Design and Implementation of a Strong Representation System for Network Policies

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Motivation

- Managing networking policies remains hard:
 - One has to fully understand the policy
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Can we provide a network policy management experience comparable to that on a database?

















P_R	dest	path	condition			
	1.2.3.4	l x	x=[ABC]∨x=[ADC			
	5.6.7.8	3 y	y=[ABC]∨y=[ADC]]		
	Ι	dest	path			
		1.2.3.4	4 [ABC]			
		1.2.3.4	4 [ADC]			
Represent		5.6.7.8	8 [ABC]			
		5.6.7.8	8 [ADC]			



P_R	dest	path	condition
	1.2.3.4	Х	x=[ABC]∨x=[ADC]
	5.6.7.8	У	y=[ABC]∨y=[ADC]
С	onditior	ble	

• Traffic balance policy





• Traffic balance policy

Balance traffic to					
1.2.0.4 and 0.0.7.0	<i>P</i> ₃	dest	path	flag	
		1.2.3.4	[ABC]	u	u = 1
		5.6.7.8	[ABC]	u	u ≠ 1
		1.2.3.4	[ADC]	V	v = 1
		5.6.7.8	[ADC]	V	v ≠ 1









• Static route and filter policy

P_1	dest	path	
	1.2.3.4	Х	x=[ABC]
	у	Z	y≠1.2.3.5

• Static route and filter policy

Assign a static route [ABC] to destination 1.2.3.4

P_1	dest	path	
	1.2.3.4	Х	x=[ABC]
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Assign a static route [ABC] to destination 1.2.3.4



Manipulating Network Policies

- Manipulate relational database table:
 - relational algebra : selection, projection, union, join, ...
 - implement SQL query on regular table

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- Manipulate relational database table:
 - relational algebra : selection, projection, union, join, ...
 - implement SQL query on regular table
- Our contribution:

Manipulating network policies by simply implementing SQL query on conditional table

- Static route and filter policy
- Traffic balance policy

P_1	dest	path		<i>P</i> ₃	dest	path	flag	
	1.2.3.4	Х	x=[ABC]		1.2.3.4	[ABC]	u	u = 1
	у	Z	y≠1.2.3.5∧y≠1.2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
					1.2.3.4	[ADC]	V	v = 1
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- Static route and filter policy
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P_1	dest	path			P_3	dest	path	flag	
	1.2.3.4	Х	x=[ABC]			1.2.3.4	[ABC]	u	u = 1
	У	Ζ	y≠1.2.3.5∧y≠1	.2.3.4		5.6.7.8	[ABC]	u	u ≠ 1
Can we generate a new policy that						1.2.3.4	[ADC]	V	v = 1
satis	sfies P_1 and	$d P_3 si$	multaneously?	P_1 join	P_3	5.6.7.8	[ADC]	V	v ≠ 1

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 Traffic balance policy

P_1	dest	path			P_3	dest	path	flag	
	1.2.3.4	Х	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	у	Ζ	y≠1.2.3.5∧y≠1.2	2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can we generate a new policy that						1.2.3.4	[ADC]	V	v = 1
satis	sfies P_1 and	d P_3 si	multaneously?	P_1 join	P_3	5.6.7.8	[ADC]	V	∨ ≠ 1

$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	Х	u	x=[ABC] ∧ u=1
	У	Z	u	y=5.6.7.8∧z=[ABC]∧u≠1
	У	Z	V	y=5.6.7.8∧z=[ADC]∧v≠1

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	1.2.3.4	Х	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	у	Ζ	y≠1.2.3.5∧y≠1.	2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we genera	ate a n	ew policy that			1.2.3.4	[ADC]	V	v = 1
satis	sfies P_1 and	$P_3 si$	multaneously?	P_1 join	P_3	5.6.7.8	[ADC]	V	V ≠ 1

$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	Х	U	x=[ABC] ∧ u=1
	У	Ζ	U	y=5.6.7.8∧z=[ABC]∧u≠1
	У	Ζ	V	y=5.6.7.8∧z=[ADC]∧v≠1
		Ι	dest	path
Valie			1.2.3.4	I [ABC]
state		ig	5.6.7.8	3 [ADC] 7

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 Traffic balance policy

P_1	dest	path			P_3	dest	path	flag	
	1.2.3.4	Х	x=[ABC]			1.2.3.4	[ABC]	U	u = 1
	у	Ζ	y≠1.2.3.5∧y≠1.	2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
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satis	sfies P_1 and	P_3 si	multaneously?	P_1 join	P_3	5.6.7.8	[ADC]	V	V ≠ 1

$P_1 \bowtie P_3$	dest	path	flag		
	1.2.3.4	Х	u	x=[ABC] \land u=1	
	У	Z	u	y=5.6.7.8∧z=[ABC]∧u≠1	
	У	Ζ	V	y=5.6.7.8∧z=[ADC]∧v≠1	
		Ι	dest	t path	
Valid forwarding			5.6.7.8 [ABC]		
state				7	













 The conventional conditional table is still restricted for the rich semantics of network policies

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$$\begin{array}{c|c|c|c|c|c|c|c|c|} P_2 & dest & path \ s(path) \\ \hline & x & y & z & l(y) \le z \end{array}$$

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Support Network Addressing

- Network addressing is a basic but critical feature to network area
- Allow sets in conditional table
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- Network addressing is a basic but critical feature to network area
- Allow sets in conditional table
 - i.e. accommodate variables and conditions
- Two methods:
 - Naive support for sets
 - Leverage SMT(Satisfiability Modulo Theories) solver for sets
Application in inter-domain routing

- We show the possibility of coordinating knowledge-driven policies from a single central point
- What about in realistic network scenarios like inter-domain routing?

We show an application - **Policy Exchange**

A Brief Picture of BGP



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A Brief Picture of BGP



Why policy exchange?

 BGP favors autonomy; Policies are made based on local preferences



Why policy exchange?

- BGP favors autonomy; Policies are made based on local preferences
- It may cause unwanted interaction between policies























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- Applications:
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Mimicking the behavior of the BGP speaker as it receives announcement from their neighbors

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- Applications:
 - Policy Exchange
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Benchmark

 Using MRT format RouteView BGP data - RIBs and UPDATEs from <u>route-view2.oregon-ix.net</u> to generate synthetic policies and realistic topologies

Performance - Relational Operators



 $\bowtie: P_1 \bowtie P_3$

 P_1 : static-route & filter policy P_3 : traffic balance policy

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Performance - Scalability



- SMT(Z3) takes 96% of the running time
- Our implementation can handle 10,000 policies in \leq 70 minutes

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Performance - More Policy Examples



 P_1 : static route policy P_3 : traffic balance policy P_4 : filter policy All: total time SMT: reasoning time

Processing time breakdown of poss(s, p)

• Primitive that validates a relational policy (p) on a given network state (s)

Performance - More Policy Examples



 P_1 : static route policy P_3 : traffic balance policy P_4 : filter policy All: total time SMT: reasoning time

- Primitive that validates a relational policy (p) on a given network state (s)
- Running time: P3 < P1 < P4

Performance - More Policy Examples



• Running time: P3 < P1 < P4

 SMT dominates the source of delay

Related Work

- Inter-domain Routing Protocols and Architectures
 - D-BGP and Trotsky: partial deployments of protocols, requires their co-existence on the global internet

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- Inter-domain Routing Protocols and Architectures
 - D-BGP and Trotsky: partial deployments of protocols, requires their co-existence on the global internet
- Declarative Networking:
 - We share database usage in networking
 - We introduce and implement a novel use of conditional tables

Thank you

http://ravel-net.org/

• Support full set of relational operators

Union: υ Selection: *σ* Difference: –

Projection: π

Rename: ρ








Other Relational Operations

Support full set of relational operators



- Static route and filter policy
 Traffic balance policy

P_1	dest	path			P_3	dest	path	flag	
	1.2.3.4	Х	x=[ABC]		1.2.3.4	[ABC]	U	u = 1	
	У	Ζ	y≠1.2.3.5∧y≠1	.2.3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we genera	ate a n	ew policy that		1.2.3.4	[ADC]	V	v = 1	
satis	fies P_1 and	d P_3 si	multaneously?	P_3	5.6.7.8	[ADC]	V	V ≠ 1	

$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	Х	u	x=[ABC] ∧ u=1
	1.2.3.4	Х	V	$x=[ABC] \land x=[ADC] \land v=1$
	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ABC] \land u = 1$
	У	Z	u	y≠1.2.3.5∧y≠1.2.3.4∧y=5.6.7.8∧z=[ABC]∧u≠1
	У	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ADC] \land v = 1$
	У	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ADC] \land v \neq 1$
	1.2.3.4 y y y y y	X Z Z Z Z	v u u v v	$x=[ABC] \land x=[ADC] \land v=1$ y=1.2.3.5 \y=1.2.3.4 \y=1.2.3.4 \z=[ABC] \u y=1.2.3.5 \y=1.2.3.4 \y=5.6.7.8 \z=[ABC] \u y=1.2.3.5 \y=1.2.3.4 \y=1.2.3.4 \z=[ADC] \u y=1.2.3.5 \y=1.2.3.4 \y=5.6.7.8 \z=[ADC] \v

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	1.2.3.4	Х	V	$x=[ABC] \land x=[ADC] \land v=1$
	У	Z	u	y≠1.2.3.5∧y≠1.2.3.4∧y=1.2.3.4∧z=[ABC]∧u=1
	У	Z	u	y≠1.2.3.5∧y≠1.2.3.4∧y=5.6.7.8∧z=[ABC]∧u≠1
	У	Z	V	y≠1.2.3.5∧y≠1.2.3.4∧y=1.2.3.4∧z=[ADC]∧v=1
	у	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [ADC] \land v \neq 1$

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Can	we genera			1.2.3.4	[ADC]	V	v = 1			
satisfies P_1 and P_3 simultaneously?							5.6.7.8	[ADC]	V	v ≠ 1

$P_1 \bowtie P_3$	dest	path	flag	
	1.2.3.4	Х	u	x=[ABC] ∧ u=1
	1.2.3.4	Х	V	$x=[ABC] \land x=[ADC] \land v=1$
Г	У	Z	u	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ABC] \land u = 1$
	У	Z	u	y≠1.2.3.5∧y≠1.2.3.4∧y=5.6.7.8∧z=[ABC]∧u≠1
	У	Z	V	$y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 1.2.3.4 \land z = [ADC] \land v = 1$
L	у	Z	V	y≠1.2.3.5∧y≠1.2.3.4∧y=5.6.7.8∧z=[ADC]∧v≠1

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 Traffic balance policy

¹ dest path ¹ dest path fla											
1.2.3.4 x x=[ABC] 1.2.3.4 [ABC] u	u = 1										
y z y≠1.2.3.5∧y≠1.2.3.4 [ABC] u	u ≠ 1										
Can we generate a new policy that [1.2.3.4 [ADC] v	v = 1										
satisfies P_1 and P_3 simultaneously? P_1 join P_3 [5.6.7.8 [ADC] v v \neq 1											
$P_1 \bowtie P_2$ dest nath flag											
$1.2.3.4$ X U X=[ABC] \land U=1											
1.2.3.4 x v $x = [ABC] \land x = [ADC] \land v = 1$											
y z u y≠1.2.3.5∧y≠1.2.3.4∧y=1.2.3.4∧z=[A	BC]∧u=1										
y z u y≠1.2.3.5∧y≠1.2.3.4∧y=5.6.7.8∧z=[A	BC]∧u≠1										
y z v y≠1.2.3.5∧y≠1.2.3.4∧y=1.2.3.4∧z=[A	DC]∧v=1										
y z v $y \neq 1.2.3.5 \land y \neq 1.2.3.4 \land y = 5.6.7.8 \land z = [A]$	DC]∧v≠1										

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P_1	dest	path			<i>P</i> ₃	dest	path	flag					
	1.2.3.4	Х	X=	[ABC]		1.2.3.4	[ABC]	U	u = 1				
	У	Z Y=	≠1.2.3.	5∧y≠1.2.3.4	ŀ	5.6.7.8	[ABC]	U	u ≠ 1				
Can	we genera	te a new	policy		1.2.3.4	[ADC]	V	v = 1					
satis	satisfies P_1 and P_3 simultaneously? P_1 join P_3 [5.6.7.8 [ADC] v v \neq 1												
$P_1 \bowtie$	$P_{\rm o}$ doct	nath	flog			Contra	dictory						
		patri	nay										
	1.2.3.4	k x	u	x=[ABC] \land	u=1								
	1.2.3.4	<u> х</u>		x=IABC1 ^		$1 \sqrt{-1}$							
				The same and the second s									
	y y	Z	u	y≠1.2.3.5∧	y≠1.2.3	3.4∧y≢1.	2.3.4∧z	=[ABC	}]∧u=1				
	y y	Z Z	u u	y≠1.2.3.5∧ y≠1.2.3.5∧	y≠1.2.3 y≠1.2.3	3.4∧y=1. 3.4∧y=5.	<mark>2.3.4∧z</mark> 6.7.8∧z	=[ABC z=[ABC	}]∧u=1)]∧u≠1				
	y y	Z Z 7	u u	$y \neq 1.2.3.5$ ∧ $y \neq 1.2.3.5$ ∧	y≠1.2.3 y≠1.2.3 y≠1.2.3	$3.4 \land y = 1.$ $3.4 \land y = 5.$ $4 \land y = 5.$	2.3.4∧z 6.7.8∧z 2.3.4∧z	=[ABC z=[ABC) ∧u=1)]∧u≠1				
	y y y	Z Z Z	u u V	y≠1.2.3.5∧ y≠1.2.3.5∧ y≠1.2.3.5∧	y≠1.2.3 y≠1.2.3 y≠1.2.3	3.4∧y=1. 3.4∧y=5. 3.4∧y=1.	2.3.4∧z 6.7.8∧z 2.3.4∧z	=[ABC z=[ABC =[ABC	}]∧u=1)]∧u≠1 }]∧v=1				
	y y y y y	Z Z Z Z	u u V V	y≠1.2.3.5∧ y≠1.2.3.5∧ y≠1.2.3.5∧ y≠1.2.3.5∧	y≠1.2.3 y≠1.2.3 y≠1.2.3 y≠1.2.3	3.4∧y=1. 3.4∧y=5. 3.4∧y=1. 3.4∧y=1.	2.3.4∧z 6.7.8∧z 2.3.4∧z 6.7.8∧z	=[ABC z=[ABC <u>=[ADC</u> z=[ADC)∧u=1)]∧u≠1)]∧v=1)]∧v≠1				

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$\underline{P_1 \bowtie P_3}$	dest	path	flag	
	1.2.3.4	Х	u	x=[ABC] ∧ u=1
	У	Ζ	u	y≠1.2.3.5∧y≠1.2.3.4∧y=5.6.7.8∧z=[ABC]∧u≠1
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P_1	dest	path			<i>P</i>	3	dest	path	flag	
	1.2.3.4	Х	Х	=[ABC]			1.2.3.4	[ABC]	U	u = 1
	У	Ζ	y≠1.2.3	8.5∧y≠1.2.3	3.4		5.6.7.8	[ABC]	U	u ≠ 1
Can	we gener	rate a n	ew policy		1.2.3.4	[ADC]	V	v = 1		
satis	fies P_1 an	d P_3 sir	multaneo		5.6.7.8	[ADC]	V	V ≠ 1		
Рм			la flaa				Redur	idancy		
1 1 M	a desi	pat	in flag							
	1.2.3.	.4 x	u	x=[ABC]	\land u=1					
	У	Z	u	y≠1.2.3.5	5∧y≠1./	2.3	.4∧y=5.	6 .7.8 ∧z	=[ABC	C]∧u≠1
	у	Z	V	y≠1.2.3.5	5∧y≠1.	2.3	.4∧y=5.	6. 7.8 ∧z	=[ADC	C]∧v≠1