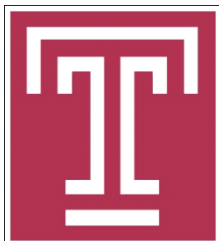
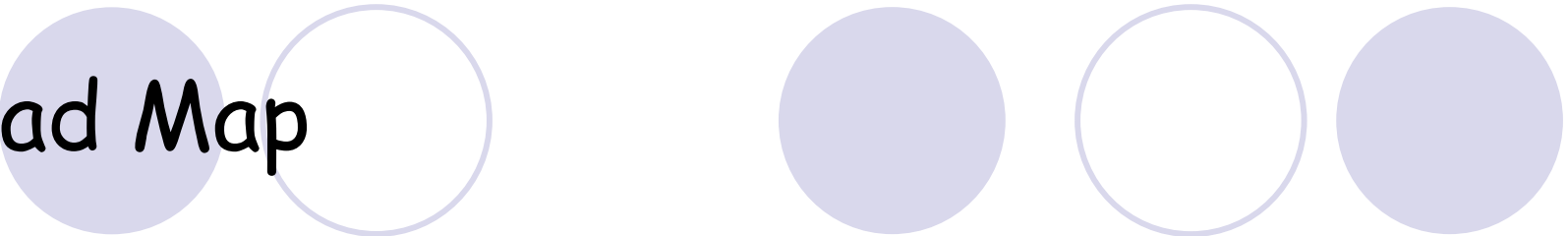


Collaborative Mobile Charging and Coverage in WSNs

Jie Wu

Computer and Information Sciences
Temple University



A decorative graphic at the top of the slide consists of two groups of three circles. The first group on the left has a solid light purple circle on the left, a white circle with a light purple outline in the middle, and another solid light purple circle on the right. The second group on the right has a solid light purple circle on the left, a white circle with a light purple outline in the middle, and another solid light purple circle on the right.

Road Map

1. Introduction
2. Mobile Chargers
3. State of the Arts
4. Challenges
5. Collaborative Coverage & Charging
6. Conclusions



1. Introduction

Need for basic research

John F. Kennedy

- ... progress in technology depends on progress in theory ...
The vitality of a scientific community springs from its passion to answer science's most fundamental questions.

Ronald Reagan

- ... although basic research does not begin with a particular practical goal ..., it ends up being one of most practical things government does.

My Two Cents

- How to select a research problem

- Simple definition
- Elegant solution
- Room for imagination

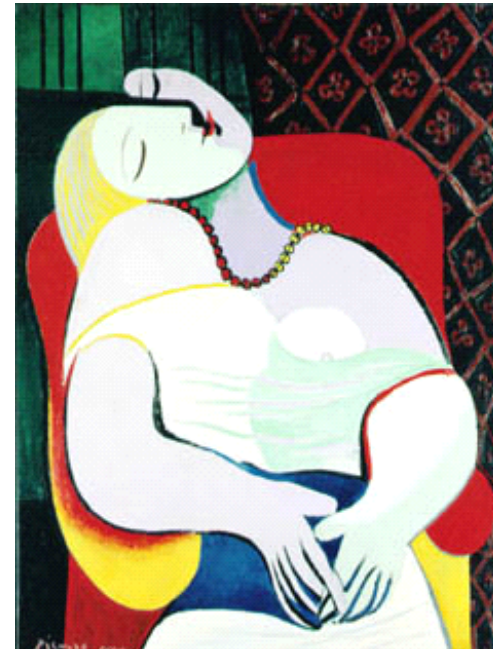


Blue Nude II₄

Picasso & Matisse

- Know how to make appropriate **abstractions** - ask the right questions
- Many CS students use excessive amounts of math to explain simple things
- The Art of Living, Time, Sept. 23, 2012
Senior people can be creative without worry the "utility" of their work

Le Rêve (the Dream)



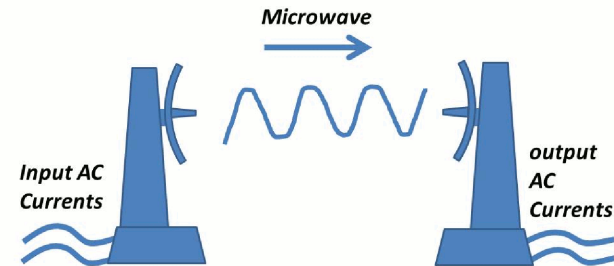


ENERGY: A Special Utility

- Limited lifetime of battery-powered WSNs
- Possible solutions
 - Energy conservation
 - Cannot compensate for energy depletion
 - Energy harvesting (or scavenging)
 - Unstable, unpredictable, uncontrollable ...
 - Sensor reclamation
 - Costly, impractical (deep ocean, bridge surface ...)

2. Mobile Chargers

- The enabling technology
 - Wireless energy transfer (Kurs '07)
 - Wireless Power Consortium

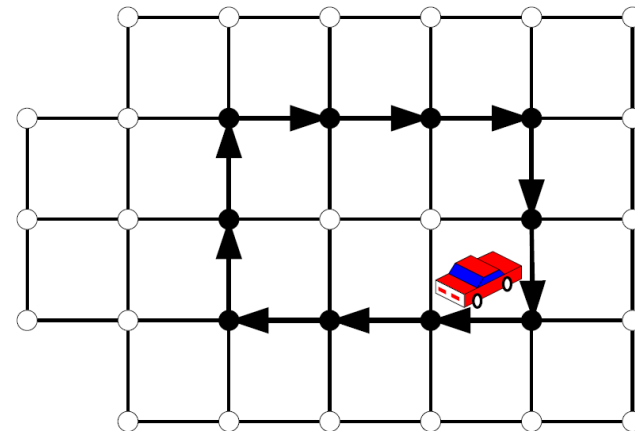


- Mobile chargers (MC)
 - MC moves from one location to another for wireless charging
 - Extended from **mobile sink** in WSNs and **ferry** in DTNs
 - Energy consumption
 - The movement of MC
 - The energy charging process

3. State of the Art

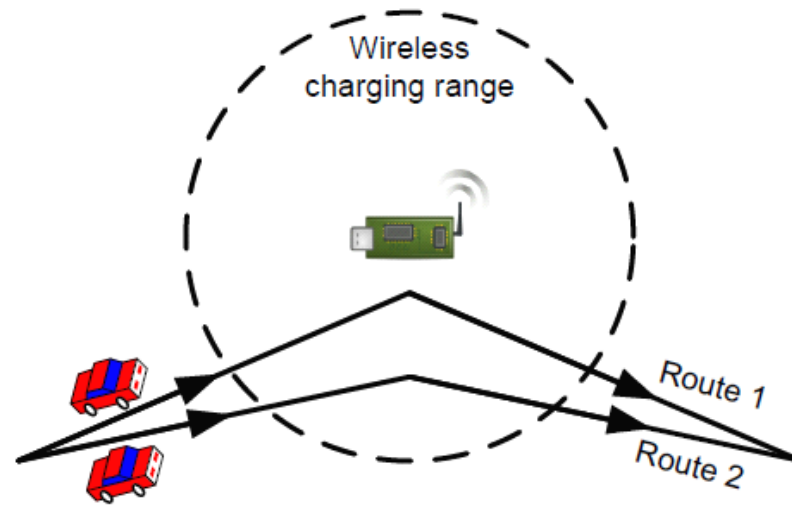
- Traveling-Salesmen Problem (TSP)
 - A minimum cost tour of n cities: the salesman travels from an origin city, visits each city exactly one time, then returns to the origin
- Covering Salesman Problem (CSP, Ohio State '89)
 - The least cost tour of a subset of cities such that every city not on the tour is within some predetermined covering distance

- Extended CSP
 - Connected dominating set (FAU '99)
 - Qi-ferry (UDelaware '13)



Charging Efficiency

- Location of charging



- Bundle and rotation (Kurs '10)
 - Charging multiple devices that are clustered together

Mobile Sinks and Chargers

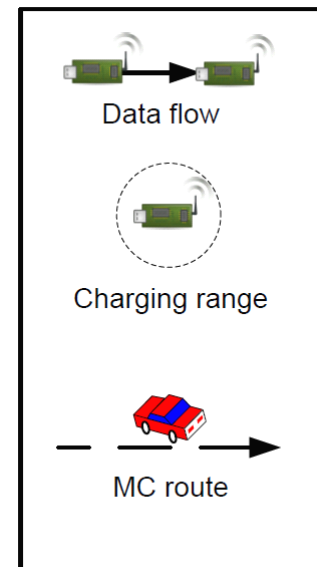
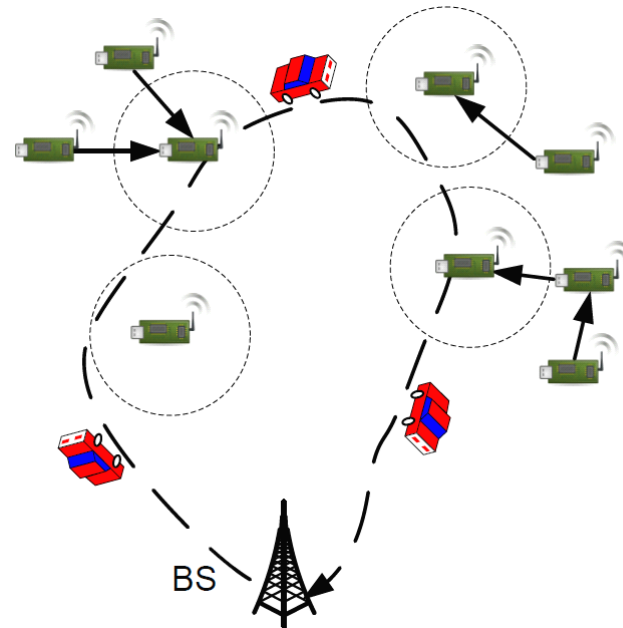
- Local trees

- Data collections at all roots
- Periodic charging to all sensors

- Base station (BS)

- Objectives

- Long vocation at BS (VT '11-13)
- Energy efficiency with deadline (Stony Brook '13)



4. Challenges



- Most existing methods

- An MC is fast enough to charge all sensors in a cycle
- An MC has sufficient energy to replenish an entire WSN (and return to BS)

- Collaborative approach using multiple MCs

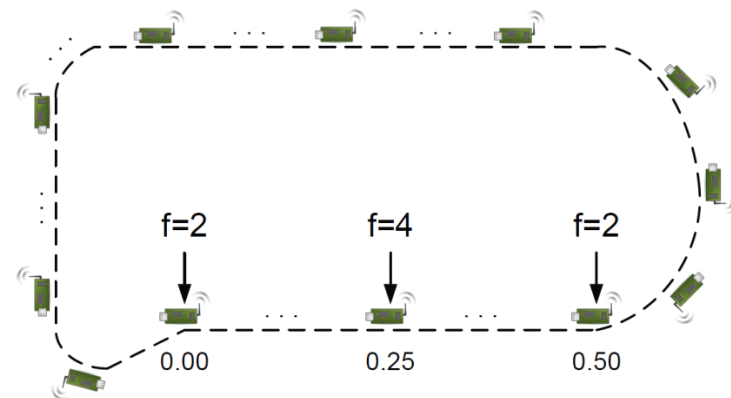
- Problem 1: MCs with unrestricted capacity but limitations on speed
- Problem 2: MCs with limited capacity and speed, and have to return to BS

5. Collaborative Coverage & Charging

- **Problem 1:** Determine the minimum number of MCs (unrestricted capacity but limitations on speed) to cover a line/ring of sensors with uniform/non-uniform recharge frequencies

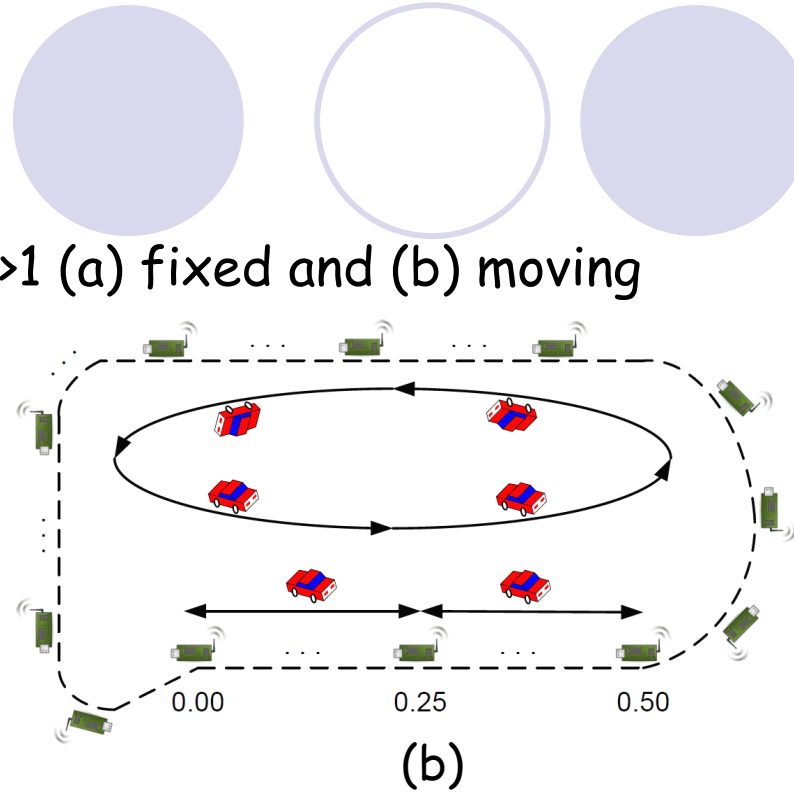
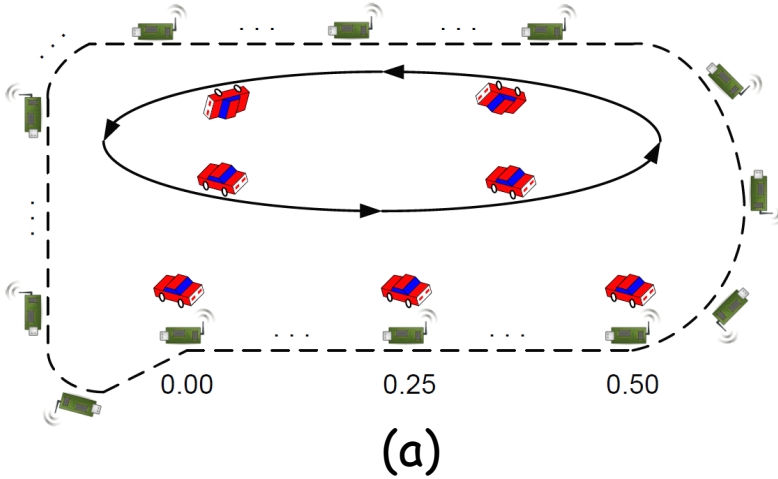
- **A toy example**

- A circle track with circumference 3.75 densely covered with sensors with frequency $f=1$ for recharge
- A sensor with $f=2$ at 0 and 0.5
- A sensor with $f=4$ at 0.25
(MC's max speed is 1)
- What is the trajectory planning of these MCs?

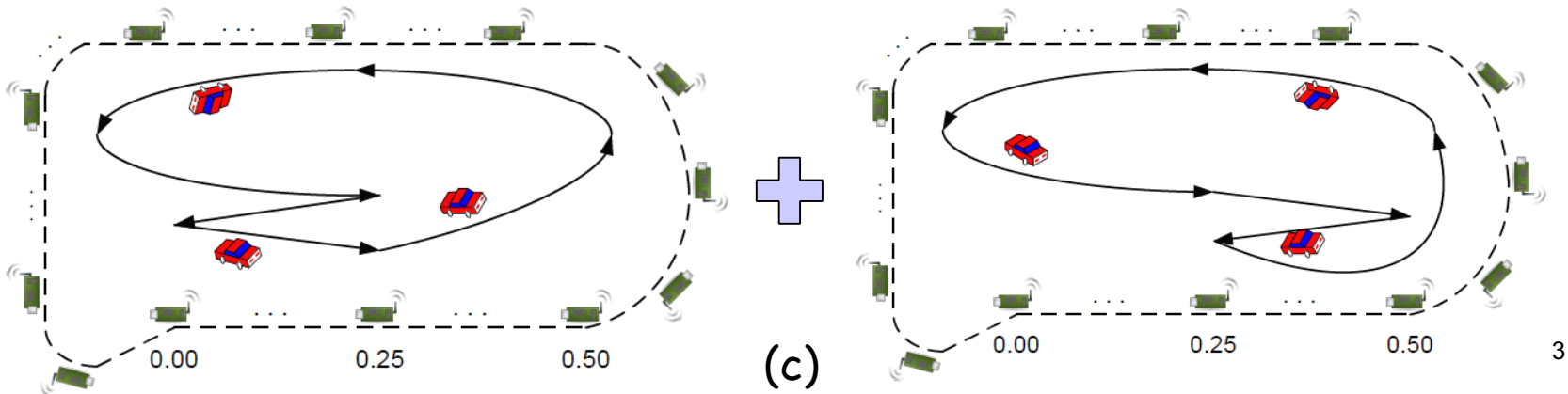


Possible Solutions

- Assigning cars for sensors with $f > 1$ (a) fixed and (b) moving

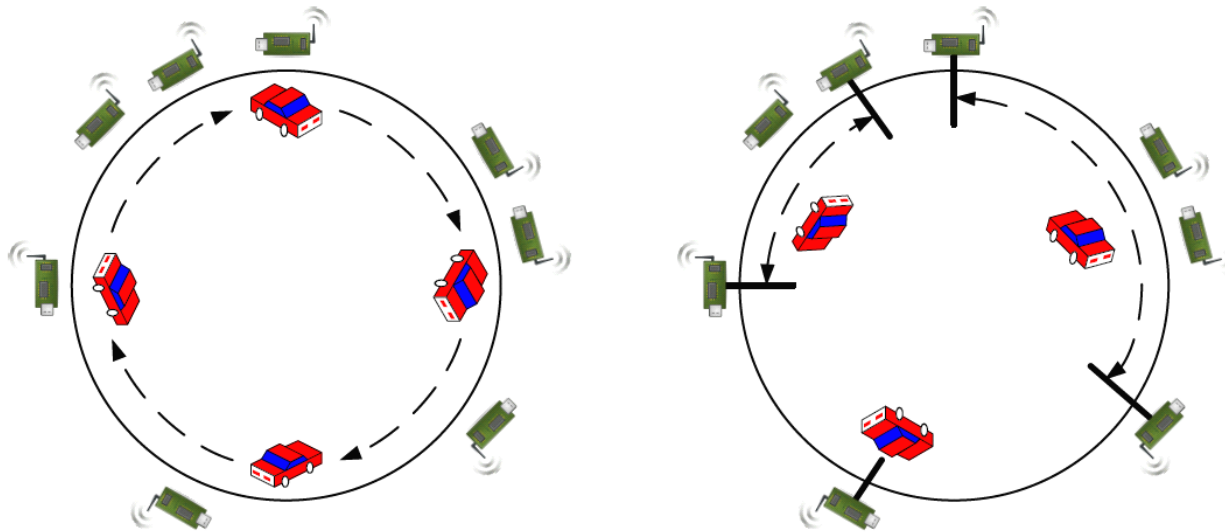


- Combining odd and even car circulations (c)



Optimal Solution (uniform frequency)

- M_1 : There are C_1 MCs moving continuously around the circle
- M_2 : There are C_2 MCs moving inside the fixed interval of length $\frac{1}{2}$ so that all sensors are covered
- **Combined method**: It is either M_1 or M_2 , so $C = \min \{C_1, C_2\}$





Properties

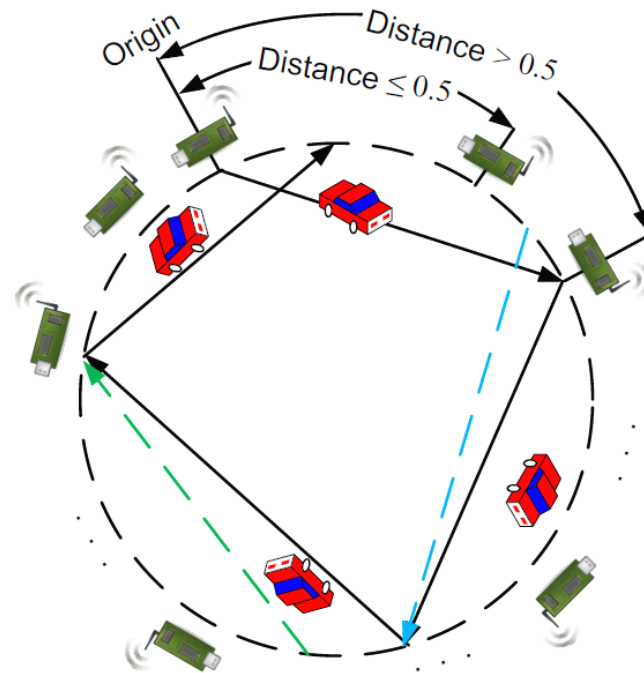
- **Theorem 1:** The combined method is optimal in terms of the minimum number of MCs used
- Scheduling
 - Find an appropriate **breakpoint** to convert a circle to a line; M_2 in the optimal solution is then followed
 - A **linear solution** is used to determine the breakpoint

Linear Solution

- Directed Interval Graph

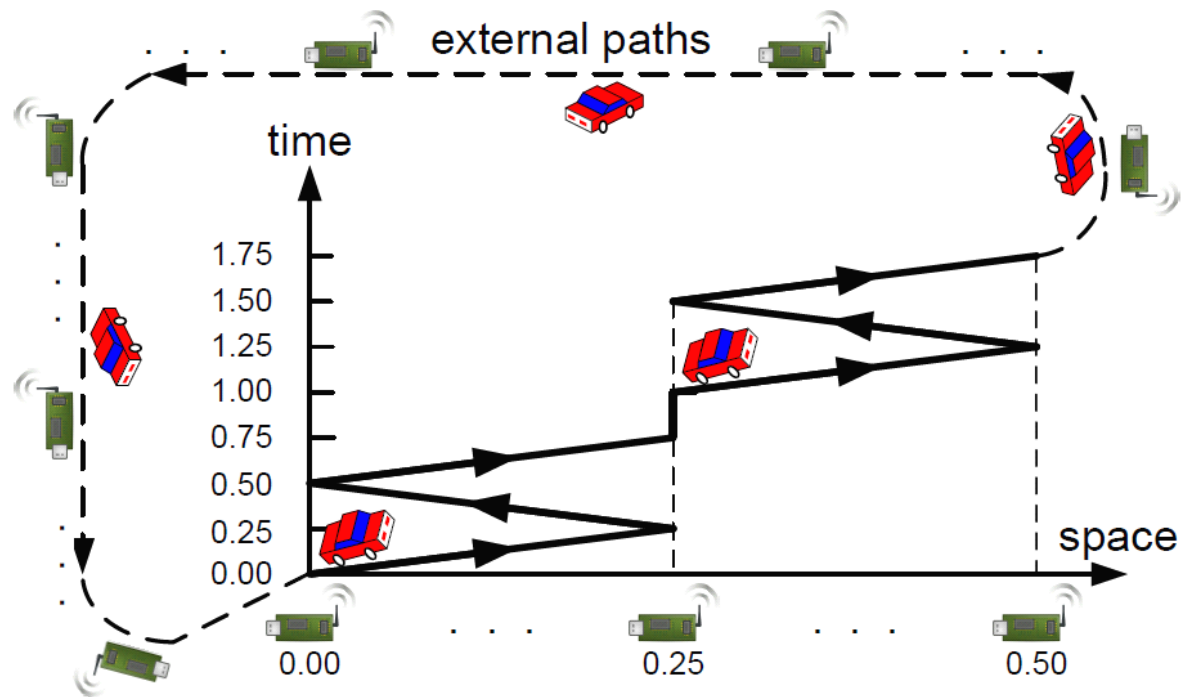
- Each directed link points from the start to the end of an interval (i.e., the first sensor beyond distance 0.5)

- The number of intervals in two solutions differ by one
- Each sensor has one outgoing and multiple incoming links
- The process stops when a path with fewer or more intervals is found or all sensors (with their outgoing links) are examined



Solution to the Toy Example

- 5 cars only, including a stop at 0.25 for 15 seconds



- **Challenges:** time-space scheduling, plus speed selection

Greedy Solution (non-uniform frequency)

- Coverage of sensors with non-uniform frequencies

serve($x_1, \dots, x_n; f_1, \dots, f_n$):

When $n \neq 0$, generate an MC that goes back and forth
as

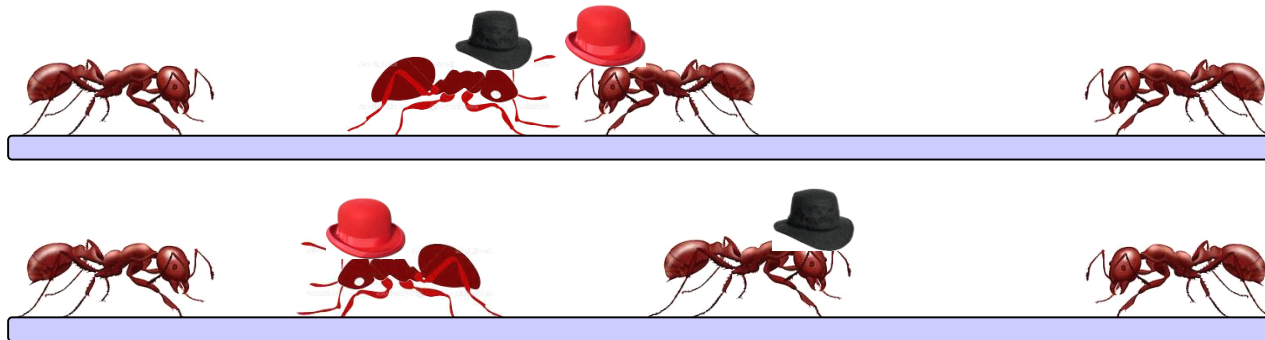
far as possible at full speed (covering x_1, \dots, x_{i-1});

serve($x_i, \dots, x_n; f_i, \dots, f_n$)

- **Theorem 2:** The greedy solution is within a factor of 2 of the optimal solution

The Ant Problem: An Inspiration

- Ant Problem, Comm. of ACM, March 2013
 - Ant Alice and her friends always march at 1 cm/sec in whichever direction they are facing, and reverse directions when they collide
 - Alice stays in the middle of 25 ants on a 1 meter-long stick
 - How long must we wait before we are sure Alice has fallen off the stick?

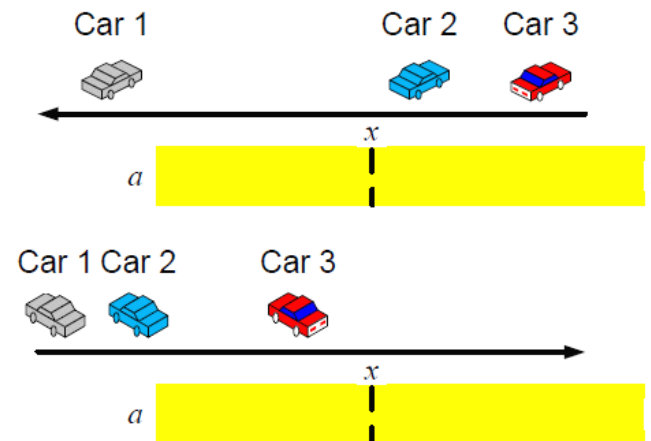
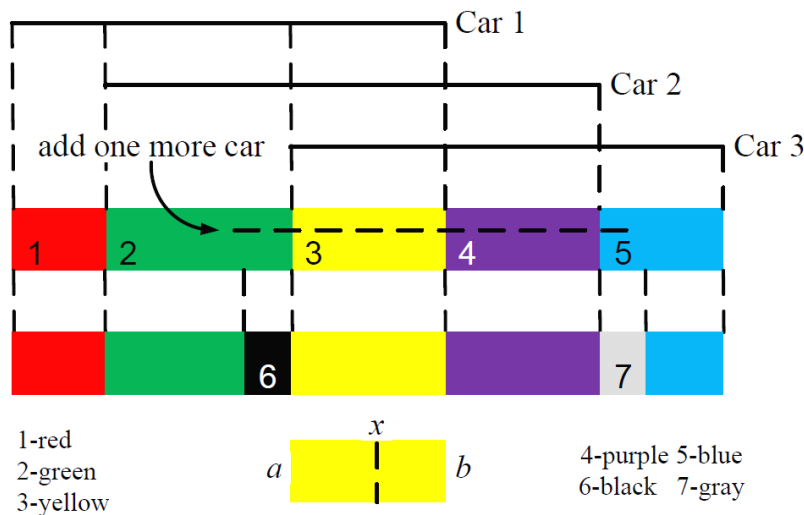


- Exchange "hats" when two ants collide

Proof of Theorem 2



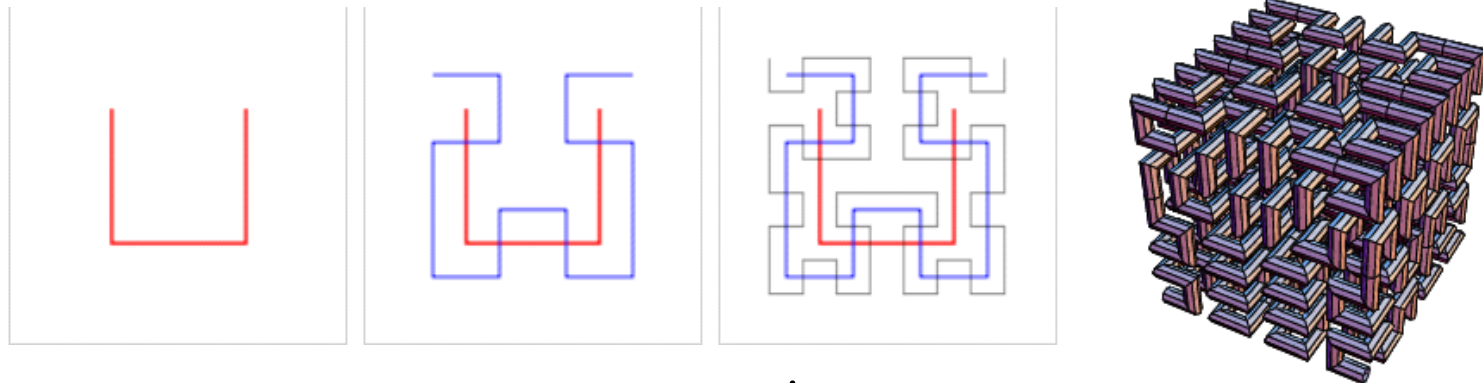
- Two cars never meet or pass each other
- Partition the line into $2k-1$ sub-regions based on different car coverage (k is the optimal number of cars)
- Each sub-region can be served by one car at full speed
- One extra car is used when a circle is broken to a line



Imagination

- Hilbert curve for k-D

- Mapping from 2-D to 1-D for preserving locality fairly well



- Charging time: converting to distance

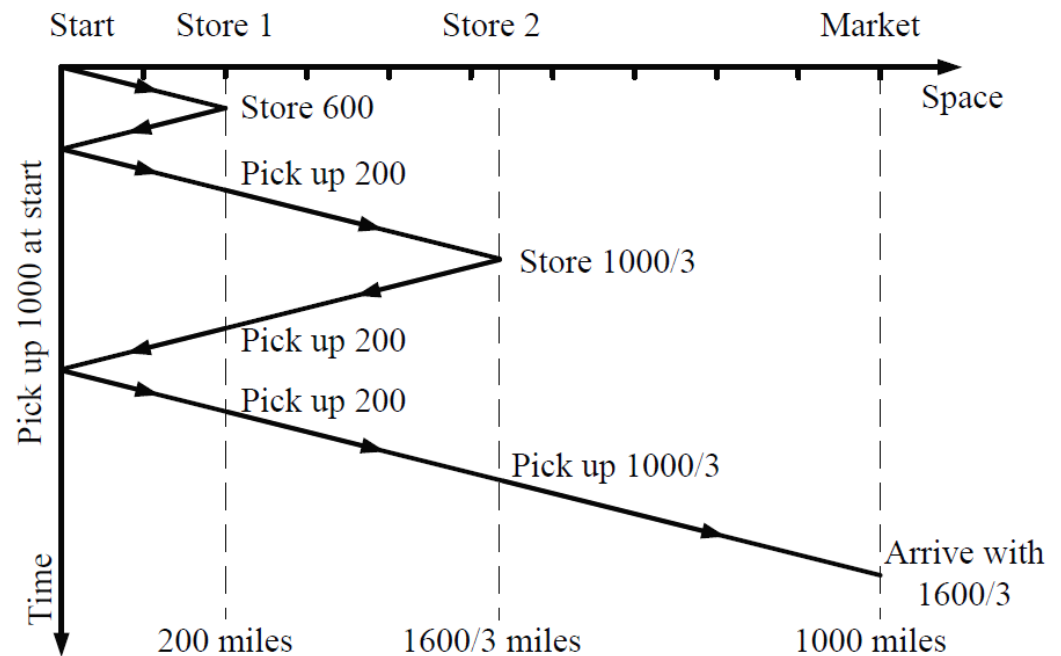
- Limited capacity: using cooperative charging

- BS to MC
- MC to MC

Bananas and a Hungry Camel

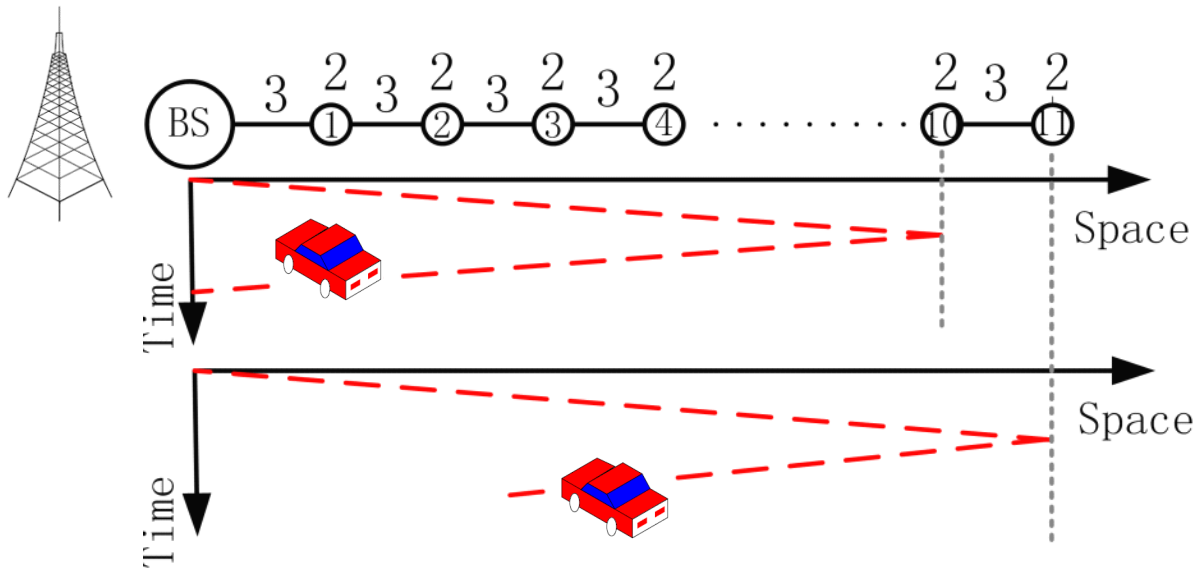
- A farmer grows 3,000 bananas to sell at market 1,000 miles away. He can get there only by means of a camel. This camel can carry a maximum of 1,000 bananas at a time, but needs to eat a banana to refuel for every mile that he walks

What is the maximum number of bananas that the farmer can get to market?



Charging a Line (with limited capacity)

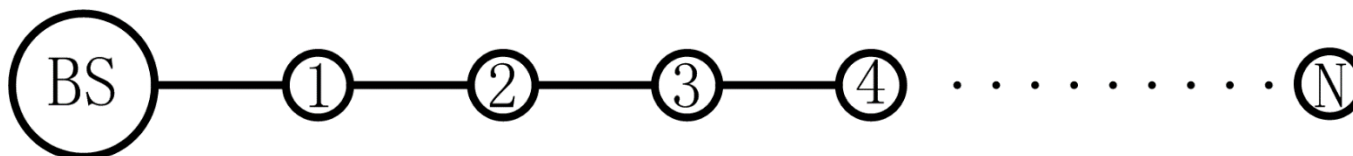
- Charge battery capacity: 80J; charger cost: 3J per unit traveling distance; sensor battery capacity: 2J



- One MC cannot charge more than 10 sensors
- Even a dedicated MC cannot charge the 14th sensor, since $14 * 3 + 2 + 14 * 3 = 86 > 80$

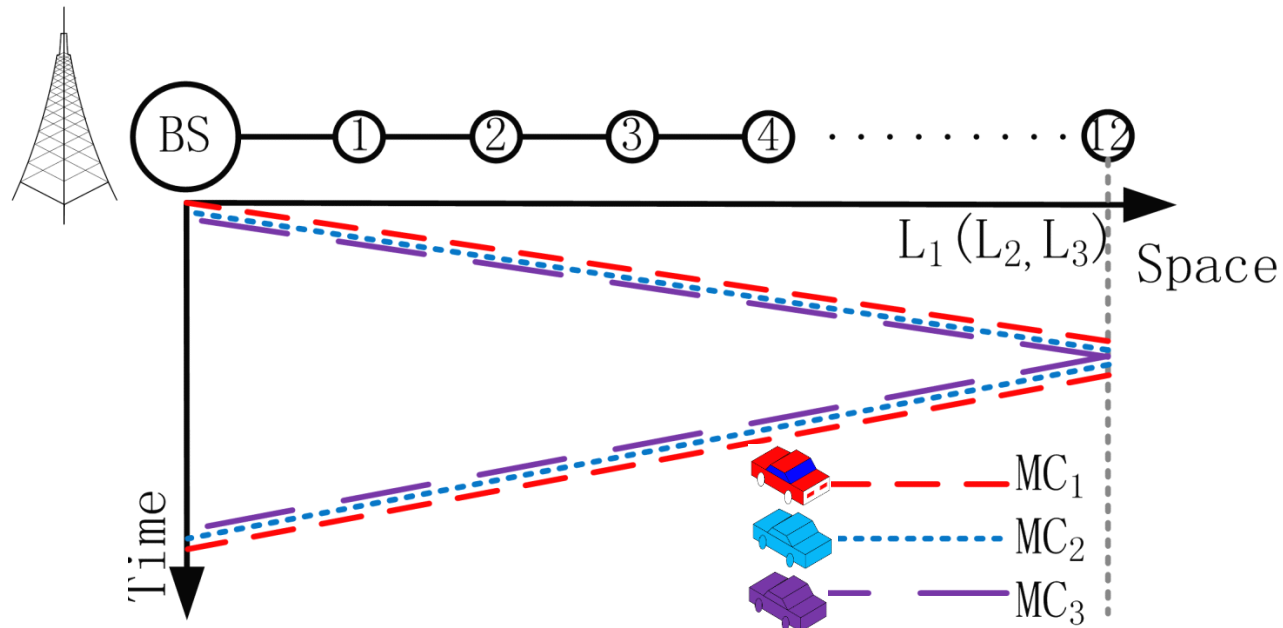
Problem Description

- **Problem 2** (IEEE MASS'12): Given k MCs with limited capacity, determine the furthest sensor they can recharge while still being able to go back to the BS
- WSN
 - N sensors, unit distance apart, along a line
 - Battery capacity for each sensor : b
 - Energy consumption rate for each sensor: r
- MC
 - Battery capacity: B
 - Energy consumption rate due to travelling: c



Motivation Example (1)

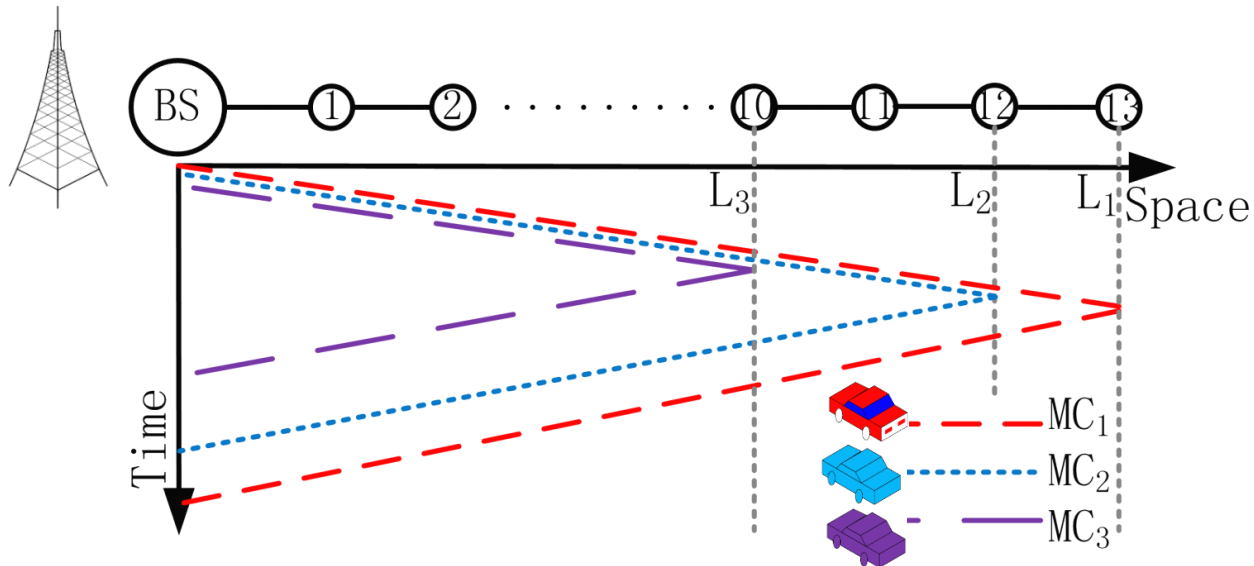
$B=80\text{J}$, $b=2\text{J}$, $c=3\text{J/m}$, $K=3$ MCs



- **Scheme I:** (equal-charge) each MC charges a sensor b/M J
- **Conclusion:** covers **12** sensors, and max distance is $< B/2c$
(as each MC needs a round-trip traveling cost)

Motivational Example (2)

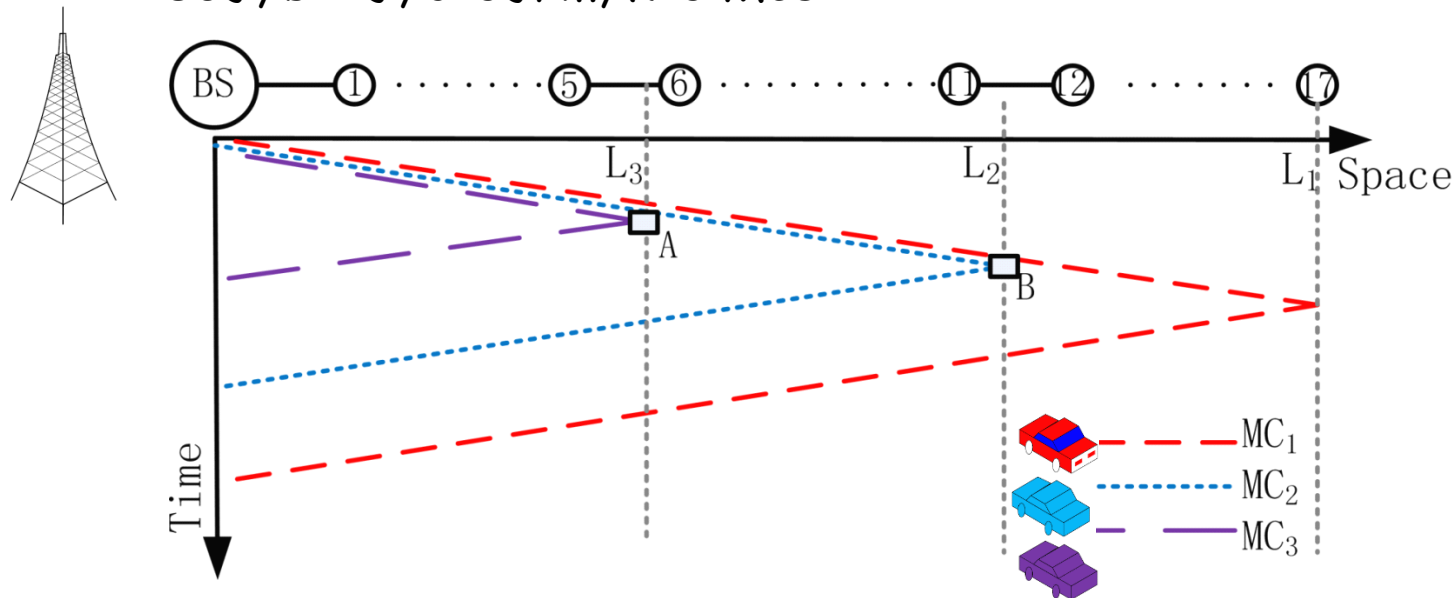
$B=80\text{J}$, $b=2\text{J}$, $c=3\text{J/m}$, $K=3$ MCs



- **Scheme II:** (one-to-one) each sensor is charged by one MC
- **Conclusion:** covers **13** sensors, and max distance is still $< B/2c$
(as the last MC still needs a round-trip traveling cost)

Motivational Example (3)

$B = 80\text{J}$, $b=2\text{J}$, $c=3\text{J/m}$, $K=3$ MCs



- **Scheme III:** (collaborative-one-to-one-charge) same as Scheme II, except each MC transfers some energy to other MCs at rendezvous points (A and B in the example)
- **Conclusion:** covers **17** sensors, and max distance is $< B/c$ (Last MC still needs a return trip without any charge)

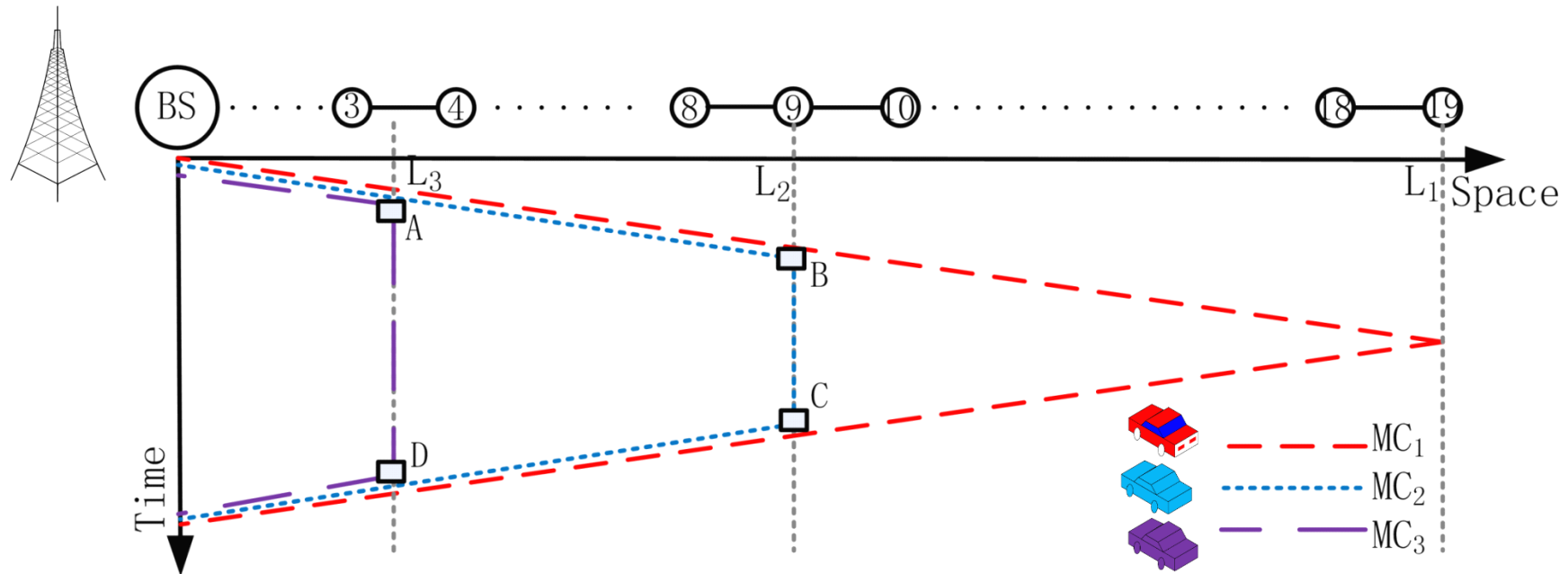
Details on Scheme III

- MC_i charges battery to all sensors between L_{i+1} and L_i
- MC_i ($1 \leq i \leq K$) transfers energy to $MC_{i-1}, MC_{i-2}, \dots, MC_1$ to their full capacity at L_i
- Each MC_i has just enough energy to return to the BS



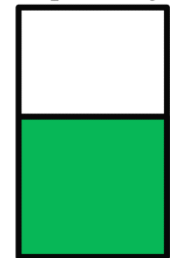
Motivational Example (4): Global Coverage

$B = 80\text{J}$, $b=2\text{J}$, $c=3\text{J/m}$, $K=3$ MCs



- "Push": limit as few chargers as possible to go forward
- "Wait": efficient use of battery "room" through *two charges*
- Conclusion: covers **19** sensors, and max distance is ∞

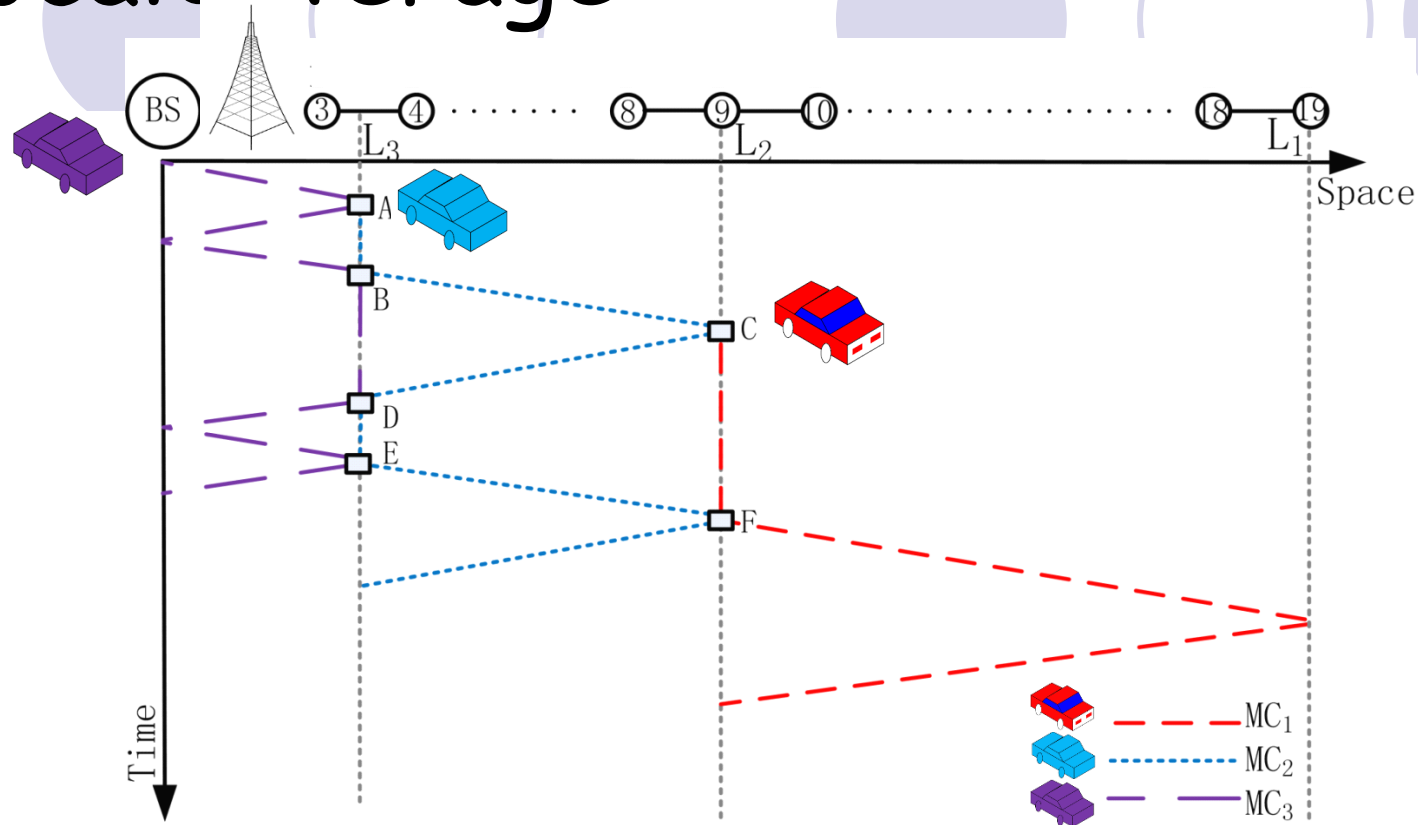
Capacity



Details on GlobalCoverage

- MC_i charges battery to all sensors between L_{i+1} and L_i
- MC_i ($1 \leq i \leq K$) transfers energy to MC_{i-1}, \dots, MC_1 to their full capacity at L_i
- MC_i waits at L_i , while all other MCs keep moving forward
- After $MC_i, MC_{i-1}, \dots, MC_1$ return to L_i , MC_i evenly balances energy among them (including itself)
- Each $MC_i, MC_{i-1}, \dots, MC_1$ has just enough energy to return to L_{i+1}

LocalCoverage



- Each MC moves and charges (is charged) between two adjacent rendezvous points
- Imagination: MC_i of LocalCoverage "simulates" $MC_i, MC_{i-1}, \dots, MC_1$ of GlobalCoverage

Properties

- **Theorem 3 (Optimality):** GlobalCoverage has the maximum ratio of payload energy to overhead energy
- **Theorem 4 (Infinite Coverage):** GlobalCoverage can cover an infinite line

○ Summation of segment length ($L_i - L_{i+1}$)

$$\begin{aligned} \sum_{i=1}^K \frac{B}{2 \cdot c \cdot i + b} &> \sum_{i=i_0}^K \frac{B}{2 \cdot c \cdot i + b} \text{ (let } 2 \cdot c \cdot i_0 \geq b) \\ &> \sum_{i=i_0}^K \frac{B}{4 \cdot c \cdot i} = \frac{B}{4 \cdot c} \sum_{i=i_0}^K \frac{1}{i} \text{ (harmonic series)} \end{aligned}$$



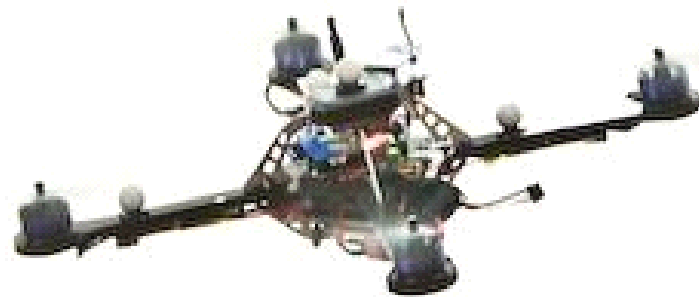
Imagination: extensions

- Simple extensions (while keeping optimality)
 - Non-uniform distance between adjacent sensors
 - Same algorithm
 - Smaller recharge cycle (than MC round-trip time)
 - Pipeline extension
- Complex extensions
 - Non-uniform frequency for recharging
 - Two- or higher-dimensional space

Imagination: applications

- Robotics

- Flying robots
- Google WiFi Balloon

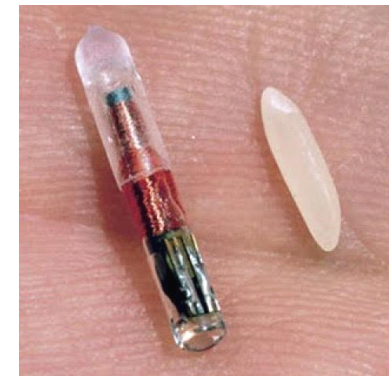


- WSNs

- Mobile sensor repairs with spares

- Passive RFID

- Energy transfer through readers



6. Conclusions

- Wireless energy transfer
- Collaborative mobile charging & coverage:
 - Unlimited capacity vs. limited capacity (with BS)
 - Charging type: BS-to-MC, MC-to-MC, and MC-to-Sensor
- Other extensions
 - Charging efficiency, MCs as mobile sinks for BS...
- **Simplicity + Elegance + Imagination = Beauty**

Acknowledgements

- Richard Beigel
- Sheng Zhang
- Huanyang Zheng

