

**Figure 4.13** The Bellman-Ford algorithm for single-source shortest paths in general graphs.

procedure shortest-paths( $G, l, s$ )

Input: Directed graph  $G = (V, E)$ ;  
 edge lengths  $\{l_e : e \in E\}$  with no negative cycles;  
 vertex  $s \in V$

Output: For all vertices  $u$  reachable from  $s$ ,  $\text{dist}(u)$  is set  
 to the distance from  $s$  to  $u$ .

for all  $u \in V$ :  
 $\text{dist}(u) = \infty$   
 $\text{prev}(u) = \text{nil}$

$\text{dist}(s) = 0$

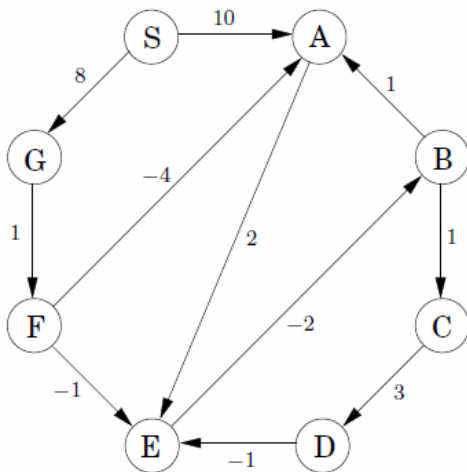
repeat  $|V| - 1$  times:

for all  $e \in E$ :  
 update( $e$ )

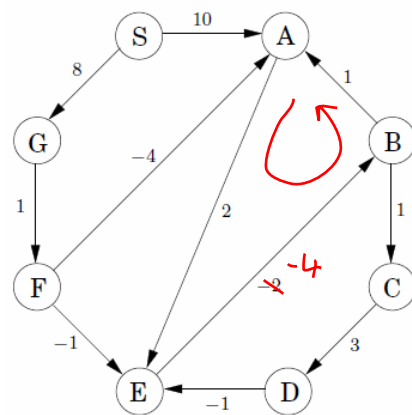
procedure update( $(u, v) \in E$ )

$\text{dist}(v) = \min\{\text{dist}(v), \text{dist}(u) + l(u, v)\}$

**Figure 4.14** The Bellman-Ford algorithm illustrated on a sample graph.



Node	Iteration							
	0	1	2	3	4	5	6	7
S	0	0	0	0	0	0	0	0
A	$\infty$	10	10	5	5	5	5	5
B	$\infty$	$\infty$	$\infty$	10	6	5	5	5
C	$\infty$	$\infty$	$\infty$	$\infty$	11	7	6	6
D	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	14	10	9
E	$\infty$	$\infty$	12	8	7	7	7	7
F	$\infty$	$\infty$	9	9	9	9	9	9
G	$\infty$	8	8	8	8	8	8	8



Negative cycle: