



Building Java Programs

Chapter 13

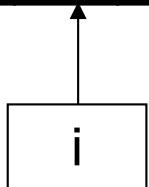
Searching and Sorting

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Sequential search

- **sequential search:** Locates a target value in an array/list by examining each element from start to finish.
 - How many elements will it need to examine?
 - Example: Searching the array below for the value **42**:

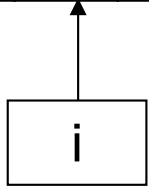
index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



- Notice that the array is sorted. Could we take advantage of this?

Sequential search

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



```
search(A[], thingToFind)
```

```
  start at the beginning
```

```
  for each item in A:
```

```
    if the item is what we're looking for:
```

```
      return its location
```

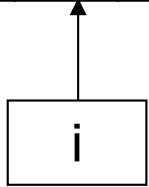
```
  if we got this far without returning already,
```

```
  what we're looking for isn't found
```

```
  return failure
```

Sequential search

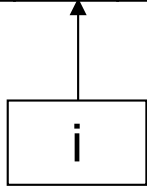
index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



```
public static int search(int A[], int thingToFind) {  
    for (int i=0; i<A.length; i++) {  
        if (A[i]==thingToFind) {  
            return i;  
        }  
    }  
    /* we didn't find it. return failure */  
    return -1;  
}
```

When Array Is Sorted

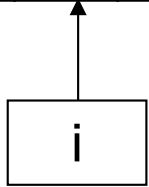
index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



- This array is already sorted
- Do we really need to go through entire thing before quitting?
 - Suppose we're searching for 38:
 - we know once we've reached $A[10]$ that we didn't find it

When Array Is Sorted

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



```
/* version 2. ONLY WORKS if A[] IS SORTED */
```

```
search(A[], thingToFind)
```

```
    start at the beginning
```

```
    for each item in A <= thingToFind:
```

```
        if the item is what we're looking for:
```

```
            return its location
```

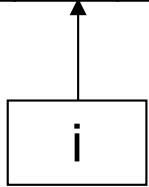
```
if we got this far without returning already,
```

```
what we're looking for isn't found.
```

```
return failure
```

When Array Is Sorted

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103



```
/* version 2. ONLY WORKS if A[] IS SORTED */  
public static int search(int A[], int thingToFind) {  
    for (int i=0; i<A.length && A[i]<=thingToFind; i++) {  
        if (A[i]==thingToFind) {  
            return i;  
        }  
    }  
    return -1;  
}
```

Searching a Sorted List

- What if we're ordering a pizza (and it's 1998)?
- Looking for phone number of Sammy's Pizza
- Start with the A's, then the B's, *etc.*?

Binary search (13.1)

- **binary search:** Locates a target value in a *sorted* array/list by successively eliminating half of the array from consideration.
 - How many elements will it need to examine?
 - Example: Searching the array below for the value **42**:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

Diagram illustrating a binary search on a sorted array. The array is shown with indices 0 to 16 and corresponding values. The value 42 is highlighted in yellow at index 10. Below the array, three boxes labeled 'min', 'mid', and 'max' have arrows pointing to the values -4, 30, and 103 respectively, indicating the current search range.

The Arrays class

- Class `Arrays` in `java.util` has many useful array methods:

Method name	Description
<code>binarySearch(array, value)</code>	returns the index of the given value in a <i>sorted</i> array (or <code>< 0</code> if not found)
<code>binarySearch(array, minIndex, maxIndex, value)</code>	returns index of given value in a <i>sorted</i> array between indexes <i>min</i> / <i>max</i> - 1 (<code>< 0</code> if not found)
<code>copyOf(array, length)</code>	returns a new resized copy of an array
<code>equals(array1, array2)</code>	returns <code>true</code> if the two arrays contain same elements in the same order
<code>fill(array, value)</code>	sets every element to the given value
<code>sort(array)</code>	arranges the elements into sorted order
<code>toString(array)</code>	returns a string representing the array, such as <code>"[10, 30, -25, 17]"</code>

- Syntax: `Arrays.methodName(parameters)`

Arrays.binarySearch

```
// searches an entire sorted array for a given value
// returns its index if found; a negative number if not found
// Precondition: array is sorted
```

```
Arrays.binarySearch(array, value)
```

```
// searches given portion of a sorted array for a given value
// examines minIndex (inclusive) through maxIndex (exclusive)
// returns its index if found; a negative number if not found
// Precondition: array is sorted
```

```
Arrays.binarySearch(array, minIndex, maxIndex, value)
```

- The `binarySearch` method in the `Arrays` class searches an array very efficiently if the array is sorted.
 - You can search the entire array, or just a range of indexes (useful for "unfilled" arrays such as the one in `ArrayIntList`)
 - If the array is not sorted, you may need to sort it first

Using `binarySearch`

```
// index    0  1  2  3  4  5  6  7  8  9  10  11  12  13  14  15
int[] a = {-4, 2, 7, 9, 15, 19, 25, 28, 30, 36, 42, 50, 56, 68, 85, 92};

int index = Arrays.binarySearch(a, 0, 16, 42); // index1 is 10
int index2 = Arrays.binarySearch(a, 0, 16, 21); // index2 is -7
```

- `binarySearch` returns the index where the value is found
- if the value is *not* found, `binarySearch` returns:
 - (`insertionPoint` + 1)
 - where `insertionPoint` is the index where the element *would* have been, if it had been in the array in sorted order.
 - To insert the value into the array, negate `insertionPoint` + 1

```
int indexToInsert21 = -(index2 + 1); // 6
```

Binary search code

```
// Returns the index of an occurrence of target in a,  
// or a negative number if the target is not found.  
// Precondition: elements of a are in sorted order  
public static int binarySearch(int[] a, int target) {  
    int min = 0;  
    int max = a.length - 1;  
  
    while (min <= max) {  
        int mid = (min + max) / 2;  
        if (a[mid] < target) {  
            min = mid + 1;  
        } else if (a[mid] > target) {  
            max = mid - 1;  
        } else {  
            return mid;    // target found  
        }  
    }  
  
    return -(min + 1);    // target not found  
}
```

Recursive binary search (13.3)

- Write a recursive `binarySearch` method.
 - If the target value is not found, return its negative insertion point.

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	10	15	20	22	25	30	36	42	50	56	68	85	92	103

```
int index = binarySearch(data, 42); // 10
int index2 = binarySearch(data, 66); // -14
```

Exercise solution

```
// Returns the index of an occurrence of the given value in
// the given array, or a negative number if not found.
// Precondition: elements of a are in sorted order
public static int binarySearch(int[] a, int target) {
    return binarySearch(a, target, 0, a.length - 1);
}

// Recursive helper to implement search behavior.
private static int binarySearch(int[] a, int target,
                                int min, int max) {
    if (min > max) {
        return -1;           // target not found
    } else {
        int mid = (min + max) / 2;
        if (a[mid] < target) {           // too small; go right
            return binarySearch(a, target, mid + 1, max);
        } else if (a[mid] > target) {    // too large; go left
            return binarySearch(a, target, min, mid - 1);
        } else {
            return mid;           // target found; a[mid] == target
        }
    }
}
}
```

Binary search and objects

- Can we `binarySearch` an array of Strings?
 - Operators like `<` and `>` do not work with `String` objects.
 - But we do think of strings as having an alphabetical ordering.
- **natural ordering**: Rules governing the relative placement of all values of a given type.
- **comparison function**: Code that, when given two values *A* and *B* of a given type, decides their relative ordering:
 - $A < B$, $A == B$, $A > B$

The compareTo method (10.2)

- The standard way for a Java class to define a comparison function for its objects is to define a `compareTo` method.
 - Example: in the `String` class, there is a method:

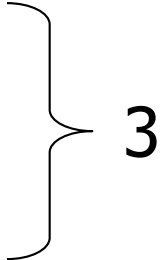
```
public int compareTo(String other)
```
- A call of `A.compareTo(B)` will return:
 - a value < 0 if **A** comes "before" **B** in the ordering,
 - a value > 0 if **A** comes "after" **B** in the ordering,
 - or 0 if **A** and **B** are considered "equal" in the ordering.

Runtime Efficiency (13.2)

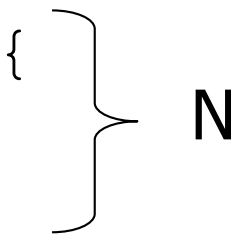
- **efficiency:** A measure of the use of computing resources by code.
 - can be relative to speed (time), memory (space), etc.
 - most commonly refers to run time
- Assume the following:
 - Any single Java statement takes the same amount of time to run.
 - A method call's runtime is measured by the total of the statements inside the method's body.
 - A loop's runtime, if the loop repeats N times, is N times the runtime of the statements in its body.

Efficiency examples

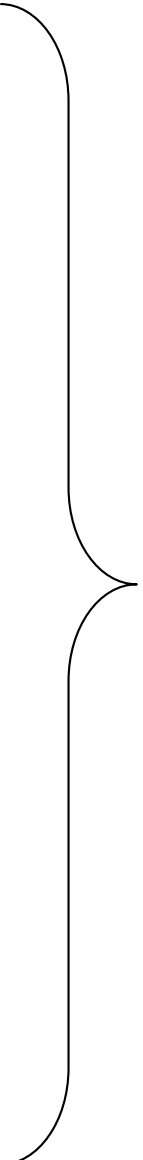
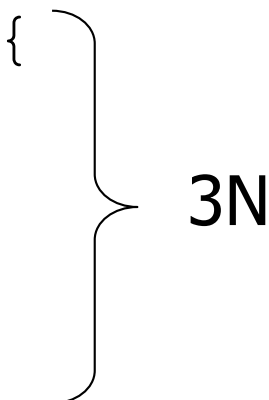
```
statement1;  
statement2;  
statement3;
```



```
for (int i = 1; i <= N; i++) {  
    statement4;  
}
```



```
for (int i = 1; i <= N; i++) {  
    statement5;  
    statement6;  
    statement7;  
}
```



$4N + 3$

Efficiency examples 2

```
for (int i = 1; i <= N; i++) {  
    for (int j = 1; j <= N; j++)  
        statement1;  
}
```

} N^2

```
for (int i = 1; i <= N; i++) {  
    statement2;  
    statement3;  
    statement4;  
    statement5;  
}
```

} $4N$

} $N^2 + 4N$

- How many statements will execute if $N = 10$? If $N = 1000$?

Algorithm growth rates (13.2)

- We measure runtime in proportion to the input data size, N .
 - **growth rate**: Change in runtime as N changes.
- Say an algorithm runs $0.4N^3 + 25N^2 + 8N + 17$ statements.
 - Consider the runtime when N is *extremely large* .
 - We ignore constants like 25 because they are tiny next to N .
 - The highest-order term (N^3) dominates the overall runtime.
 - We say that this algorithm runs "on the order of" N^3 .
 - or **$O(N^3)$** for short ("Big-Oh of N cubed")

Complexity classes

- **complexity class:** A category of algorithm efficiency based on the algorithm's relationship to the input size N .

Class	Big-Oh	If you double N , ...	Example
constant	$O(1)$	unchanged	10ms
logarithmic	$O(\log_2 N)$	increases slightly	175ms
linear	$O(N)$	doubles	3.2 sec
log-linear	$O(N \log_2 N)$	slightly more than doubles	6 sec
quadratic	$O(N^2)$	quadruples	1 min 42 sec
cubic	$O(N^3)$	multiplies by 8	55 min
...
exponential	$O(2^N)$	multiplies drastically	$5 * 10^{61}$ years

Binary search (13.1, 13.3)

- **binary search** successively eliminates half of the elements.
 - *Algorithm:* Examine the middle element of the array.
 - If it is too big, eliminate the right half of the array and repeat.
 - If it is too small, eliminate the left half of the array and repeat.
 - Else it is the value we're searching for, so stop.
 - Which indexes does the algorithm examine to find value **22**?
 - What is the runtime complexity class of binary search?

<i>index</i>	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
<i>value</i>	-4	-1	0	2	3	5	6	8	11	14	22	29	31	37	56

Binary search runtime

- For an array of size N , it eliminates $\frac{1}{2}$ until 1 element remains.
 $N, N/2, N/4, N/8, \dots, 4, 2, 1$
 - How many divisions does it take?

- Think of it from the other direction:
 - How many times do I have to multiply by 2 to reach N ?
 $1, 2, 4, 8, \dots, N/4, N/2, N$
 - Call this number of multiplications "x".

$$2^x = N$$

$$\mathbf{x = \log_2 N}$$

- Binary search is in the **logarithmic** complexity class.

Sorting

- **sorting**: Rearranging the values in an array or collection into a specific order (usually into their "natural ordering").
 - one of the fundamental problems in computer science
 - can be solved in many ways:
 - there are many sorting algorithms
 - some are faster/slower than others
 - some use more/less memory than others
 - some work better with specific kinds of data
 - some can utilize multiple computers / processors, ...
 - *comparison-based sorting* : determining order by comparing pairs of elements:
 - `<`, `>`, `compareTo`, ...

Sorting methods in Java

- The `Arrays` class in `java.util` has a static method `sort` that sorts the elements of an array

```
String[] words = {"foo", "bar", "baz", "ball"};
Arrays.sort(words);
System.out.println(Arrays.toString(words));
// [ball, bar, baz, foo]
```

Selection sort

- **selection sort:** Orders a list of values by repeatedly putting the smallest or largest unplaced value into its final position.

The algorithm:

- Look through the list to find the smallest value.
- Swap it so that it is at index 0.
- Look through the list to find the second-smallest value.
- Swap it so that it is at index 1.
- ...
- Repeat until all values are in their proper places.

Selection sort example

- Initial array:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	22	18	12	-4	27	30	36	50	7	68	91	56	2	85	42	98	25

- After 1st, 2nd, and 3rd passes:

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	18	12	22	27	30	36	50	7	68	91	56	2	85	42	98	25

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	12	22	27	30	36	50	7	68	91	56	18	85	42	98	25

index	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
value	-4	2	7	22	27	30	36	50	12	68	91	56	18	85	42	98	25

Selection sort code

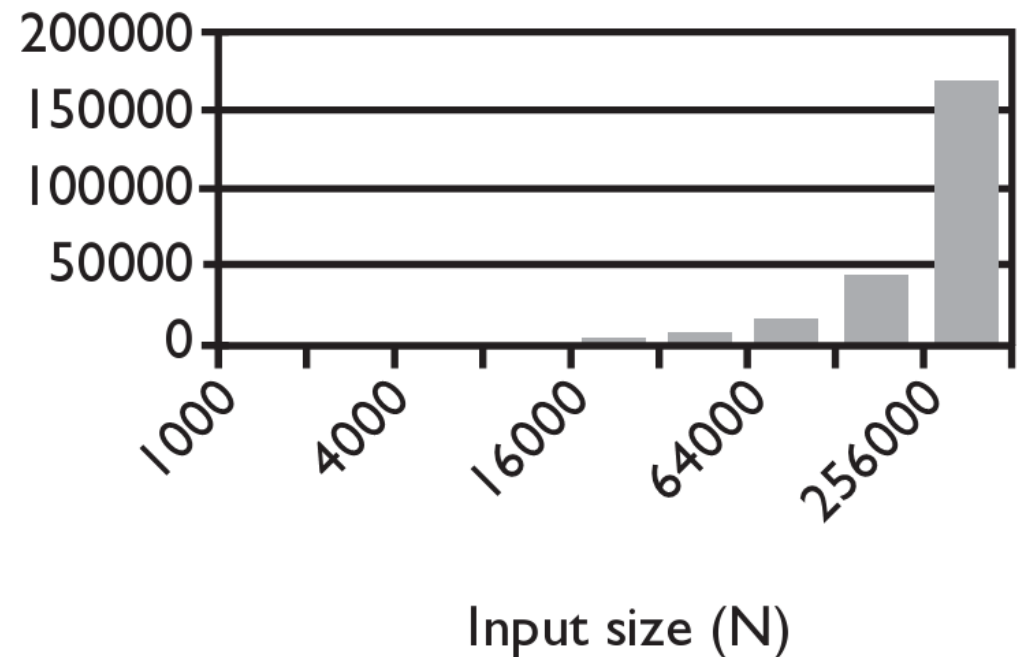
```
// Rearranges the elements of a into sorted order using  
// the selection sort algorithm.
```

```
public static void selectionSort(int[] a) {  
    for (int i = 0; i < a.length - 1; i++) {  
        // find index of smallest remaining value  
        int min = i;  
        for (int j = i + 1; j < a.length; j++) {  
            if (a[j] < a[min]) {  
                min = j;  
            }  
        }  
  
        // swap smallest value its proper place, a[i]  
        swap(a, i, min);  
    }  
}
```

Selection sort runtime (Fig. 13.6)

- What is the complexity class (Big-Oh) of selection sort?

N	Runtime (ms)
1000	0
2000	16
4000	47
8000	234
16000	657
32000	2562
64000	10265
128000	41141
256000	164985



Merge sort

- **merge sort:** Repeatedly divides the data in half, sorts each half, and combines the sorted halves into a sorted whole.

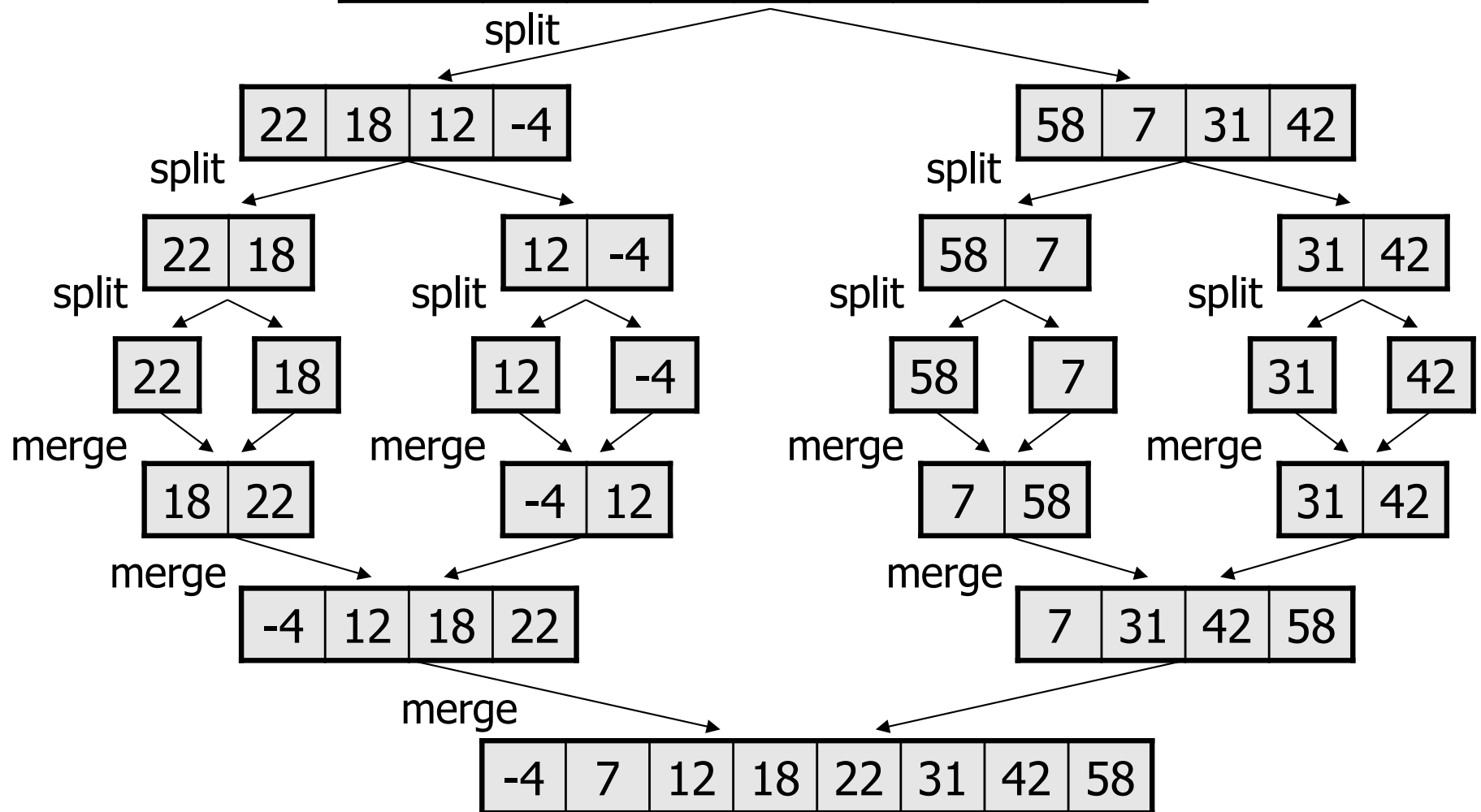
The algorithm:

- Divide the list into two roughly equal halves.
- Sort the left half.
- Sort the right half.
- Merge the two sorted halves into one sorted list.

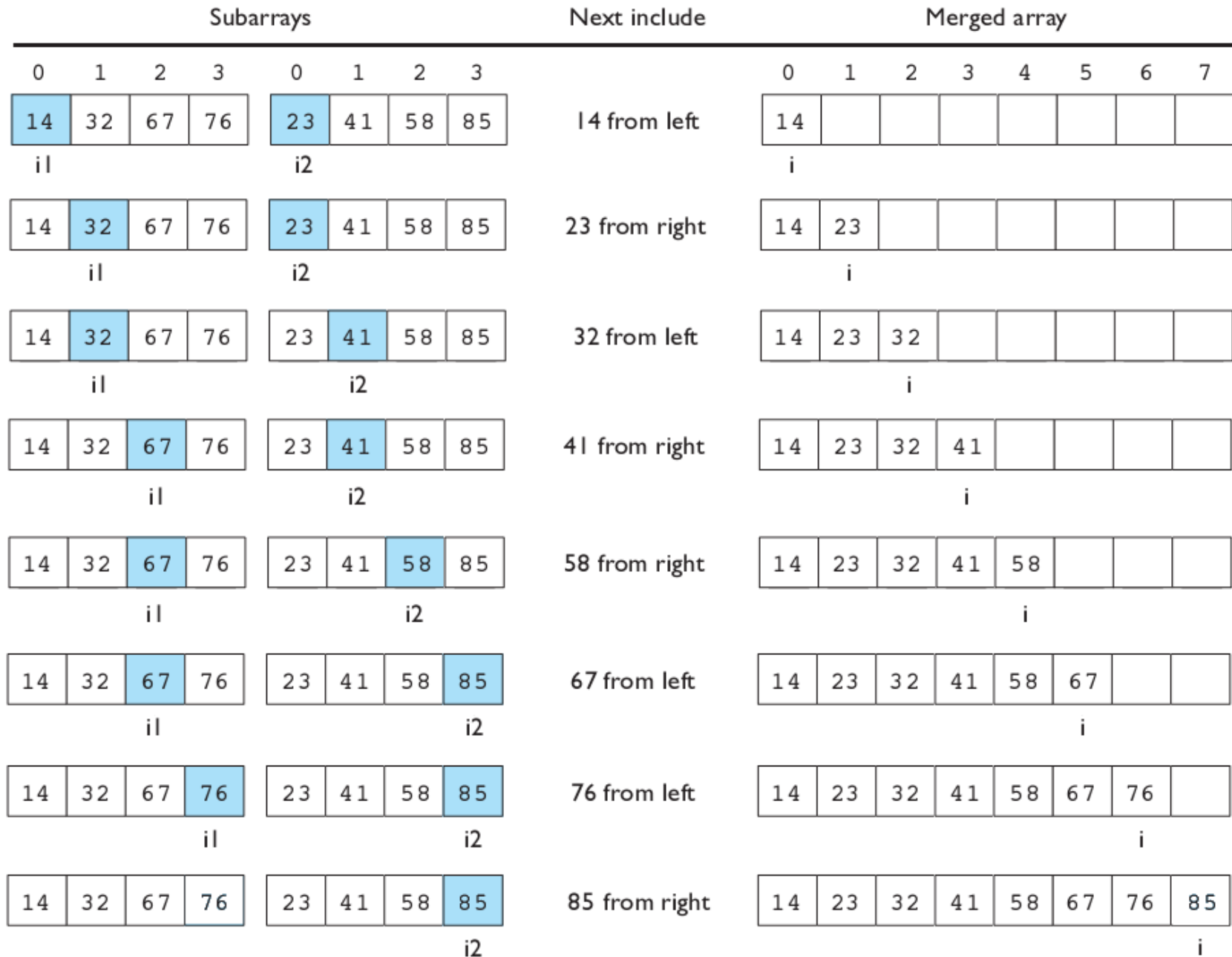
- Often implemented recursively.
- An example of a "divide and conquer" algorithm.
 - Invented by John von Neumann in 1945

Merge sort example

index	0	1	2	3	4	5	6	7
value	22	18	12	-4	58	7	31	42



Merging sorted halves



Merge halves code

```
// Merges the left/right elements into a sorted result.
// Precondition: left/right are sorted
public static void merge(int[] result, int[] left,
                        int[] right) {
    int i1 = 0;    // index into left array
    int i2 = 0;    // index into right array

    for (int i = 0; i < result.length; i++) {
        if (i2 >= right.length ||
            (i1 < left.length && left[i1] <= right[i2])) {
            result[i] = left[i1];    // take from left
            i1++;
        } else {
            result[i] = right[i2];    // take from right
            i2++;
        }
    }
}
```

Merge sort code

```
// Rearranges the elements of a into sorted order using
// the merge sort algorithm.
public static void mergeSort(int[] a) {
    // split array into two halves
    int[] left  = Arrays.copyOfRange(a, 0, a.length/2);
    int[] right = Arrays.copyOfRange(a, a.length/2, a.length);

    // sort the two halves
    ...

    // merge the sorted halves into a sorted whole
    merge(a, left, right);
}
```

Merge sort code 2

```
// Rearranges the elements of a into sorted order using
// the merge sort algorithm (recursive).
public static void mergeSort(int[] a) {
    if (a.length >= 2) {
        // split array into two halves
        int[] left  = Arrays.copyOfRange(a, 0, a.length/2);
        int[] right = Arrays.copyOfRange(a, a.length/2, a.length);

        // sort the two halves
        mergeSort(left);
        mergeSort(right);

        // merge the sorted halves into a sorted whole
        merge(a, left, right);
    }
}
```

Merge sort runtime

- What is the complexity class (Big-Oh) of merge sort?

N	Runtime (ms)
1000	0
2000	0
4000	0
8000	0
16000	0
32000	15
64000	16
128000	47
256000	125
512000	250
1e6	532
2e6	1078
4e6	2265
8e6	4781
1.6e7	9828
3.3e7	20422
6.5e7	42406
1.3e8	88344

