Artificial Intelligence: Graph Based Search Techniques

Presented by Michael Otte



These slides were prepared with help from: "Artificial Intelligence, A Modern Approach" by Stuart Russell and Peter Norvig

Graph Based Search

- What I am going to talk about in the next three classes:
 - Uninformed (blind) Search
 - Informed (heuristic based) Search
 - Current applications to artificial intelligence
 - Robotics
 - Computer Games
 - Part 3 of your project
 - Implement a search algorithm (A*) that will be used for navigation on the LAGR robot.

Graph Based Search

- What is a graph?
 - Nodes (states)
 - Initial State
 - Goal State
 - Arcs (the links between the states)
 - Directed
 - Undirected
 - Cost
 - Other stuff:
 - Successor Function
 - Goal Test

Graph Based Search

• A simple Graph:



I am going to do more interesting example on the board (with corresponding solution tree).

Graph Based Search: Uninformed Search

- Solving a graph can be thought of as searching through a tree until a goal state is found.
- Any traditional tree search methods can be used (but some may work better than others)
 - Depth-First Search
 - Breadth-First Search
 - Depth-Limited Search
 - Iterative Deepening Depth-First Search
 - Bidirectional Search
 - Dijkstra's Algorithm

Uninformed Search: Depth-First Search

- Depth: The number of nodes along the path from the initial state.
- Depth-first search always expands the deepest node in the current *fringe* of the search tree.
 - Fringe: the set of nodes that the search has found
- It is possible to implement depth first search by storing the fringe in a stack.
- Remember to mark a node as visited when it is expanded.

Uninformed Search: Depth-First Search

- Depth first Requires O(bm) memory.
 - *b*: the (maximum) branching factor
 - *m*: the maximum depth of the tree
- A variant called *backtracking search* reduces memory requirements to O(m) by letting each node remember which of its children have been explored.
- Problem: what happens if the graph contains an infinite number of nodes?
- Depth-first search is not complete or optimal

Uninformed Search: Breadth-First Search

- Breadth-first search expands nodes in the order of their depth.
 - That is, every node of depth n is expanded before any node of depth n+1 is expanded.
 - This can be implemented in a first-in-first-out queue
- Memory and time: O(b^{d+1})
 - b: branching factor
 - d: depth of goal
- Breadth-first search is complete, it is optimal if all steps have the same cost.

Uninformed Search: Depth-limited and Iterative Deepening

- Depth-limited search is (usually depth-first) search with a constraint on the length of a path.
 - Choose the constraint carefully or risk failure
- Iterative Deepening expands this idea:
 - Perform depth-limited search with depth 1,2,3...
 - Memory (depth-first based): O(bd)
 - Time: O(b^d)
 - This is faster than breadth-first search! Why?
 - Because most nodes exist on the bottom level of the tree and breadth-first search inserts an extra layer of nodes into the fringe.

Uninformed Search: Dijkstra's Shortest Path Algorithm

- Used when arcs have arbitrary non-negative costs.
- Basic idea: expand the shortest path first and store the distance required to reach each node at that node, updating as necessary.
 - runtime (if implemented carefully with a Fibonacci heap) is O(E + V log V)
 - *E*: total edges in the graph
 - V: total vertices in the graph
- Optimal and Complete

Graph Based Search: Informed Search

- Sometimes we can use additional information about the world to help guide our search.
- The idea is that we can impart a little bit of 'common sense' about what is good or bad to the system using a *heuristic function*.
- Algorithms:
 - Hill-climbing search
 - Genetic Algorithms
 - Greedy best-first search
 - A* Search (You will have to implement this one)

Informed Search (Local Search) Hill-Climbing Search

- Suppose that we are told to find the best possible state in a space with respect to a particular quantity.
 - But, the space is huge and we can only make a few measurements a time.
 - Picture a mouse that is trying to find the top of a roller-coaster track.
- Example on board...

Informed Search (Local Search) Hill-Climbing Search

- Start at a random location and go up
 - If space is convex, then this will always find a global solution, if not it will find a local solution.
 - Steepest ascent: choose the highest neighbour
 - Stochastic Hill Climbing: choose randomly from uphill moves (this makes more sense in > 2D)
 - Random Restart: try a bunch of different things.
 - Now we have a pack of mice instead of just one.
 - Local Beam Search: try a bunch of different things, and allow them to communicate.
 - If one of the mice finds a good state, it tells its friends about it.
 - Simulated Annealing: temperature and probability
 - So much for the mouse metaphor...

Informed Search (Local Search) Genetic Algorithms

- With some probability:
 - Independently successful features of individuals are randomly combined to form a new individual.
 - A new individual is created by mutating an old individual.
 - A cool example:

http://www.rennard.org/alife/english/gavgb.html

Back to global path finding...

Graph Based Search: Uninformed Search Revisited

On the Quiz from last time:

- Depth First Search: 75%
- Breadth First Search: 75%
- Dijkstra's: 25%

I think that I'll work through a few more examples of these things before we move on.

- Another way of thinking about it:
 - Recursively expand Nodes based on the order that you discover them in.
- That means
 - Do not return back to the parent of a node before all of the children of the current node have been explored.
- Memory Requirement Confusion from last time:
 - O(db) if we store the unexpanded nodes
 - d=depth
 - b=branching factor
 - O(d) if we let each node remember which children have been expanded.



Suppose our successor function breaks ties by expanding the left most node

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What we know





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Suppose our successor function breaks ties by expanding nodes counter clockwise, starting at 12:00.

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- Another way of thinking about it:
 - First, look at all nodes 1 arc away from the start
 - Then, look at all nodes 2 arcs away from the start
 - Next, look at all nodes 3 arcs away from the start
 - etc ...
- When you are told about nodes that are too far away, save them for later.



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- Fibonacci heap high level concepts (Wikipedia):
 - A forest of trees
 - The trees are arranged in a ring (doubly linked list)
 - Each tree is a doubly linked list
 - Always keep a pointer to the over all minimum
 - Lazy operations (some things are saved for later)
 - 'Merge' just combines forests into a larger ring
 - 'extract min' basically merges trees along the way
 - Idecrease key' makes new trees along the way
 - 'delete' sets the node's value to -inf then extracts min
 - Special constraints:
 - Size of a subtree rooted at node of degree k is at least F_{k+2}, where F_k is the kth Fibonacci number.
 - In their time analysis they have a notion of saving time for later.

- Another way of thinking about it:
 - Always expand the node that is the closest to the start based on cost.
 - When you expand a node, save the back pointer along the cheapest path.
 - If during an expansion you find a cheaper way to get to a child (neighbor), update the child so that it knows about the cheapest way back to the start.



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Informed Search Greedy Best-First Search

- Greedy best-first search: expands the node on the fringe that is closest to the goal.
- In 2D or 3D Cartesian space, if cost = distance
 - Expand the node that is closest to the goal first.
 - f(n) = h(n)
 - Trip navigation example on the board.

- You are going to implement this as part 3 of your project.
- Idea: minimize the total estimated solution cost.
- Expand Nodes that have a low value of:

$$f(n) = g(n) + h(n)$$

- g(n): actual cost from the start to a node n
- h(n): estimated cost from node n to the goal
 - If h(n) = 0 this reduces to Dijkstra's

- (This is part 3 of your project).
- Idea: minimize the total estimated solution cost.
- Expand the nodes that has the lowest value of: f(n) = g(n) + h(n)
- g(n): actual cost from the start to a node n
- h(n): estimated cost from node n to the goal
 - If h(n) is an admissible and consistent then A* is optimal.
 - That is, if h(n) never overestimates the cost to the goal, and maintains the triangle inequality, then A* will always find the best solution.
 - How fast this happens depends on how good your heuristic is.

- How to apply this to navigation in:
 - real environments (robotics)
 - simulated environments (computer games)
- Step 1: model the world as a 4 or 8-connected graph (i.e. a chessboard).
 - Graph locations are defied by row and column.
 - Neighbours are row ±1 and column ±1
 - Let each node contain cost information about the world—now the graph is called an occupancy grid.
- Step 2: think of a good heuristic
 - Any ideas...
- Step 3: implement A* with your heuristic.

- Step 4: make some code that does this:
 - Sense or input information about your environment
 - figure out your start and goal locations
 - run A* from start to goal on your occupancy grid
 - Move along the resulting path a little bit
 - Repeat
- This is already implemented on the LAGR robot
 - All you have to do is implement A* with your own heuristic.
- Example on board...
 - Cat chasing a mouse in an environment that has cost given by: dirt = 1, grass 3, and water = 10

A* pseudo-code (found on Wikipedia):

- function A*(start,goal)
- **var** closed := the empty set
- var q := make_queue(path(start))
- while q is not empty
- **var** p := remove_first(q)
- **var** x := the last node of p
- if x in closed
- continue
- if x = goal
- return p
- add x to closed
- **foreach** y **in** successors(x)
- enqueue(q, p, y)
- **return** failure

- Grading for part 3
 - 50 % Does it work off-line on maps that I give it?
 - 25 % Explanation of why you chose a particular heuristic.
 - 25 % Does it run in real-time on the LAGR robot and make good decisions?
 - You'll need to have it running at about 1 frame per second on a map of 100 X 100 grids