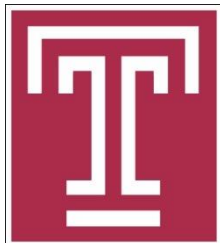


2-Dominant Resource Fairness: Fairness-Efficiency Tradeoffs in Multi-resource Allocation

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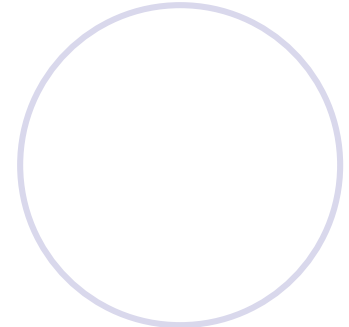
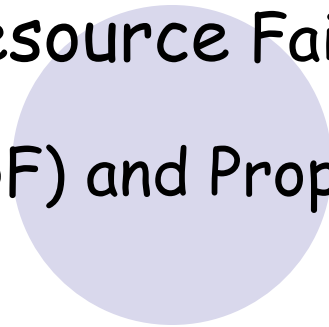
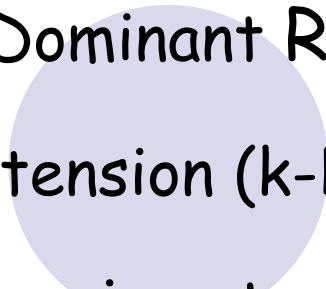
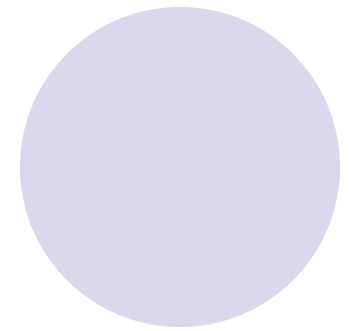
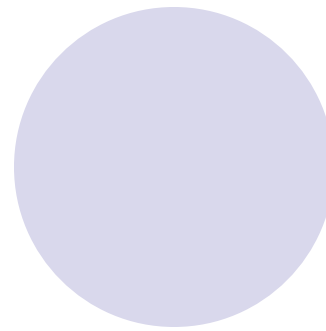
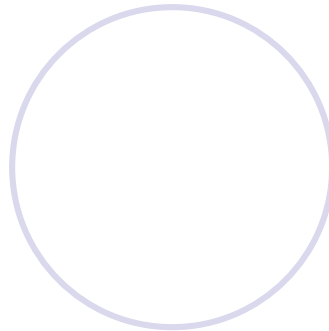
Temple University, USA



Road Map



- Introduction
- Motivation
- 2-Dominant Resource Fairness (2-DF)
- Extension (k-DF) and Properties
- Experiment
- Conclusions



1. Introduction

Multi-resource allocation

- ✦ Sharing more than one type of resource
Bandwidth, Memory, CPU, etc

- ✦ Users have heterogeneous resource demands

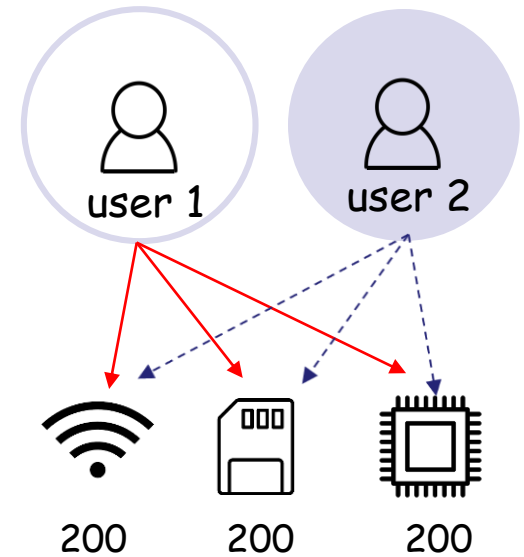
EX.1 Three resources: Bandwidth, Memory, CPU

Total resource: $\langle 200, 200, 200 \rangle$ units

User 1 requires $\langle 40, 8, 8 \rangle$ units / task

User 2 requires $\langle 8, 5, 1 \rangle$ units / task

- ✦ How to fairly/efficiently allocate all resources among users



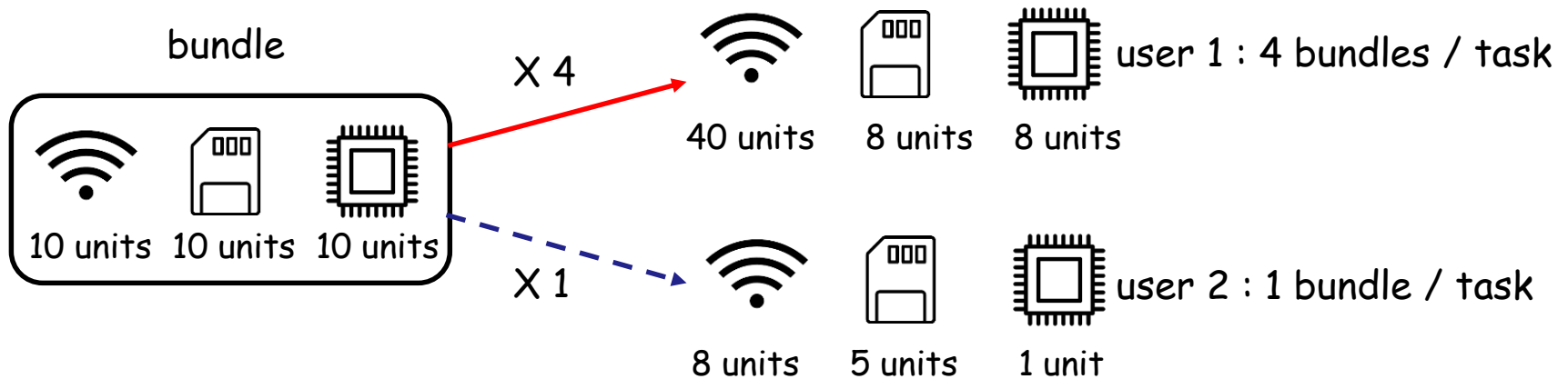
Resource Abstraction

🔗 All resources are partitioned into bundles

- ✦ each bundle has fixed amounts of different resources
- ✦ multiple resources are abstracted as a single resource

🔗 Drawbacks

- ✦ ignore different demands of heterogeneous users
- ✦ cannot always match nicely with users' demands



Dominant Resource Fairness ^[1] (DRF)

↳ Dominant resource

✦ the resource that a user has the biggest share of

↳ Dominant share

✦ the fraction of the dominant share a user is allocated

↳ DRF allocation mechanism

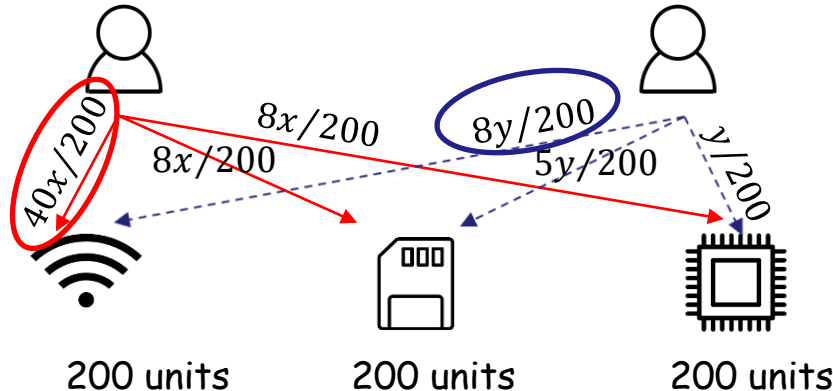
✦ applying max-min fairness to dominant shares

user 1: x tasks

user 2: y tasks

maximize

x, y



subject to

$$\begin{cases} 40x/200 = 8y/200 \\ 40x + 8y \leq 200 \\ 8x + 5y \leq 200 \\ 8x + y \leq 200 \end{cases}$$

$$x = 2.5$$

$$y = 12.5$$

[1] A. Ghodsi, M. Zaharia, B. Hindman, A. Konwinski, S. Shenker, and I. Stoica. "Dominant resource fairness: fair allocation of multiple resource types." In NSDI, 2011.

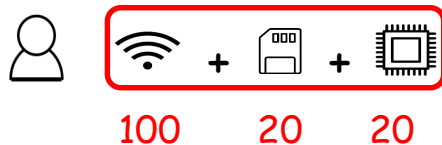
Dominant Resource Fairness (DRF)

Properties

- ✦ sharing incentive
- ✦ strategy proof
- ✦ envy-free
- ✦ Pareto efficient

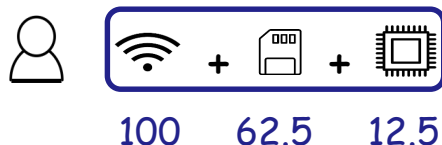
DRF Allocation

user 1 2.5 tasks




100 20 20

user 2 12.5 tasks




100 62.5 12.5

user 1 2.5 tasks




100 100 100

user 2 12.5 tasks




100 100 100

user 1 2.5 tasks



100 20 20

user 2 12.5 tasks




100 90 90

< 8, 5, 1 >


< 10, 9, 9 >

user 1 1.6 tasks




100 62.5 12.5

user 2 4 tasks




100 20 20

user 1 1 task



40 8 8

user 2 20 tasks



160 100 20

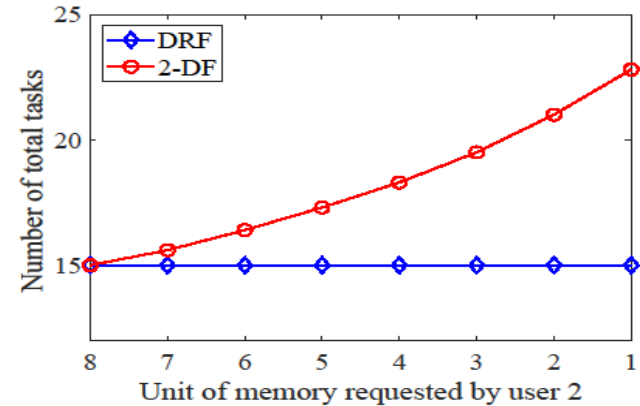
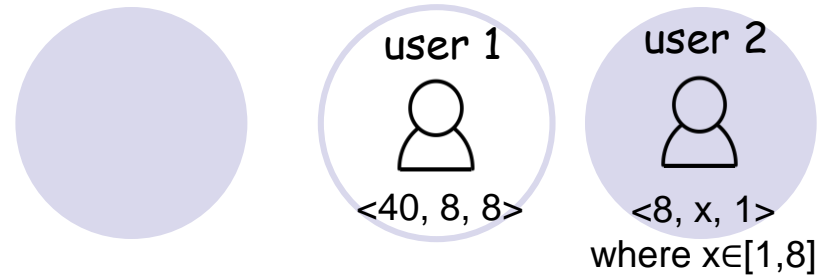
2. Motivation

Fairness dispute in DRF

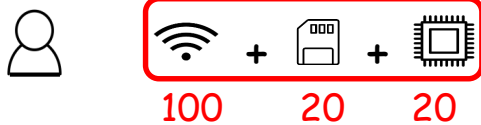
Focus on one resource

Efficiency loss in DRF

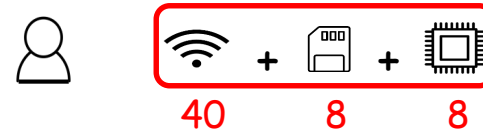
DRF less efficiently uses resources [2]



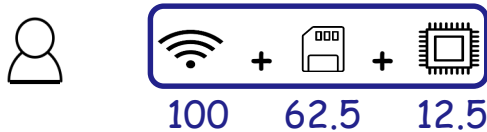
user 1 2.5 task



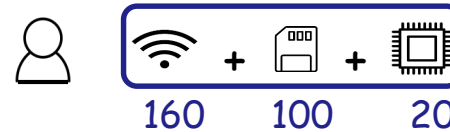
user 1 1 task



user 2 12.5 tasks



user 2 20 tasks



[2] Y. Jin and M. Hayashi. "Efficiency comparison between proportional fairness and dominant resource fairness with two different type resources." in CISS 2016.

Metrics on Fairness and Efficiency

Fairness

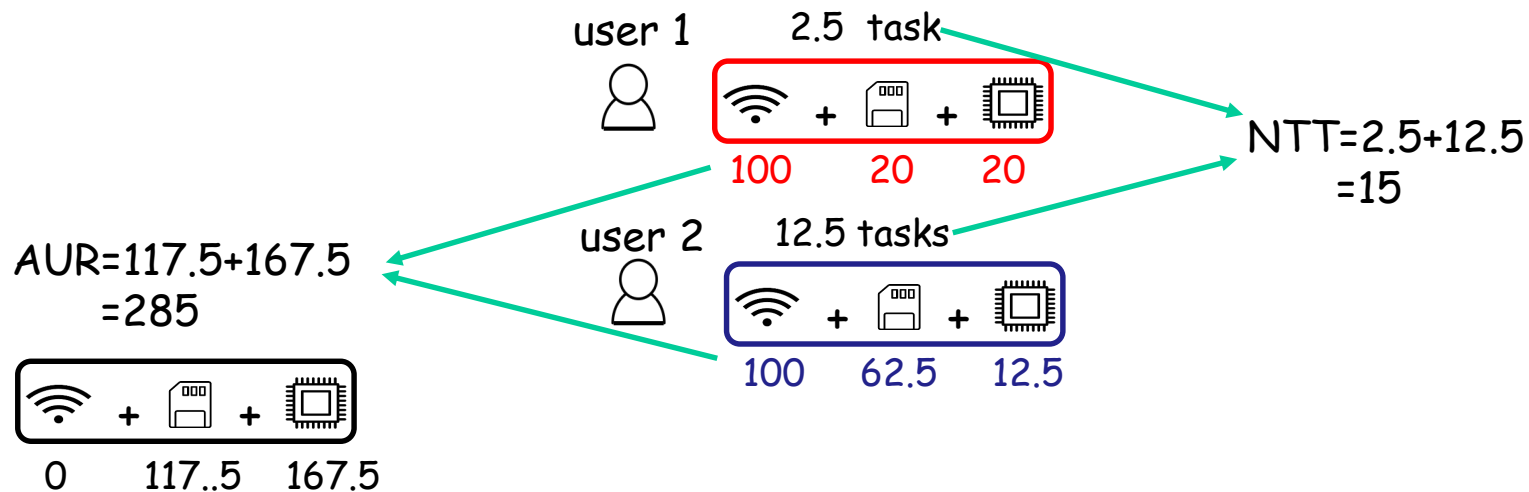
✦ Desirable Properties

sharing incentive, strategy proof, envy-free, Pareto efficient

Efficiency

✦ Two measurements

1. the number of total tasks completed (NTT)
2. the amount of unused resources (AUR)



3. 2-Dominant Resource Fairness

& Model

- ✧ r resources and n users
- ✧ Resource j 's capacity: C_j
- ✧ User i 's request vector: D_i
- ✧ User i 's final allocation vector: A_i

& Objective

- ✧ Design a new fairshare function for max-min fairness for efficiency improvement
- ✧ A fairshare function on
 - multiple resources (instead of dominant resource alone)
 - weighting factors among different resources

2-DF Allocation Mechanism

2-DF fairshare function

- ✦ 2-dominant share: $s_i = \varphi_i \cdot d_{i1} \cdot d_{i2}$ where
- ✦ φ_i : number of user i 's tasks
- ✦ $d_{i1} = \max \left\{ \frac{a_{ij}}{c_j} \right\}$: user i 's first dominant ratio among all resources j
- ✦ $d_{i2} = \max \left\{ \frac{a_{ij}}{c_j} \right\} - \{d_{i1}\}$: user i 's second dominant ratio among all resources j

Allocation mechanism

$$s_1 = \varphi_1 \cdot d_{11} \cdot d_{12} \\ = \varphi_1 / 125$$

$$s_2 = \varphi_2 \cdot d_{21} \cdot d_{22} \\ = \varphi_2 / 1000$$

maximize φ_1, φ_2

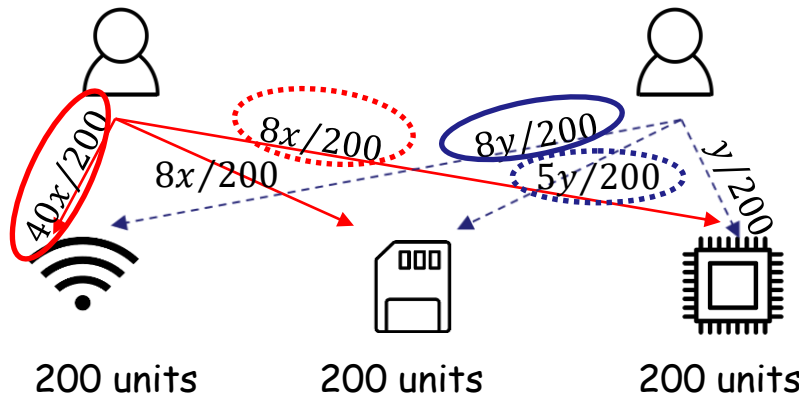
subject to

$$\begin{cases} \varphi_1 / 125 = \varphi_2 / 1000 \\ 40\varphi_1 + 8\varphi_2 \leq 200 \\ 8\varphi_1 + 5\varphi_2 \leq 200 \\ 8\varphi_1 + \varphi_2 \leq 20 \end{cases}$$

$$\varphi_1 = 2 \quad \varphi_2 = 15$$

user 1: φ_1 tasks

user 2: φ_2 tasks

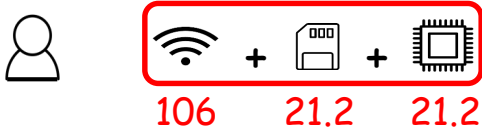


4. Properties and Extension (k-DF)

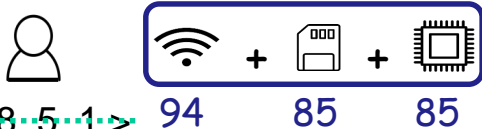
Properties

- ✦ Strategy proof
- ✦ Pareto-efficient

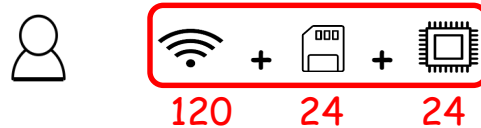
user 1 2.65 tasks



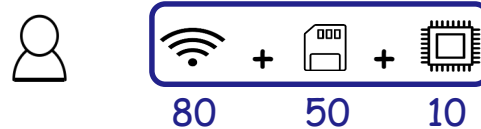
user 2 11.75 tasks



user 1 3 tasks



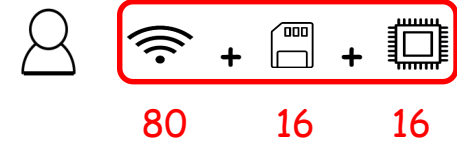
user 2 10 tasks



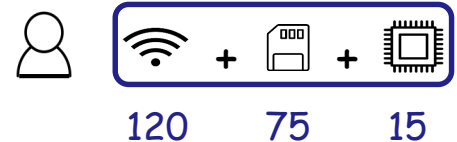
$\langle 8, 5, 1 \rangle$
 $\langle 10, 9, 9 \rangle$

2-DF Allocation

user 1 2 tasks



user 2 15 tasks



k-DF mechanism

- ✦ k-dominant share: $s_i = \varphi_i \prod_{l=1}^k w_{il} d_{il}$ where w_{il} is a weight
- ✦ consider k dimensions of resources

5. Experiment -- First Scenario

↳ Setting

- ✧ A data center with 3 resources and 3 users
- ✧ Resource capacity $\langle C, C, C \rangle$ where $C \in \{3, 5\}$
- ✧ User 1's request vector $\langle d_{11}, d_{12}, d_{13} \rangle$ where $d_{1i} \in [1, C]$
- ✧ User 2's request vector $\langle d_{21}, d_{22}, d_{23} \rangle$ where $d_{2j} \in [1, C]$
- ✧ User 3's request vector $\langle d_{31}, d_{32}, d_{33} \rangle$ where $d_{3k} \in [1, C]$

↳ Three comparison algorithms

- ✧ No Fairness Constraints (NFC)
- ✧ Dominant Resource Fairness (DRF)
- ✧ 2-Dominant Resource Fairness (2DF)

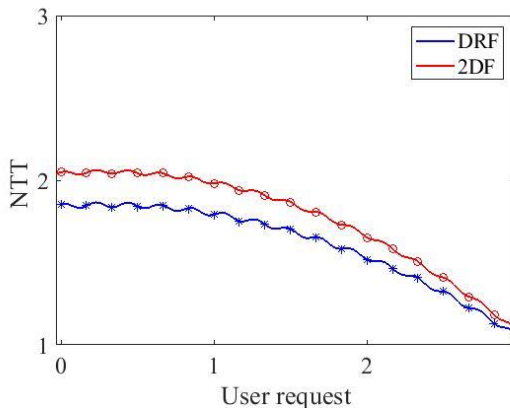
First Scenario

Efficiency -- NNT

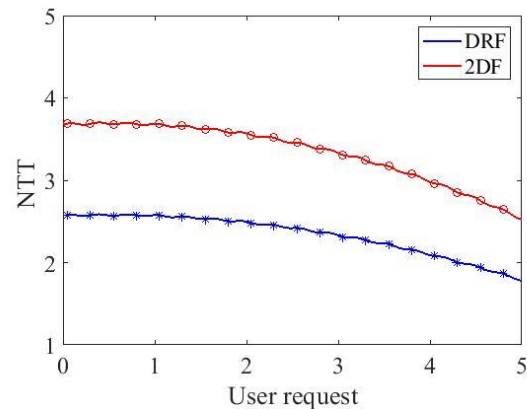
Average NTT under different capacities

Capacity	NFC	DRF	2DF
3	1.626	1.342	1.387
5	1.823	1.481	1.545

NNT with request vectors increasing



(a) Capacity of 3



(b) Capacity of 5

5. Experiment -- Second Scenario

Setting

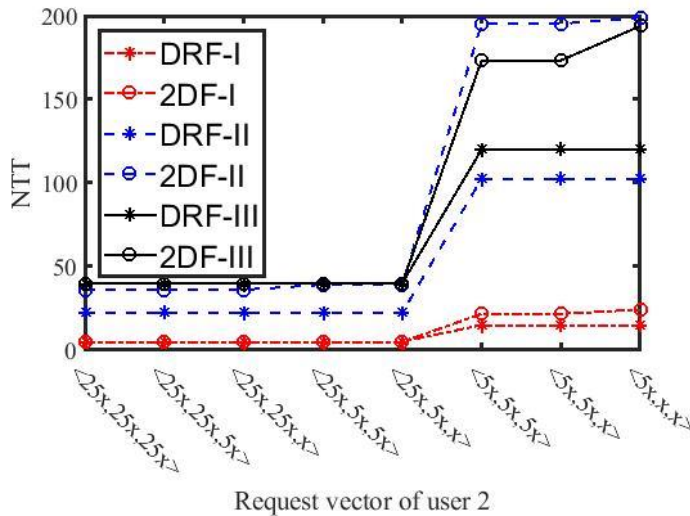
- ✧ A data center with 3 resources and 2 users
- ✧ Resource Capacity $\langle 1000, 1000, 1000 \rangle$
- ✧ Two user request types: heavy and light
 - a request $D_i = \langle d_{i1}, d_{i2}, d_{i3} \rangle$ is heavy if $\forall d_{ij} \in \{25x_1, 5x_1, x_1\}$
 - a request $D_i = \langle d_{i1}, d_{i2}, d_{i3} \rangle$ is light if $\forall d_{ij} \in \{25x_2, 5x_2, x_2\}$
 - where $x_1 \sim N(8, 0)$ and $x_2 \sim N(1, 0)$
- ✧ Three combinations of user request types

Combination	user 1	user 2
I	heavy	heavy
II	heavy	light
III	light	light

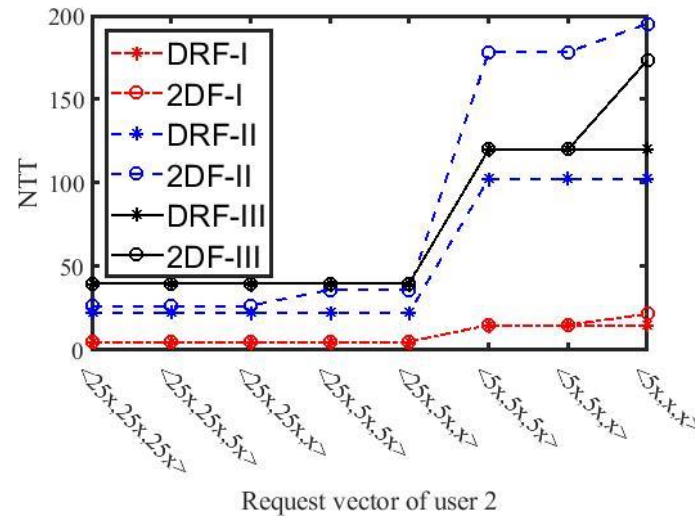
Second Scenario

Efficiency

- ✧ NNT : When user 1 has big tasks and user 2 has small tasks, the improvement of NNT is most obvious.
- ✧ Average increase: 45%

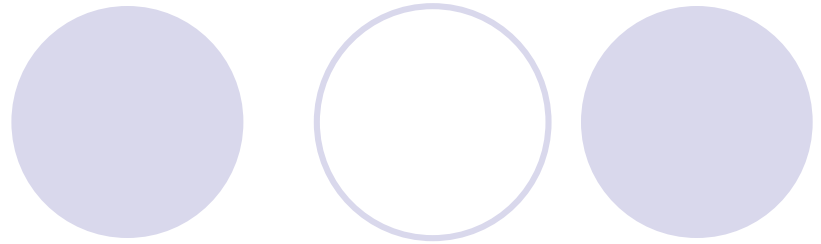


- (a) I: user 1 $\langle 25x_1, 25x_1, 25x_1 \rangle$
 II: user 1 $\langle 25x_1, 25x_1, 25x_1 \rangle$
 III: user 1 $\langle 25x_2, 25x_2, 25x_2 \rangle$



- (b) I: user 1 $\langle 25x_1, 5x_1, x_1 \rangle$
 II: user 1 $\langle 25x_1, 5x_1, x_1 \rangle$
 III: user 1 $\langle 25x_2, 5x_2, x_2 \rangle$

Second Scenario



Efficiency

✧ AUR : 2DF consumes **less resources** while yields **more tasks**

Average AUR	DRF	2DF
I	418	654
II	418	802
III	418	654

Average AUR	DRF	2DF
I	1218	1271
II	1218	998
III	1218	1271

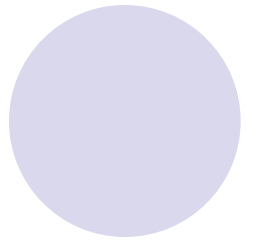
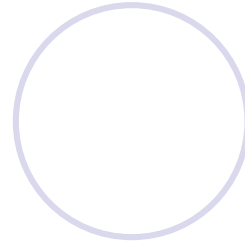
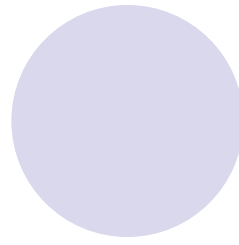
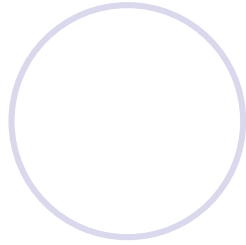
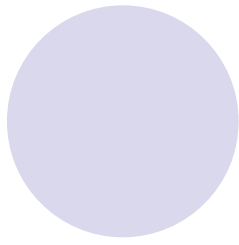
(a) I: user 1 $\langle 25x_1, 25x_1, 25x_1 \rangle$
II: user 1 $\langle 25x_1, 25x_1, 25x_1 \rangle$
III: user 1 $\langle 25x_2, 25x_2, 25x_2 \rangle$

(b) I: user 1 $\langle 25x_1, 5x_1, x_1 \rangle$
II: user 1 $\langle 25x_1, 5x_1, x_1 \rangle$
III: user 1 $\langle 25x_2, 5x_2, x_2 \rangle$



6. Conclusion

- ↳ DRF suffers from serious fairness concerns without utility guarantees
- ↳ 2-DF seeks to balance fairness and efficiency
 - ✧ Extension from 2-DF to k -DF
 - ✧ strategy proof and Pareto efficient



Q & A