

1. SELECT name FROM instructor
 WHERE
 dept_name = "Biology";

$\Pi_{name} (\sigma_{dept_name = \text{"Biology"}} (\text{instructor}))$

$\{t \mid \exists s \in \text{instructor} (t[\text{name}] = s[\text{name}] \wedge s[\text{dept_name}] = \text{"Biology"})\}$

$\{ \langle n \rangle \mid \exists i, d, s (\langle i, n, d, s \rangle \in \text{instructor} \wedge d = \text{'Biology'}) \}$

2.
 SELECT title FROM course
 WHERE dept_name = "Comp. Sci." AND credits = 3;

$\Pi_{title} (\sigma_{dept_name = \text{"Comp. Sci."} \wedge credits=3} (\text{course}))$

$\{t \mid \exists s \in \text{course} (t[\text{title}] = s[\text{title}] \wedge s[\text{dept_name}] = \text{"Comp. Sci."} \wedge credits = 3)\}$

$\{ \langle t \rangle \mid \exists c, d, cr (\langle c, t, d, cr \rangle \in \text{course} \wedge d = \text{'Comp. Sci.'} \wedge cr = 3) \}$

3.

SELECT course.course_id, title
 FROM course, takes
 WHERE course.course_id = takes.course_id
 AND takes.ID = 12345;

$\Pi_{course.course_id, course.title} (\sigma_{takes.ID = 12345} (\text{course} \bowtie_{course.course_id=takes.course_id} \text{takes}))$

$\{t \mid \exists s \in \text{course} (t[\text{course_id}] = s[\text{course_id}] \wedge t[\text{title}] = s[\text{title}] \wedge \exists u \in \text{takes} (u[\text{course_id}] = s[\text{course_id}] \wedge u[\text{ID}] = 12345))\}$

$$\{\langle c, t \rangle \mid \exists d, cr(\langle c, t, d, cr \rangle \in \text{course} \wedge \exists i, ci, si, s, y, g(\langle i, ci, si, s, y, g \rangle \in \text{takes} \wedge i = 12362 \wedge c = ci))\}$$

4.

```
SELECT SUM(course.credits)
FROM course, takes
WHERE course.course_id = takes.course_id
AND takes.ID = 17424;
```

We did not cover aggregate queries in RA, Tuple Calculus, and Domain Calculus.

5.

```
SELECT takes.ID, SUM(course.credits)
FROM course, takes
WHERE course.c
ourse_id = takes.course_id
GROUP BY takes.ID
```

We did not cover aggregate queries in RA, Tuple Calculus, and Domain Calculus.

6.

```
SELECT DISTINCT S.name
FROM takes T, course C, student S
WHERE C.dept_name = 'Comp. Sci.' and T.course_id = C.course_id and T.ID = S.ID
```

$$\Pi_{name}(\sigma_{course.dept_name='Comp. Sci.' \wedge takes.course_id=course.course_id \wedge takes.ID=student.ID} (course \times takes) \times student)$$

$$\{t \mid \exists s \in \text{student}(t[\text{name}] = s[\text{name}] \wedge \exists u \in \text{takes}(s[\text{ID}] = u[\text{ID}] \wedge \exists v \in \text{course}(u[\text{course_id}] = v[\text{course_id}] \wedge v[\text{dept_name}] = \text{'Comp. Sci. '}))\}$$

$$\{ \langle n \rangle \mid \exists i, d, t (\langle i, n, d, t \rangle \in \text{student} \wedge \exists id, c, si, s, y, g (\langle id, c, si, s, y, g \rangle \in \text{takes} \wedge i = id \wedge \exists ci, ti, dn, cr (\langle ci, ti, dn, cr \rangle \in \text{course} \wedge c = ci \wedge dn = \text{'Comp. Sci. '}))))) \}$$

7.

select id from instructor except
(select teaches.id from teaches, instructor where teaches.id = instructor.id)

$$\Pi_{ID}(\text{instructor}) - \Pi_{ID}(\text{teaches})$$

$$\{t \mid \exists s \in \text{instructor}(t[ID] = s[ID] \wedge \forall u \in \text{teaches}(s[ID] \neq u[ID]))\}$$

$$\{ \langle i \rangle \mid \exists n, d, s (\langle i, n, d, s \rangle \in \text{instructor} \wedge \forall id, c, si, se, y (\langle id, c, si, se, y \rangle \in \text{teaches} \wedge i \neq id)) \}$$

Intermediate SQL queries

1. SELECT min(enrollment) as min_enrol, max(enrollment) as maxEnrol
FROM (SELECT count(*) as enrollment
FROM takes
GROUP BY course_id, sec_id, semester, year) as countBySection

We did not cover aggregate queries in RA, Tuple Calculus, and Domain Calculus. That being said, MAX and MIN can be computed with the standard calculus.

2.
SELECT course_id, sec_id, semester, year, count(*) AS MaxEnrollment
FROM takes
GROUP BY course_id, sec_id, semester, year

```
HAVING count(*) = (SELECT MAX(count)
FROM (SELECT COUNT(ID) AS count
FROM
takes
GROUP BY course_id, sec_id, semester, year)
AS studentCount)
```

3.

```
SELECT * FROM course
WHERE course_id LIKE "CS-1%";
```

Advanced SQL queries

1. CREATE VIEW faculty AS (SELECT ID, name, dept_name FROM instructor);

Views can be expressed with the assignment operator.

$$\text{faculty} \leftarrow \Pi_{\text{ID, name, dept_name}}(\text{instructor})$$

$$\text{faculty} \leftarrow \{t \mid \exists s \in \text{instructor} (t[\text{ID}] = s[\text{ID}] \wedge t[\text{name}] = s[\text{name}] \wedge t[\text{dept_name}] = s[\text{dept_name}])\}$$

$$\text{faculty} \leftarrow \{\langle i, n, d \rangle \mid \langle i, n, d, s \rangle \in \text{instructor}\}$$

2.

```
CREATE VIEW CSinstructors AS (SELECT *
FROM instructor
WHERE dept_name = "Comp. Sci.");
```

$$\text{CSinstructors} \leftarrow \sigma_{\text{dept_name} = \text{"Comp. Sci."}}(\text{instructor})$$

$$\text{CSInstructor} \leftarrow \{t \mid t \in \text{instructor} \wedge t[\text{dept_name}] = \text{'Comp. Sci.'}\}$$

$$\text{CSinstructors} \leftarrow \{\langle i, n, d, s \rangle \mid \langle i, n, d, s \rangle \in \text{instructor} \wedge d = \text{"Comp. Sci."}\}$$

3.