Throughput and Fairness-Aware Dynamic Network Coding in Wireless Communication Networks

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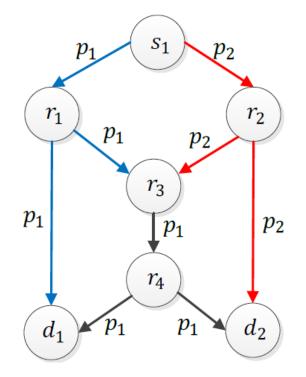


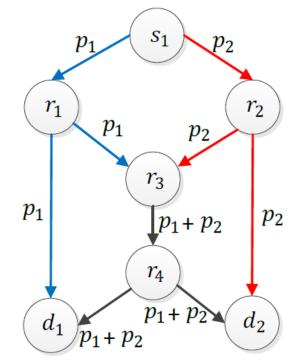
Agenda

- Introduction
- Motivation and setting
- Proposed methods
 - Dynamic network coding
 - Fair dynamic network coding
- Simulation results
- Conclusion



Bottleneck problem



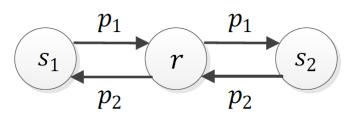


Without coding

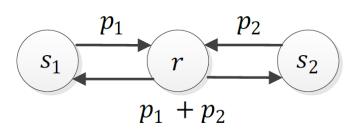
With coding

Network Coding in Wireless Networks

No coding
 4 transmissions



Coding
3 transmissions

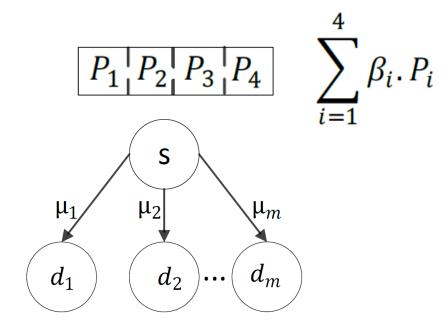


Inter-flow coding
Increases the throughput



Network Coding in Wireless Networks

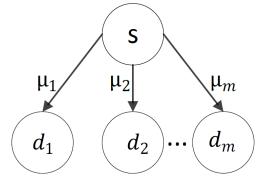
- Intra-flow coding
 - Reliability



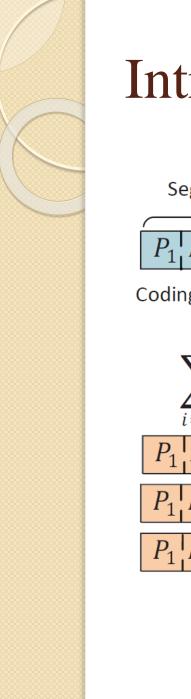


Setting

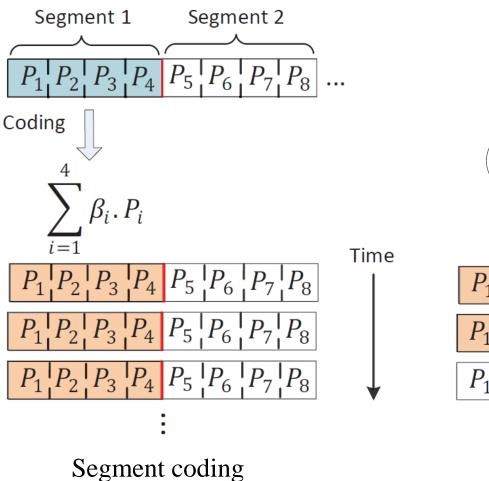
- One source
 - Broadcasts a set of packets
- Multiple destinations

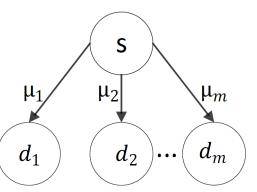


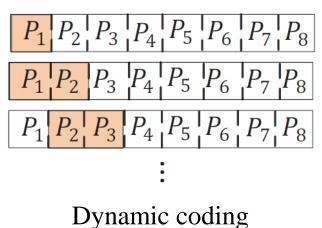
- Independent erasure channels
- Equal size time slots
 - One packet transmission per time slot
- Objective
 - Throughput



Introduction (Segment Coding)









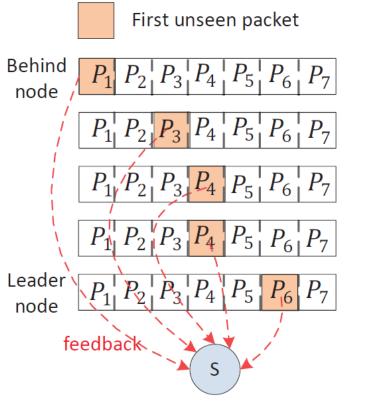
Introduction

- Seen packet (Sundararajan'08)
 - A node has seen a packet P if it can generate a linear combination of the form P + Q, using the received coded packets in its buffer
 - $\circ \ \alpha_{1,1} p_1 + \alpha_{2,1} p_2 + \alpha_{3,1} p_3$
 - $\begin{tabular}{ll} & \alpha_{1,1}p_1 + \alpha_{2,1}p_2 + \alpha_{3,1}p_3 \\ & \alpha_{1,2}p_1 + \alpha_{2,2}p_2 + \alpha_{3,2}p_3 \end{tabular} \end{tabular}$
 - Seen packets can be removed from the sender's buffer



Introduction

• ARQ with network coding (ANC)



Coded $P_1 + P_3 + P_4 + P_6$ packet:



Idea

- Behind and leader nodes
- Code packets in the range of the first unseen packets by the leader and behind nodes

Behind
node
$$P_1 P_2 P_3 P_4 P_5 P_6 P_7$$

 $P_{11} P_2 P_3 P_4 P_5 P_6 P_7$
 $P_{11} P_2 P_3 P_4 P_5 P_6 P_7$
 $P_{11} P_2 P_3 P_4 P_5 P_6 P_7$
Leader
node
 $P_{11} P_2 P_3 P_4 P_5 P_6 P_7$
feedback
s
 $P_1 P_2 P_3 P_4 P_5 P_6 P_7$



Multiple Behind and Leader Nodes

First unseen packet

Behind node

Behind node

Leader node

Leader node

 $P_1 P_2 P_3 P_4 P_5 P_6 P_7$ $P_1 | P_2 | P_3 | P_4 | P_5 | P_6 | P_7$ $P_1 P_2 P_3 P_4 P_5 P_6 P_7$

 $P_1 P_2 P_3 P_4 P_5 P_6 P_7$

 $P_1 P_2 P_3 P_4 P_5 P_6 P_7$ $P_1 P_2 P_3 P_4 P_5 P_6 P_7$

• 2 methods to deal with multiple behind and leader nodes

Dynamic NC without Overhearing

- All leaders need to transmit a feedback
 - A receiver that missed the last transmission cannot be a leader node
 - If the index of the first unseen packet is equal to the largest index included in the received coded packet, then the node is a leader node

• Behind nodes

• If all the behind nodes receive the current transmissions, they do not send any feedback messages

Dynamic NC with Overhearing

- Two feedbacks per time slot
- Just one leader and one behind node send feedback
 - Set a back-off time based on the erasure rate of the nodes
 - The receivers listen to the channel
 - Leader node finishes its back-off time
 - Send feedback if has not overheard feedback from the other leaders

Dynamic NC with Overhearing

- Two feedbacks per time slot
- Just one leader and one behind node send feedback
 - The behind nodes that have received the last transmissions do not need to transmit a feedback
 - Only one of the nodes that was a behind node in the previous slot, and missed the current transmission should send a feedback



Throughput

- In ANC each transmission has innovative information for all of the nodes
 - Achieves the maximum throughput
 - Proof
- The same approach can be used to prove that the DNC is throughput optimal

Fair Dynamic NC

- Unfairness of ANC and DNC
 - The nodes with low error rates receive more coded packets than the other nodes, and become the leaders
 - The nodes with higher error rates might not be able to decode the packet for a long time

Fair Dynamic NC

• A trade-off between fairness and throughput

$$x = (1 - w) \times L - w \times (m - L)$$

- w: fairness factor
- *L* : number of leaders
- *m* : number of users
- If x>0, the sender adds a new packet to the coded packet



Simulations (Definitions)

• Decoding delay unfairness

$$f'_D(j) = \frac{\sum_{i=1}^m |D(i,j) - \bar{D}(j)|}{m}$$

- Decoding delay fairness $f_D(j) = \frac{1}{f'_D(j)}$
- Decoding unfairnes

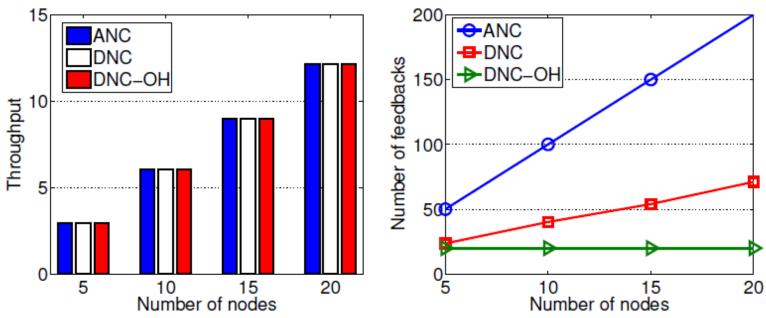
$$f'_E = \frac{\sum_{i=1}^m |E(i) - E|}{m}$$



Simulations

- ANC: ARQ with NC
- DNC: Dynamic network coding without overhearing
- DNC-OH: Dynamic network coding with overhearing

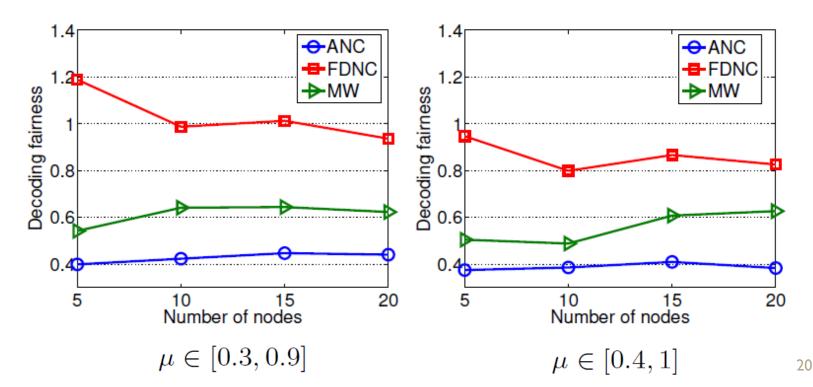
 $\mu \in [0.3,0.9]$



Simulations (Decoding Fairness)

- ANC: ARQ with NC
- FDNC: Fain dynamic network coding
- MW: Moving window

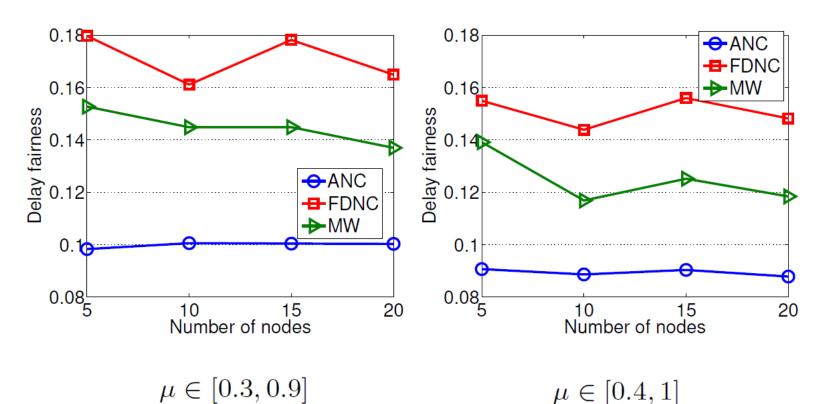
w = 0.7





Simulations (Delay Fairness)

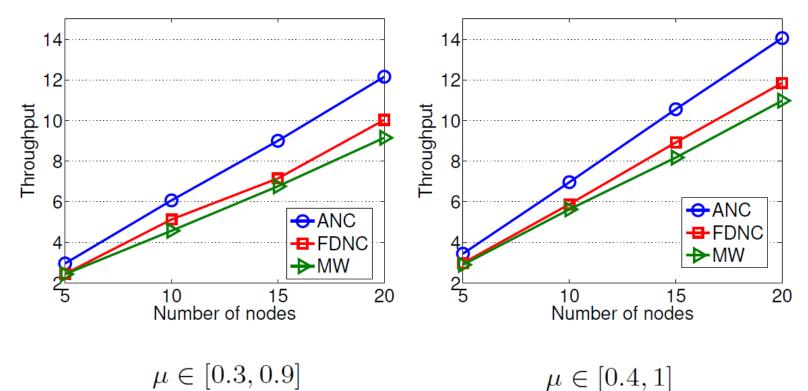
w = 0.7





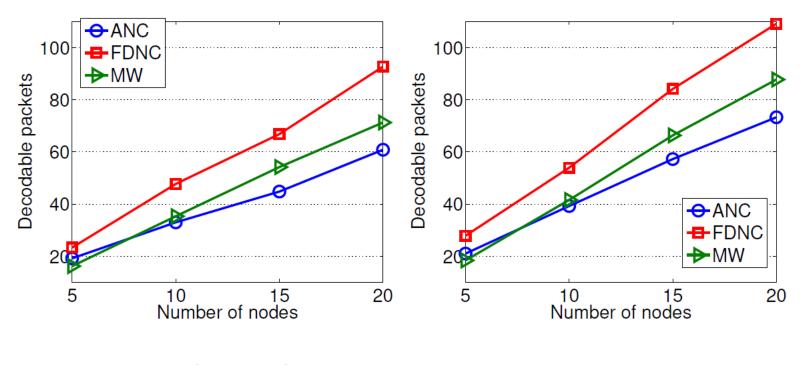
Simulations (Throughput)

w = 0.7





w = 0.7



 $\mu \in [0.3, 0.9]$

 $\mu \in [0.4, 1]$



Summary

- Dynamic coding increases the throughput of network coding
 - Too many feedback messages
- We propose the DNC and DNC-OH methods to reduce the number of feedbacks
- We propose the FDNC method to provide decoding and decoding delay fairness

Questions