varphi$ between end points have been obtained by the approach described above. Assuming that $\kappa$ pairs of end nodes are matched and $\epsilon$ end nodes are not matched, as illustrated in Fig. 5. The distance between two junction nodes $j_x \in J_x$ and $j_y \in J_y$ is

$$jd(j_x, j_y) = \frac{1}{\kappa} \left( \sum_{e \in E_m} pd(P(j_x, e), P(j_y, \varphi(e))) + \epsilon \zeta \right), \quad (3)$$

where $E_m \subseteq E_x$ is the matched end nodes set ($E_m = \{e_{x1}, e_{x2}, e_{x3}, e_{x5}, e_{x6}\}$ in Fig. 5) and $\zeta$ is the penalty factor for the unmatched end nodes which is defined as:
Figure 5: Computing the distance between junction nodes $j_x$ and $j_y$ based on the matched of end nodes. There are $\kappa = 5$ matched pairs of end nodes and $\epsilon = 2$ unmatched end nodes.

The penalty factor is defined as the mean plus one standard deviation (std) of the distances between the closest skeleton paths is to ensure the distance between the outlier and its closest skeleton path would be larger than the penalty factor, so that the outliers would be excluded from the matching with a relatively small penalty. This so defined penalty factor is inspired by the definition of jump cost in the sequence matching algorithm proposed by Latecki et al. [14]. Note that, $\zeta$ defined in Eq. 4 is not symmetric with respect to $S_x$ and $S_y$, so we switch the roles of $S_x$ and $S_y$, compute the value of $\zeta$ twice. Then the final penalty factor $\zeta$ is the average of the two values.

Then we compute all distances between the junction nodes in $J_x$ and those in $J_y$ and obtain a $mx \times my$ matrix $\Phi(J_x, J_y)$:

$$
\Phi(J_x, J_y) = 
\begin{pmatrix}
jd(j_{x1}, j_{y1}) & jd(j_{x1}, j_{y2}) & \cdots & jd(j_{x1}, j_{yny}) \\
jd(j_{x2}, j_{y1}) & jd(j_{x2}, j_{y2}) & \cdots & jd(j_{x2}, j_{yny}) \\
\vdots & \vdots & \ddots & \vdots \\
jd(j_{xnx}, j_{y1}) & jd(j_{xnx}, j_{y2}) & \cdots & jd(j_{xnx}, j_{yny}) 
\end{pmatrix}
$$

(5)
Figure 6: Common structure skeleton graph construction. Colors indicate the correspondences between the nodes in the CSSG and those in the individual skeletons. The instances belong to $u_{ex}$, $u_{jx}$ and $u_{jy}$ are $\{e_{x1}, e_{x2}, e_{x3}\}$, $\{j_{x2}, j_{x3}\}$ and $\{j_{y1}, j_{y2}, j_{y3}\}$, respectively. Hence, the junction path $P(u_{jx}, u_{ex}) = \{P(j_{x2}, e_{x2}), P(j_{x3}, e_{x3})\}$ and $P(u_{jy}, u_{ex}) = \{P(j_{y1}, e_{x1}), P(j_{y2}, e_{x2}), P(j_{y3}, e_{x3})\}$. The numbers next to the nodes in the CSSG are their weights.

Finally, we apply the Hungarian algorithm to the matrix $\Phi(J_x, J_y)$ to obtain the optimal correspondence $\psi : \{j_{x1}, \ldots, e_{jmax}\} \rightarrow \{j_{y1}, \ldots, j_{nymy}, \phi\}$ between the junction nodes and the distance between the pair of shapes $S_x, S_y$, since junction nodes do not have any natural order:

$$jpd(S_x, S_y) = \mathcal{H}(\Phi(J_x, J_y)),$$

where $\mathcal{H}(\cdot)$ is the Hungarian function. We refer to this distance measure as JPD, for junction path distance.

4.2. Common Structure Skeleton Graph Construction

The CSSG is constructed from a cluster of shapes. As shown in Fig. 6, each node in the CSSG consists of a set of matching nodes (instances) of the skeletons in this cluster. The number of the instances that belong to the node in the CSSG gives us a significant information: A node that consists of few instances is more likely an outlier, while a node which consists of many instances is an important node. Thus, we can assign different weights to different nodes according to the number of their instances. As the numbers shown in Fig. 6, for a node $v$ in the CSSG constructed by $\lambda$ shapes, its weight is $w_N(v) = \frac{|v|}{\lambda}$, where $|\cdot|$ is the cardinality of a set.

As for the edge weight in the CSSG, it is computed by simply counting the edges in the original skeleton graphs between nodes that belong to two nodes.
in the CSSG. For an edge \( P(v_x,v_y) \) between two nodes \( v_x \) and \( v_y \) in the CSSG, it consists of a set of edges that belong to the individual skeletons, whose element is \( P(v_{x_i},v_{y_i}) \), where \( v_{x_i} \in v_x, v_{y_i} \in v_y \) and the same subscript of “\( v_x \)” and “\( v_y \)” means the nodes are in the same skeleton. The weight of edge \( P(v_x,v_y) \) is simply its cardinality: \( W_E(P(v_x,v_y)) = |P(v_x,v_y)| \). Fig. 6 shows an example of computing the edge weight: The edge \( P(u_{jx},u_{ex}) \) between two nodes \( u_{jx} = \{j_{x2},j_{x3}\} \) and \( u_{ex} = \{e_{x1},e_{x2},e_{x3}\} \), then \( P(u_{jx},u_{ex}) = \{P(j_{x2},e_{x2}),P(j_{x3},e_{x3})\} \), thus, the weight \( w_E(P(u_{jx},u_{ex})) = 2 \). Similarly, the edge \( P(u_{jy},u_{ex}) = \{P(j_{y1},e_{x1}),P(j_{y2},e_{x2}),P(j_{y3},e_{x3})\} \), and its weight \( w_E(P(u_{jy},u_{ex})) = 3 \).

4.3. Distance Measure between Common Structure Skeleton Graphs

Now we extend the junction path distance to measure the distance between CSSGs. The CSSG is also a skeleton graph, therefore, as computing junction path distance, we obtain the optimal correspondence between end nodes of the CSSGs first, then based on it, determine the optimal correspondence between junction nodes and the distance between CSSGs.

Each node in the CSSG consists of a set of matched nodes of the skeletons in the cluster. Therefore, in order to measure the distance between CSSGs, we need to measure the distance between sets of corresponding nodes. For two CSSGs \( CS_x \) with \( nx \) end nodes and \( mx \) junction nodes and \( CS_y \) with \( ny \) end nodes and \( my \) junction nodes, their end nodes are \( UE_x = \{ue_{x1},...,ue_{xny}\} \) and \( UE_y = \{ue_{y1},...,ue_{yny}\} \), respectively, and their junction nodes are \( UJ_x = \{uj_{x1},...,uj_{xmx}\} \) and \( UJ_y = \{uj_{y1},...,uj_{ymy}\} \), respectively. Consider two end nodes \( ue_x = \{e_{x1},...,e_{xly}\} \subset UE_x \) and \( ue_y = \{e_{y1},...,e_{yly}\} \subset UE_y \), where \( lx \) and \( ly \) are the cardinalities of the sets. Intuitively, for two CSSGs formed by the shapes of the same class, an important node in one is less likely to match an outlier in the other. Therefore, we define the distance between \( ue_x \) and \( ue_y \) as

\[
ucd(ue_x,ue_y) = \frac{\sum_i \sum_k ed(e_{xi},e_{yj})(1 + \frac{\|w_N(ue_x) - w_N(ue_y)\|}{w_N(ue_x) + w_N(ue_y)})}{lx \cdot ly},
\]

where \( \| \cdot \| \) denotes the absolute value. Note that, by introducing the penalty factor \( \frac{\|w_N(ue_x) - w_N(ue_y)\|}{w_N(ue_x) + w_N(ue_y)} \), two important nodes are more likely to match each other and an outlier is less likely to match an important node. If both \( ue_x \) and \( ue_y \) are important nodes, then both \( w_N(ue_x) \) and \( w_N(ue_y) \) are large, thus \( \|w_N(ue_x) - w_N(ue_y)\| \) and \( w_N(ue_x) + w_N(ue_y) \) are small and large respectively.
Consequently, the penalty factor \( \frac{w_N(ue_x) - w_N(ue_y)}{w_N(ue_x) + w_N(ue_y)} \) is very small. While, without loss of generality, if \( ue_x \) and \( ue_y \) are an important node and an outlier respectively, then \( \|w_N(ue_x) - w_N(ue_y)\| \) and \( w_N(ue_x) + w_N(ue_y) \) are larger and smaller respectively. Consequently, the penalty factor \( \frac{w_N(ue_x) - w_N(ue_y)}{w_N(ue_x) + w_N(ue_y)} \) is much larger than the one computed in the former case. If both \( ue_x \) and \( ue_y \) are outliers, then both \( \|w_N(ue_x) - w_N(ue_y)\| \) and \( w_N(ue_x) + w_N(ue_y) \) are small, since there is no hypothesis that the outliers are prone to match each other. Then we compute a \( nx \times ny \) matrix of the distances of end nodes as we did in Sec. 4.1:

\[
\Phi(U E_x, U E_y) = \\
\begin{pmatrix}
ued(ue_{x_1}, ue_{y_1}) & ued(ue_{x_1}, ue_{y_2}) & \cdots & ued(ue_{x_1}, ue_{y_n}) \\
ued(ue_{x_2}, ue_{y_1}) & ued(ue_{x_2}, ue_{y_2}) & \cdots & ued(ue_{x_2}, ue_{y_n}) \\
\vdots & \vdots & \ddots & \vdots \\
ued(ue_{x_n}, ue_{y_1}) & ued(ue_{x_n}, ue_{y_2}) & \cdots & ued(ue_{x_n}, ue_{y_n})
\end{pmatrix}.
\] (8)

Similarly, we apply OSB to the matrix \( \Phi(U E_x, U E_y) \) to obtain the optimal correspondence \( \varphi : \{ue_{x_1}, \ldots, ue_{x_n}\} \rightarrow \{ue_{y_1}, \ldots, ue_{y_n}, \phi\} \) between end nodes of the two CSSGs.

As for the optimal correspondence between junction nodes and the pairwise distance between CSSGs, we apply the algorithm with similar form as introduced in Sec. 4.1. Each edge in a CSSG corresponds to multiple skeleton paths, so path distance defined in [3] can be adapted to compute the distance between the edges in CSSGs. For two junction nodes \( uj_x = \{j_{x_1}, \ldots, j_{x_n}\} \subseteq U J_x \) and \( uj_y = \{j_{y_1}, \ldots, j_{y_n}\} \subseteq U J_y \), where \( ix \) and \( iy \) are the cardinalities of the sets. For two junction paths \( P(uj_x, ue_x) \) and \( P(uj_y, ue_y) \), the distance between them is

\[
upd(P(uj_x, ue_x), P(uj_y, ue_y)) =
\sum_{e_{x_i} \in P(uj_x, ue_x)} \sum_{e_{y_k} \in P(uj_y, ue_y)} \frac{Pd(P(j_{x_1}, e_{x_1}), P(j_{y_1}, e_{y_k}))}{w_E(P(uj_x, ue_x)) \cdot w_E(P(uj_y, ue_y))},
\] (9)

where the same subscript of “j” and “e” means the junction node and the end node are in the same skeleton graph.

For the two CSSGs \( CS_x \) and \( CS_y \), assuming that \( \kappa \) pairs of end nodes are matched and \( \epsilon \) end nodes \( U E_x \) are skipped (not matched). Skipping an outlier is trivial to the matching, while skipping an important node increases the dissimilarity between the two CSSGs. Therefore, we would like to assign
Figure 7: Computing the distance between junction nodes $uj_x$ and $uj_y$ based on the matched of end nodes. There are $\kappa = 5$ matched pairs of end nodes and $\epsilon = 2$ unmatched end nodes $UE = \{ue_{x4}, ue_{y6}\}$.

A penalty to each skipped node according to their weights. As shown in Fig. 7, the distance between $uj_x$ and $uj_y$ is

$$ujd(uj_x, uj_y) = \frac{1}{\kappa} \left( \sum_{ue \subseteq UE_m} upd(P(uj_x, ue), P(uj_y, \varphi(ue))) + \zeta \sum_{ue \notin UE} w_N(ue)(1 + \frac{||w_N(uj_x) - w_N(uj_y)||}{w_N(uj_x) + w_N(uj_y)}) \right),$$

(10)

where $UE_m \subseteq UE_x$ is matched end nodes set ($UE_m = \{ue_{x1}, ue_{x2}, ue_{x3}, ue_{x5}, ue_{x6}\}$ in Fig. 7) and $\zeta$ is the penalty factor for the unmatched end nodes and defined as:

$$\zeta = \text{mean}_{uj_x \in UJ_x, ue_x \in UE_x} \left( \min_{uj_y \in UJ_y, ue_y \in UE_y} upd(P(uj_x, ue_x), P(uj_y, ue_y)) \right) +$$

$$\sqrt{\text{std}_{uj_x \in UJ_x, ue_x \in UE_x} \left( \min_{uj_y \in UJ_y, ue_y \in UE_y} upd(P(uj_x, ue_x), P(uj_y, ue_y)) \right)},$$

(11)

We also switch the roles of $CS_x$ and $CS_y$, compute the value of $\zeta$ twice, and use the average as the final penalty factor $\zeta$. Notice that Eq. 3 is the special case of Eq. 10 when the common structure only corresponds to a single instance.

Then we compute all distances between the junction nodes in $UJ_x$ and those in $UJ_y$ and obtain a $mx \times my$ matrix $\Phi(UJ_x, UJ_y)$:
\[ \Phi(UJ_x, UJ_y) = \begin{pmatrix}
ujd(u_{jx_1}, u_{jy_1}) & ujd(u_{jx_1}, u_{jy_2}) & \cdots & ujd(u_{jx_1}, u_{jy_N}) \\
ujd(u_{jx_2}, u_{jy_1}) & ujd(u_{jx_2}, u_{jy_2}) & \cdots & ujd(u_{jx_2}, u_{jy_N}) \\
\vdots & \vdots & \ddots & \vdots \\
ujd(u_{jx_N}, u_{jy_1}) & ujd(u_{jx_N}, u_{jy_2}) & \cdots & ujd(u_{jx_N}, u_{jy_N})
\end{pmatrix}. \quad (12) \]

Finally, we apply the Hungarian algorithm to the matrix \( \Phi(UJ_x, UJ_y) \) to obtain the optimal correspondence \( \psi : \{u_{jx_1}, \ldots, u_{jx_N}\} \rightarrow \{u_{jy_1}, \ldots, u_{jy_N}\} \) between the junction nodes and the distance between the pair of CSSGs \( CS_x, CS_y \):

\[ csd(CS_x, CS_y) = \mathcal{H}(\Phi(UJ_x, UJ_y)). \quad (13) \]

We refer to this distance measure between clusters as CSD, for common structure distance.

### 4.4. Clustering Scheme

Now we demonstrate our clustering scheme. Given an unlabeled shape set \( X = \{x_1, x_2, \ldots, x_N\} \) to be clustered into \( \Theta \) clusters \( \mathcal{C} = \{C_1, C_2, \ldots, C_\Theta\} \), where \( N \) is the number of shapes, and their skeletons are \( SS = \{s_1, s_2, \ldots, s_N\} \). The agglomerative shape clustering scheme is shown in Fig. 8.

<table>
<thead>
<tr>
<th>Procedure Clustering (Input SS, ( \Theta ), Output ( \mathcal{C} ))</th>
</tr>
</thead>
<tbody>
<tr>
<td>1. Initialize ( \mathcal{C} = {C_1, C_2, \ldots, C_N} ) and ( \theta = N ), where ( C_i = {s_i}_{i=1}^N ).</td>
</tr>
<tr>
<td>2. While ( \theta &gt; \Theta )</td>
</tr>
<tr>
<td>3. Let ( CS_i ) be the CSSG of ( C_i (i = 1, 2, \ldots, \theta) ), compute its node weights and edge weights.</td>
</tr>
<tr>
<td>4. Find the closest two common structures ( CS_i ) and ( CS_k ), ( (i, k = 1, 2, \ldots, \theta, i \neq k) ).</td>
</tr>
<tr>
<td>5. Merge ( C_i ) and ( C_k ) into one: ( \hat{C} = C_i \cup C_k ).</td>
</tr>
<tr>
<td>6. Update ( \mathcal{C} : \mathcal{C} = \mathcal{C}\setminus C_i, \mathcal{C} = \mathcal{C}\setminus C_k, \mathcal{C} = \mathcal{C} \cup \hat{C}, \theta = \theta - 1. )</td>
</tr>
<tr>
<td>7. End</td>
</tr>
</tbody>
</table>

Figure 8: The agglomerative shape clustering algorithm.

We refer to the proposed shape clustering method as CSD + AHC, for agglomerative hierarchical clustering based on common structure distance.
5. Experimental Results

To assess the quality of the proposed clustering method, we evaluate it on five standard data sets: Torsello’s data set [32], “50 hands” data set [17], and three Aslan and Tari data sets [2, 1, 4]. The following clustering methods are considered for comparison in all our experiments: The pairwise distance between shapes by inner distance [15] with normalized cuts [23] used to cluster shapes, denoted by IDSC + Ncuts; Junction path distance with normalized cuts, denoted by JPD + Ncuts; JPD but with normalized cuts replaced by agglomerative hierarchical clustering with average linkage, denoted by JPD + AHC; the proposed CSD + AHC. Our experiments are divided into two parts. We commence by illustrating qualitative examples of the clusters obtained by different clustering methods. Then we give quantitative comparison between the clustering methods. All results are obtained by tuning the desired number of clusters to be equal to the natural number of classes of the data set. To apply Ncuts algorithm, the distances between shapes must be converted to similarities first. This usually can be done by using a Gaussian kernel. Assuming that the distance between two shapes $S$ and $S'$ is $D(S, S')$, then the similarity between them is obtained by $W(S, S') = \exp(-D^2(S, S')/\sigma^2)$. In our experiments, we set $\sigma = 10.0$ empirically.

5.1. Clustering Examples

We begin with the illustration of clusters on a small data set provided by Torsello [32] that contains 25 shapes grouped in 9 classes. We set $\alpha = 150$ for this dataset. We compare the proposed CSD + AHC to the union of attributed trees which obtains the best clustering result in [32], IDSC + Ncuts, JPD + AHC and JPD + Ncuts. Fig. 9 shows the result of comparison, we set the number of clusters $\Theta = 9$ for all the clustering methods. As shown in Fig. 9, comparing with JPD + Ncuts, JPD + AHC and IDSC + AHC, CSD + AHC achieves the perfect result which shows the advantage of common structure.

We also test our method on the Aslan and Tari data set 56 shapes [2], which includes 14 classes of articulated shapes with 4 shapes in each class. We set $\alpha = 100$ for this dataset and set the number of clusters $\Theta = 14$ for all the clustering methods. Fig. 10 shows the clustering result obtained by CSD + AHC. There are only two errors: four windmills are clustered with pentagrams and a bone is clustered into a new cluster. However, the windmill is very similar to the pentagrams so that the error seems acceptable. The
Figure 9: Comparison between different clustering methods on Torsello’s data set [32].
(a) CSD + AHC. (b) JPD + AHC. (c) JPD + Ncuts. (d) IDSC + Ncuts. (e) Union of attributed trees [32].

Cluster results obtained by IDSC + Ncuts, JPD + Ncuts, JPD + AHC are shown in Fig. 11, Fig. 12 and Fig. 13, respectively. Clearly, CSD + AHC achieves the best result. The results on the Aslan and Tari data set with 56 shapes depict the advantage of JPD for measuring the distance between articulated shapes and how CSD + AHC works in clustering. Fig. 14 illustrates several results of common structures obtained by CSD + AHC. The left column shows the common structure of clusters and the other columns show skeleton graphs from shapes which belong to the cluster. Note that, the end nodes and the junction nodes marked without number are the out-
Figure 10: Clusters of the Aslan and Tari data set with 56 shapes by CSD + AHC.

Figure 11: Clusters of the Aslan and Tari data set with 56 shapes by IDSC + Ncuts.

Outliers detected by the common structure, as the one in the skeleton of the fourth turtle. Fig. 15 depicts the hierarchical property of CSD + AHC. We show the clustering results obtained by turning different desired numbers of clusters. Note that, if the clustering is allowed to be continued a bit further (i.e. $\Theta = 13$), cat and horse shapes will be grouped together, since they are two visually similar shape classes.
5.2. Quantitative Analysis

To quantitatively analyze the clustering results, we use the normalized mutual information (NMI) as a measure of clustering results. The ground-truth class partition $\Gamma$ and the returned cluster partition $\Delta$ define a confusion matrix with each entry $n_{i}^{(j)}$ being the number of data samples in cluster $i$ and class $j$, and $n$ is the total number of samples. Then NMI is computed as follows:

$$
\text{NMI} = \frac{2 \sum_{i=1}^{I} \sum_{j=1}^{J} \frac{n_{i}^{(j)}}{n} \log \frac{n_{i}^{(j)}n}{\sum_{k=1}^{J} n_{k}^{(j)} \sum_{k=1}^{I} n_{k}^{(i)}}}{H(\Gamma) + H(\Delta)}
$$

where $I$ is the number of clusters and $J$ is the number of classes. $H(\Gamma) = -\sum_{i=1}^{I} \frac{n_{i}}{n} \log \frac{n_{i}}{n}$ and $H(\Delta) = -\sum_{j=1}^{J} \frac{n_{j}}{n} \log \frac{n_{j}}{n}$ are the entropies of parti-
The common structure of several shape classes in the Aslan and Tari data set 56 shapes [2]. The end nodes (in red) and junction nodes (in green) are marked with the numbers. The same number indicates the correspondences between nodes in the sample shapes and nodes in the common structure. The junction node and the end node in the skeleton of the fourth turtle marked without numbers have low weight, therefore, we do not show them in the common structure.
Figure 15: The hierarchical property of CSD + AHC shown on the Aslan and Tari data set 56 shapes [2]. The shapes in the same rectangular region are grouped into a cluster by CSD + AHC. The rectangular regions are marked by two colors alternately to make them visible. From top to bottom, the desired number of clusters is decreased: (a) Θ = 36. (b) Θ = 24. (c) Θ = 14. (d) Θ = 13.

Table 1: The NMIs of the clustering results on Torsello's data set [32].

<table>
<thead>
<tr>
<th></th>
<th>CSD + AHC</th>
<th>JPD + AHC</th>
<th>JPD + Ncuts</th>
<th>IDSC + Ncuts</th>
<th>[32]</th>
</tr>
</thead>
<tbody>
<tr>
<td>NMI</td>
<td>1.0000</td>
<td>0.9618</td>
<td>0.7778</td>
<td>0.6431</td>
<td>0.9313</td>
</tr>
</tbody>
</table>

Table 2: The NMIs of the clustering results on Aslan and Tari data set with 56 shapes [2].

<table>
<thead>
<tr>
<th></th>
<th>CSD + AHC</th>
<th>JPD + AHC</th>
<th>JPD + Ncuts</th>
<th>IDSC + Ncuts</th>
</tr>
</thead>
<tbody>
<tr>
<td>NMI</td>
<td>0.9734</td>
<td>0.8674</td>
<td>0.6174</td>
<td>0.5660</td>
</tr>
</tbody>
</table>

tion Γ and ∆, respectively. A high value of NMI indicates that Γ and ∆ match well.

We list the NMIs of the clustering results on Torsello’s data set and the Aslan and Tari data set with 56 shapes in Tab. 1 and Tab. 2, respectively.

“50 hands” data set [17] consists of 50 boundary shapes of hands in different postures. The natural number of clusters in this data set is (depending on subjective interpretation) 8 or 9, with cluster sizes between 2 and 8 elements. The groundtruth clusters of these two cases are illustrated in Fig. 16. We set α = 100 for this dataset. Several clustering methods are compared on this data set, including the method in [17], the method in [13], IDSC + Ncuts, JPD + Ncuts, JPD + AHC and our CSD + AHC. As shown in Tab. 3, whether the desired number of clusters Θ = 9 or Θ = 8, CSD + AHC achieves the perfect result.
The groundtruth clusters of data set '50 hands', taken from Mio et al. [17]. The natural numbers of clusters are 9 and 8 in (a) and (b) respectively.

Table 3: The NMIs of the clustering results on “50 hands” data set [17].

<table>
<thead>
<tr>
<th>Method</th>
<th>NMI (Θ = 9)</th>
<th>NMI (Θ = 8)</th>
</tr>
</thead>
<tbody>
<tr>
<td>CSD + AHC</td>
<td>1.0000</td>
<td>1.0000</td>
</tr>
<tr>
<td>[17]</td>
<td>1.0000</td>
<td>-</td>
</tr>
<tr>
<td>[13]</td>
<td>0.6055</td>
<td>1.0000</td>
</tr>
<tr>
<td>IDSC + Ncuts</td>
<td>1.0000</td>
<td>0.6732</td>
</tr>
<tr>
<td>JPD + Ncuts</td>
<td>0.8403</td>
<td>0.8568</td>
</tr>
<tr>
<td>JPD + AHC</td>
<td>0.9726</td>
<td>0.9331</td>
</tr>
</tbody>
</table>

We also test cluster methods on a larger data set, the Aslan and Tari data set with 180 shapes, which includes 30 classes of articulated shapes with 6 shapes in each class. We set \( \alpha = 100 \) for this data set and set the number of clusters \( \Theta = 30 \) for all the clustering methods. Tab. 4 shows the clustering results obtained by IDSC + Ncuts, JPD + Ncuts, JPD + AHC and CSD + AHC. Our result is much better than the results of other methods.

Finally, we test cluster methods on a much larger data set, the Aslan and Tari data set with 1000 shapes, which includes 50 classes of articulated shapes with 20 shapes in each class. We set \( \alpha = 100 \) for this data set and set the number of clusters \( \Theta = 50 \) for all the clustering methods. Tab. 5 shows the clustering results obtained by IDSC + Ncuts, JPD + Ncuts, JPD + AHC, CSD + AHC and Foreground Focus [37]. The result of Foreground Focus is quoted from [9]. Other results in [9] are not obtained by tuning the desired number of clusters to be equal to the natural number of classes of the data set, so we do not quote them here. In Tab. 5, our result is still the best, which illustrates the effectiveness of the common structure on large data set.

To assess the ability of the proposed clustering algorithm to classify the
Table 4: The NMI values of the clustering results on Aslan and Tari data set with 180 shapes.

<table>
<thead>
<tr>
<th>Method</th>
<th>NMI</th>
</tr>
</thead>
<tbody>
<tr>
<td>CSD + AHC</td>
<td>0.9694</td>
</tr>
<tr>
<td>IDSC + Ncuts</td>
<td>0.5423</td>
</tr>
<tr>
<td>JPD + Ncuts</td>
<td>0.5785</td>
</tr>
<tr>
<td>JPD + AHC</td>
<td>0.8793</td>
</tr>
</tbody>
</table>

Table 5: The NMI values of the clustering results on Aslan and Tari data set with 1000 shapes.

<table>
<thead>
<tr>
<th>Method</th>
<th>NMI</th>
</tr>
</thead>
<tbody>
<tr>
<td>CSD + AHC</td>
<td>0.8096</td>
</tr>
<tr>
<td>IDSC + Ncuts</td>
<td>0.5433</td>
</tr>
<tr>
<td>JPD + Ncuts</td>
<td>0.4549</td>
</tr>
<tr>
<td>JPD + AHC</td>
<td>0.7693</td>
</tr>
<tr>
<td>Foreground Focus</td>
<td>0.7329</td>
</tr>
</tbody>
</table>

We perform experiments on an increasing number of shapes in the two Aslan and Tari data sets. As for the Aslan and Tari data set with 56 shapes, we commence with the 8 shapes from the first 2 classes and then increase the number of shape classes under consideration until the full set of 56 shapes was included. We plot the curves of NMI as the number of shape classes is increased in Fig. 17. Similarly, we plot the curves obtained by all the cluster methods on the Aslan and Tari data set with 180 shapes in Fig. 18. We commence with 30 shapes from the first 5 classes and then increase the number of shape classes until all 180 shapes were included. In both Fig. 17 and Fig. 18, CSD + AHC appears as the curve marked with green “square” and clearly outperforms other clustering methods.

5.3. Parameter Discussion

There is a parameter introduced in this paper, the threshold $T_m$ for merging junction points. Bigger threshold means more junction points would be merged together. $T_m$ is a constant in all our experiments: $T_m = 5$. There is another parameter introduced in [3], the weight factor $\alpha$. In Fig. 19, we show that the clustering result is not sensitive to the choice for $\alpha$.

6. Conclusion

In this paper, we present an agglomerative hierarchical algorithm for shape clustering based on a common structure formed by the shapes belonging to the same clusters. Unlike the general clustering methods whose results only depend on the pairwise similarities, the proposed clustering method extracts the common structure which captures the intrinsic intra-class structural information of the cluster of shapes. Consequently, it can be used to
Figure 17: The curves of NMI as the number of shape classes is increased on the Aslan and Tari data set with 56 shapes [2].

deliver more robust distance measure between clusters. The presented experimental results demonstrate that our shape clustering algorithm significantly outperforms other state-of-the-art methods.

The high time complexity is the limitation of our method, since the overall time required for a basic agglomerative hierarchical clustering is $O(N^2 \log(N))$ [31], where $N$ is the number of data points. However, we can speed up the clustering by merging more than two closest clusters in each iteration, which will be the topic of our future work.

7. Acknowledgements

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Figure 18: The curves of NMI as the number of shape classes is increased on the Aslan and Tari data set with 180 shapes [1].

Figure 19: The curves of NMI as the weight factor $\alpha$ is varied on the four data sets.


8. Appendix

Preprocessing. A merging strategy given in [34] says that any pair of junction points are merged when the distance between them are less than a threshold. However, this strategy is not described clearly. For example, both the distance between junction points $j_1$ and $j_2$ and the one between
junction points \( j_1 \) and \( j_2 \) are less than the threshold, however the distance between junction points \( j_2 \) and \( j_3 \) is larger than the threshold. Should \( j_2 \) and \( j_3 \) be merged together? Here, we propose to cluster the junction points by agglomerative hierarchical clustering with single linkage. At first, each junction point starts in its own cluster, then in each iteration, two closest clusters of junctions are merged as one moves up the hierarchy. The distance measure between clusters is always based on the distance measure between pairs of closest instances.

For two junction points \( j_1 \) and \( j_2 \) on the skeleton \( S \) with \( n \) end points \( E = \{ e_1, \ldots, e_n \} \), the distance between them is

\[
d(j_1, j_2) = \begin{cases} \frac{1}{n} \sum_{i=1}^{n} pd(P(j_1, e_i), P(j_2, e_i)) & \text{if } j_1 \sim j_2 \\ \infty & \text{otherwise} \end{cases}.
\]

(14)

To cluster the junction points on skeleton \( S \), we should give the distance measure between clusters. For two cluster junction points \( JP, JQ \) on skeleton \( S \), the distance between them determined by the single linkage criterion is

\[
cd(JP, JQ) = \begin{cases} \min_{j_p \in JP, j_q \in JQ} d(j_p, j_q) & \text{if } JP \sim JQ \\ \infty & \text{otherwise} \end{cases},
\]

(15)

where \( JP \sim JQ \) means \( JP \) and \( JQ \) are connected directly, i.e., there are no junction points of other clusters in any paths between the junction points of \( JP \) and \( JQ \).

Suppose that all junctions had been merged into \( m \) clusters \( \{ JP_i \}_{i=1}^m \), the stopping criterion of the merge process is

\[
\min_{i,k} cd(JP_i, JP_k) \geq T_m, (i,k = 1, 2, \ldots, m, i \neq k),
\]

(16)

where \( T_m \) is a threshold. The larger threshold means less clusters, i.e. more junction points would be merged together. Since the path distance is scale invariant, the threshold \( T_m \) is scale invariant too. We set \( T_m = 5 \) empirically in all our experiments.

The junction points in one shape merged into the same cluster are considered as a whole for the next computation. This means a junction node in a skeleton graph corresponds a cluster of junction points. To compute the distance between two junction nodes \( JP, JQ \), actually, we choose \( jp^* \in JP \) and \( jq^* \in JQ \) to represent \( JP \) and \( JQ \), respectively, if satisfying

\[
(jp^*, jq^*) = \arg \min_{jp \in JP, jq \in JQ} jd(jp, jq).
\]

(17)
Then the distance between $JP$ and $JQ$ is $jd(jp^*, jq^*)$. Other distances related to junction nodes are computed by the similar way.
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1. A method is proposed to discover the common structure of a cluster of shapes.
2. A cluster of shapes are represented by their CSSG.
3. Assigning weights to nodes and edges in CSSGs to help measure their distances.
4. An agglomerative strategy is proposed for shape clustering based on CSSGs.