



An Online Approach for DNN Model Caching and Processor Allocation in Edge Computing

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1. Introduction

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- 2. Contributions
- 3. Problem Formulation
- 4. Algorithm Design
- 5. Evaluation
- 6. Conclusion







1. Introduction



- Edge Computing
 - object detection
 - virtual reality
 - intelligent cameras



- Deep Neural Networks (DNN) inference
 - VGG, ResNet
- Cache DNN models on the edge brings benefits
 - efficiency
 - privacy
 - security







- Motivation
 - The difference between cloud and edge
 - Cloud: high computing capacity, long transmission delay
 - Edge: short transmission delay, more caching DNN cost
 - Computing is fine-grained on the edge
 - The number of processors assigned to each service will affect the service delay
 - The user request distribution for mobility
 - The user connection information that counts the history on each server can be obtained





- Achieve the trade-off between user perception delay and energy consumption cost
 - consider DNN model caching and processor allocation in EC
 - NP-Complete
- Propose a novel online algorithm called DMCPA-GS-Online
 - leverages the Lyapunov framework
 - expanded edition of Gibbs Sampling
 - near-optimal
- Evaluate the performance of DMCPA-GS-Online
 - a trace dataset from the real world

3. Problem Formulation



- A user request will encounter two caching hit conditions:
 - Edge-Hit: directly processed by the edge
 - Cloud-Hit: forwarded to the cloud server



Fig. 1. An overview of our problem.

 $x_{i,i}(t) \in \{0,1\}$

cache decision vector

 $y_{i,i}(t) \in \mathbb{Z}$

processor allocation vector







- Energy Consumption Cost
 - The consumption cost to maintain processors: $e_{i,j}$
 - The initialization cost for DNN model: a_j
- Overall

$$T^{e}(t) = \sum_{S_{i} \in S} \sum_{M_{j} \in M} e_{i,j} y_{i,j} + a_{j} \max\{x_{i,j}(t) - x_{i,j}(t-1), 0\}$$

The Total Cost

 $T(t) = \alpha T^{d}(t) + \beta T^{e}(t)$

3. Problem Formulation



$$\mathcal{P}^{0}: \min_{\forall t, x(t), y(t)} \lim_{T \to \infty} \frac{1}{T} \sum_{t=0}^{T-1} \mathbb{E}[\mathcal{T}(t)]$$

s.t. $C_{1}: \mathcal{T}(t) = \alpha \mathcal{T}^{d}(t) + \beta \mathcal{T}^{e}(t),$
 $C_{2}: \sum_{M_{j} \in \mathbb{M}} x_{i,j}(t) w_{j} \leq W_{i}, \forall S_{i} \in \mathbb{S},$
 $C_{3}: \sum_{M_{j} \in \mathbb{M}} y_{i,j}(t) \leq \phi_{i}, \forall S_{i} \in \mathbb{S},$
 $C_{4}: x_{i,j}(t) \in \{0, 1\}, \forall S_{i} \in \mathbb{S}, \forall M_{j} \in \mathbb{M},$

$$C_{5}: \quad y_{i,j}(t) \in \mathbb{Z}, \forall S_{i} \in \mathbb{S}, \forall M_{j} \in \mathbb{M},$$
$$C_{5}: \quad \min\{u_{i}(t) \in \mathbb{Q}, \forall S_{i} \in \mathbb{S}, \forall M_{j} \in \mathbb{M}, u_{i}(t) \in u_{i}(t)\} \in \mathbb{N},$$

$$C_6: \min\{y_{i,j}(t), 0\} \le x_{i,j}(t) \le y_{i,j}(t),$$

$$C_7: \lim_{T \to \infty} \frac{1}{T} \sum_{t=0}^{T-1} \mathbb{E}[L_j^P(t)] \le D_j, \forall M_j \in \mathbb{M}.$$



This is a long-term average optimization problem









Define the queues in each time slot:

$$Q_{j}(t+1) = \max \{Q_{j}(t) + L_{j}^{P}(t) - D_{j}, 0\}$$

- Define the quadratic Lyapunov function: $L(\Theta(t)) = \frac{1}{2} \sum_{j} Q_{j}(t)^{2}, \quad \text{let } \Theta(t) = [Q_{1}(t), \dots, Q_{N_{m}}(t)]$
- Define the Lyapunov drift:

 $\Delta(\Theta(t)) = L(\Theta(t+1)) - L(\Theta(t))$

Using Lyapunov drift, the drift-plus-penalty algorithm can be used, for the fact:

$$\Delta(\Theta(t)) \le B + \sum_{j} Q_{j}(t) (L_{j}^{P,Q}(t) - D_{j})$$

4. Algorithm Design

 The origin problem can be converted into a series of subproblems as:

$$\mathcal{P}^3: \min_{\forall t, x(t), y(t)} \mathbb{E}[B + V \cdot \mathcal{T}(t) + \sum_{j} Q_j(t) (L_j^P(t) - D_j) |\Theta(t)]$$

$$(C_1) - (C_6),$$

one time slot problem

s.t.

Algorithm 1: The DMCPA-GS-Online Algorithm Input: $Q_j(0) \leftarrow 0, x^{prev}(0) \leftarrow 0$ for t = 0 to T do receive $d_k^{input}(t)$ from environment; update current distribution P from environment; get $x^*(t), y^*(t)$ by solving \mathcal{P}^3 using Alg. 2; $x^{prev}(t+1) \leftarrow x^*(t)$; for $M_j \in \mathbb{M}$ do $| Q_j(t+1) \leftarrow \max\{Q_j(t) + L_j^P(t) - D_j, 0\}$; end end

The algorithm for the subproblem









provec

- The expanded edition of Gibbs Sampling Algorithm for the subproblem
 - 1. initialize y randomly at the beginning
 - 2. compute the optimal caching strategy x based on y
 - 3. jump from (x, y) to (x', y') with the probability:

$$\Pr = \frac{1}{1 + e^{(T' - T)/\omega}}$$

The probability of getting the global optimal value close to 1 when w is close to 0

4. repeat jump until converging





- Evaluation Setup
 - 5 edge servers
 - at least 200 users
 - P follows the trace of a dataset from the real world
 - d_k^{input} follows Poisson distribution with expectation 100
 - ϕ_i is set to 20
 - W_i follows N(10,3) , W_i is 300
 - set $\alpha = \beta = 0.5$





Motivation



Dataset from the real world



(a) The probability variation.







Fig. 5. The impact of different parameters for DMCPA-GS.



- Cloud-Only
- Edge-Average
- Edge-Random



Fig. 6. The performance of different algorithms.





- We consider the DNN Model Caching and Processor Allocation problem.
- Our goal is to minimize user perception delay and energy consumption with careful model caching and processor allocation strategy.
- We formulate it as an Integer Nonlinear Program and the novel online algorithm DMCPA-GS-Online is proposed without future information.
- Experiments based on trace dataset from the real world demonstrate our algorithm outperforms other baselines.





Thanks! Q&A